

# **Energy Aware Consolidation for Dynamic Resource Applications**

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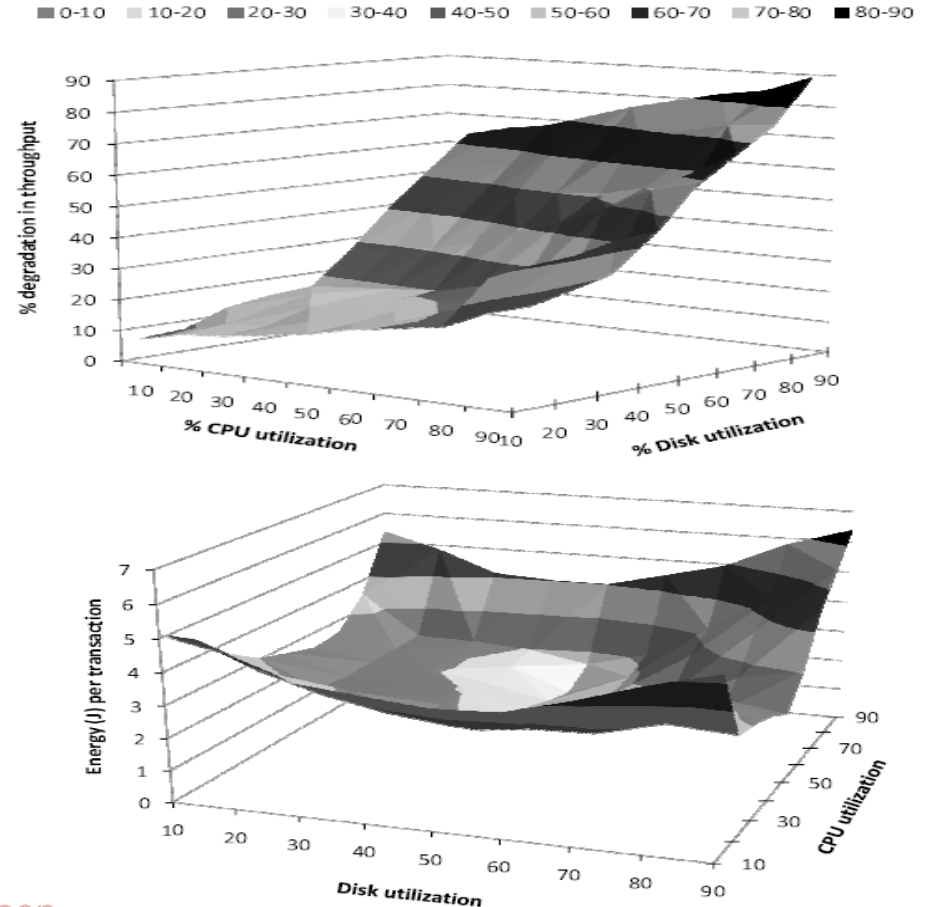
- Introduction
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- The goal is to find an method for allocating client applications on a set of common servers in a manner that minimizes energy used.
- The method proposed is an extension of the one described in the original paper *Energy Aware Consolidation in Cloud Computing*
- The original method is limited to allocating applications with static resource utilizations.
- Our extended method is capable of allocating applications with dynamic resource footprint.

## What is Energy Aware Consolidation?

- Running many dissimilar client applications on the same server cluster.
- In other words running multiple data center applications on a common set of servers.
- This allows for the consolidation of application workloads on a smaller number of servers that may be kept better utilized.

- Effective consolidation is not as trivial as packing the maximum workload in the smallest number of servers.
- Keeping resources at 100% utilization is not energy efficient.
- Goal is to minimize the energy used per unit service.
- Energy optimal resource point (vector) is 50% hard disk, and 70% cpu utilization.



- The allocation method is almost identical that the method proposed in the original paper, except for the \*.
1. For each server type, the server is subjected to various client application workloads in order to obtain Energy consumption data. From this data we obtain:
    - a. the Energy vs. resource utilization relationship for the server(s).
    - b. the Energy optimal resource point. ex) [70%, 50%]
    - c. \* the Energy vs. resource usage relationship for the applications.
      - i. either tabulated data or function.
  2. \* Allocate incoming client applications according to the Dynamic Allocation Algorithm.

- Resource utilization and consumption by servers and client applications are conveyed with resource vectors.
  - ex) 50% CPU utilization and 50% Hard Disk utilization would be written  $(0.5, 0.5) = \mathbf{r}$
- Distance between two resource points is given by the euclidean distance
 
$$\delta_e = \sqrt{(x_1^2 + x_2^2 + \dots + x_n^2)}$$
  - ex)  $\mathbf{r}_1 = (0.3, 0.8)$ ,  $\mathbf{r}_2 = (0.7, 0.5)$ ,  $\delta_e(\mathbf{r}_2 - \mathbf{r}_1) = \sqrt{(0.7 - 0.3)^2 + (0.5 - 0.8)^2} = 0.5$
- $\mathbf{r}(t)$  indicates a resource vector is dependent on time.
- $\max\{ \mathbf{r}(t) \}$  returns a set of resource vector points that are local maxima.
- $\mathbf{s}^* = (0.7, 0.5)$  = optimal resource vector for server  $s$ .
- $|\cdot|$ , is the cardinality of set, it give the number of elements in the set.

- Consider a workload  $\mathbf{W}$ , and a set of servers  $\mathbf{S}$  the allocated the algorithms as follows.

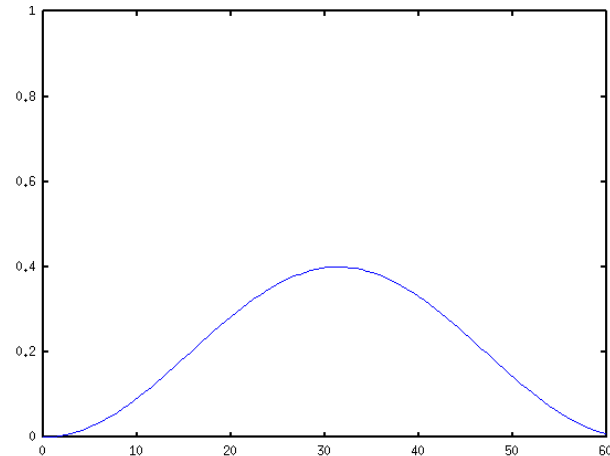
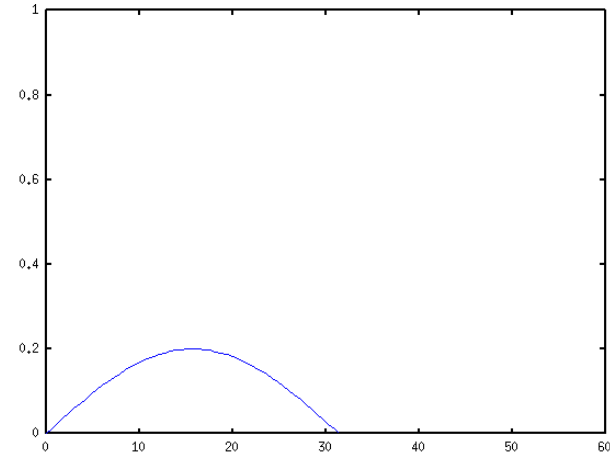
## Dynamic Allocation Algorithm

1. Let score[s] be a map of mean distances of maxima for a server.
2. Foreach client application  $w$  in  $\mathbf{W}$ 
  - a. Foreach server  $s$  in  $\mathbf{S}$ 
    - i. Compute  $\max\{s\} \cup \max\{w\} = \max\_set$
    - ii. Foreach  $t$  in  $\max\_set$ 
      1. IF  $r_s(t) + r_w(t) \leq s^*$  THEN  $\text{score}[s] += \delta_e(s^* - (r_s(t) + r_w(t)))$
      2. ELSE remove  $s$  from score and break.
    - iii.  $\text{score}[s] /= |s|$
  - b. Assign  $w$  to the maximum scoring server.



# Dynamic Allocation Algorithm Example Part 1

- We have a workload consisting of two applications with resource utilizations of the following form:
  - App 1 has a resource relationship given on the right.
    - The app has a maxima at  $t = 16$  and resource utilization of  $r_1$
  - App 2 has a resource relationship given on the right.
    - The app has a maxima at  $t = 31$  and resource utilization of  $r_2$
- The graphs are for both CPU and HD utilization vs. time for each application.



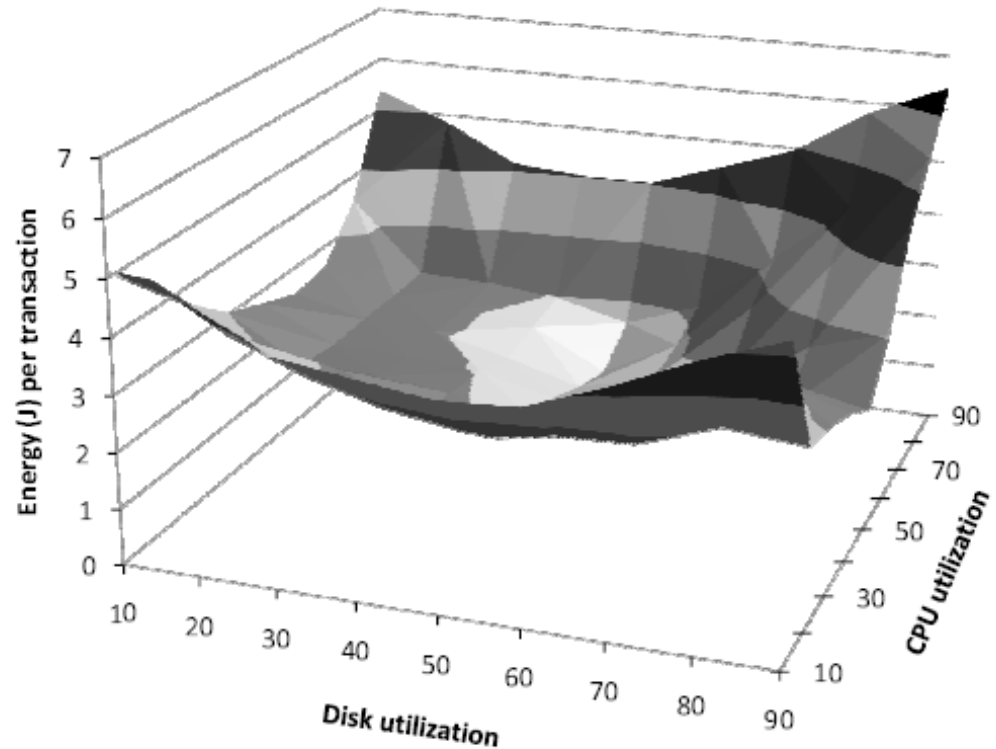
## Dynamic Allocation Algorithm Example Part 2

- Assume we have two servers, A and B that are idling at 10% utilization, so  $\mathbf{s}_a = (0.1, 0.1)$ ,  $\mathbf{s}_b = (0.2, 0.2)$ .
  - Both servers has optimal resource  $\mathbf{s}^*(t) = (0.7, 0.5)$
- First we compute the max\_set for  $\max\{\mathbf{s}_a\} \cup \max\{\mathbf{r}_1\} = \{t=16\}$
- Now we check if the  $\mathbf{r}_1(t=16) + \mathbf{s}_a(t=16) < \mathbf{s}^*(t)$ ,  $(0.2, 0.2) + (0.1, 0.1) < (0.7, 0.5)$
- Next we add the euclidean distance between the target
  - $\delta_e(\mathbf{s}_a(t) + \mathbf{r}_1(t), \mathbf{s}^*(t)) = \sqrt{(0.3 - 0.7)^2 + (0.3 - 0.5)^2} = 0.447 \rightarrow \text{score}[\mathbf{s}_a]$
  - $\delta_e(\mathbf{s}_b(t) + \mathbf{r}_1(t), \mathbf{s}^*(t)) = \sqrt{(0.4 - 0.7)^2 + (0.4 - 0.5)^2} = 0.316 \rightarrow \text{score}[\mathbf{s}_b]$
- Now we assign the application to the server with largest score, which is  $\mathbf{s}_a$ .

## Dynamic Allocation Algorithm Example Part 3

- Assume we have two servers, A and B that are idling at 10% utilization, so  $\mathbf{s}_a = (0.1, 0.1)$ ,  $\mathbf{s}_b = (0.2, 0.2)$ .
  - Both servers has optimal resource  $\mathbf{s}^*(t) = (0.7, 0.5)$
- First we compute the max\_set for  $\max\{\mathbf{s}_a\} \cup \max\{\mathbf{r}_1\} = \{t=16\}$
- Now we check if the  $\mathbf{r}_1(t=16) + \mathbf{s}_a(t=16) < \mathbf{s}^*(t)$ ,  $(0.2, 0.2) + (0.1, 0.1) < (0.7, 0.5)$
- Next we add the euclidean distance between the target
  - $\bar{\delta}_e(\mathbf{s}_a(t) + \mathbf{r}_1(t), \mathbf{s}^*(t)) = \sqrt{(0.3 - 0.7)^2 + (0.3 - 0.5)^2} = 0.447 \rightarrow \text{score}[\mathbf{s}_a]$
  - $\bar{\delta}_e(\mathbf{s}_b(t) + \mathbf{r}_1(t), \mathbf{s}^*(t)) = \sqrt{(0.4 - 0.7)^2 + (0.4 - 0.5)^2} = 0.316 \rightarrow \text{score}[\mathbf{s}_b]$
- Now we assign the application to the server with largest score, which is  $\mathbf{s}_a$ .

- In order to validate the algorithm, a simulation of the original experimental conditions was constructed.
- The simulation consists of four servers processing a workload of 8 client applications.
- The Power vs. resource utilization for the simulated servers was obtained from the Energy per transaction graph shown below.



- We want **Power** as a function of cpu and hard disk utilization.
- The original experiment gives data for the total energy used over a 60second period for constant utilizations.
- We use the thermodynamic relationship **Power \* Time = Work**, and make the approximation that **Power** is dependent only on the utilization rates. This approx. is only true for small time intervals.
  - $c$  = CPU utilization rate, and  $h$  = Hard Disk utilization rate.
- $E(c,h) / N$  is given, and we know that  $P(c,h) * 60 = E(c,h)/N$ . Solving for  $P(c,h)$  yields  $P(c,h) = E(c,h) / ( N * 60)$ .
- The **Power** as a function of resource utilizations is critical for the simulation, since  $P(c,h) * dt = \text{Energy used for the small time interval } dt$ .

- Assumption:
  - Profiling data has already been obtained for all client applications.
  - All servers are homogenous with identical **Power** vs. resource relationship.
  - All servers are awake.
- Simulation Description:
  - Processing a workload consisting of dummy applications with variable resource utilizations.
  - At simulation start, the allocation algorithm begins delegating applications to each of the servers.
  - Server processing takes place in discrete time steps  $\sim 0.1$  (ms).
    - The energy used during each timestep is calculated using
      - $P(c,h) * (\text{timestep})$ , and added to the total.
  - The simulation concludes when the workload has been processed.

- The simulation was run for three allocation algorithms using a workloads of varying utilization functions.
  - Dynamic algorithm - algorithm for allocating applications with dynamic resource utilizations.
  - Original Algorithm - algorithm for allocating applications with static resource utilizations.
  - Optimal Algorithm - algorithm that allocates applications according to a aprior calculated ordering that is optimal. This is algorithm is used for validation only.

- For a workload mix with a constant utilization of 180% CPU and 180% HD.
  - Original Algorithm ~ 4.12 [J/Op]
  - Dynamic Algorithm ~ 4.12 [J/Op]
  - Optimal Algorithm ~ 4.1 [J/Op]
- For a workload with mix of constant and sinusoidal utilizations.
  - Original Algorithm ~ 6.18 [J/Op]
  - Dynamic Algorithm ~ 4.5 [J/Op]



- Srikantaiah S et al. (2008). Energy aware consolidation for cloud computing. In: Proc of HotPower

