

Policy-preserving Middlebox Placement in SDN-Enabled Data Centers

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Some slides are from

www.cs.berkeley.edu/~randy/Courses/CS268.F08/lectures/22-policy_switching.ppt, and

www.cs.yale.edu/homes/yu-minlan/talk/sigcomm13.pptx

Overview

- What is middlebox?
- What is SDN (Software Defined Network) and NFV (Network Function Virtualization)?
- Policy-preserving middlebox placement problem in data centers
 - Problems and preliminary solutions
- Conclusions

Middleboxes

- A **middlebox**, or network appliance, is a computer networking device that transforms, inspects, filters, or otherwise manipulates traffic for purposes other than packet forwarding.
 - Intermediaries in-between the communicating hosts
 - Often without knowledge of one or both parties
- Examples
 - Network address translators
 - Firewalls
 - Load balancers
 - Intrusion detection systems
 - Transparent Web proxy caches

Problem: Middleboxes are hard to deploy

- Place on network path



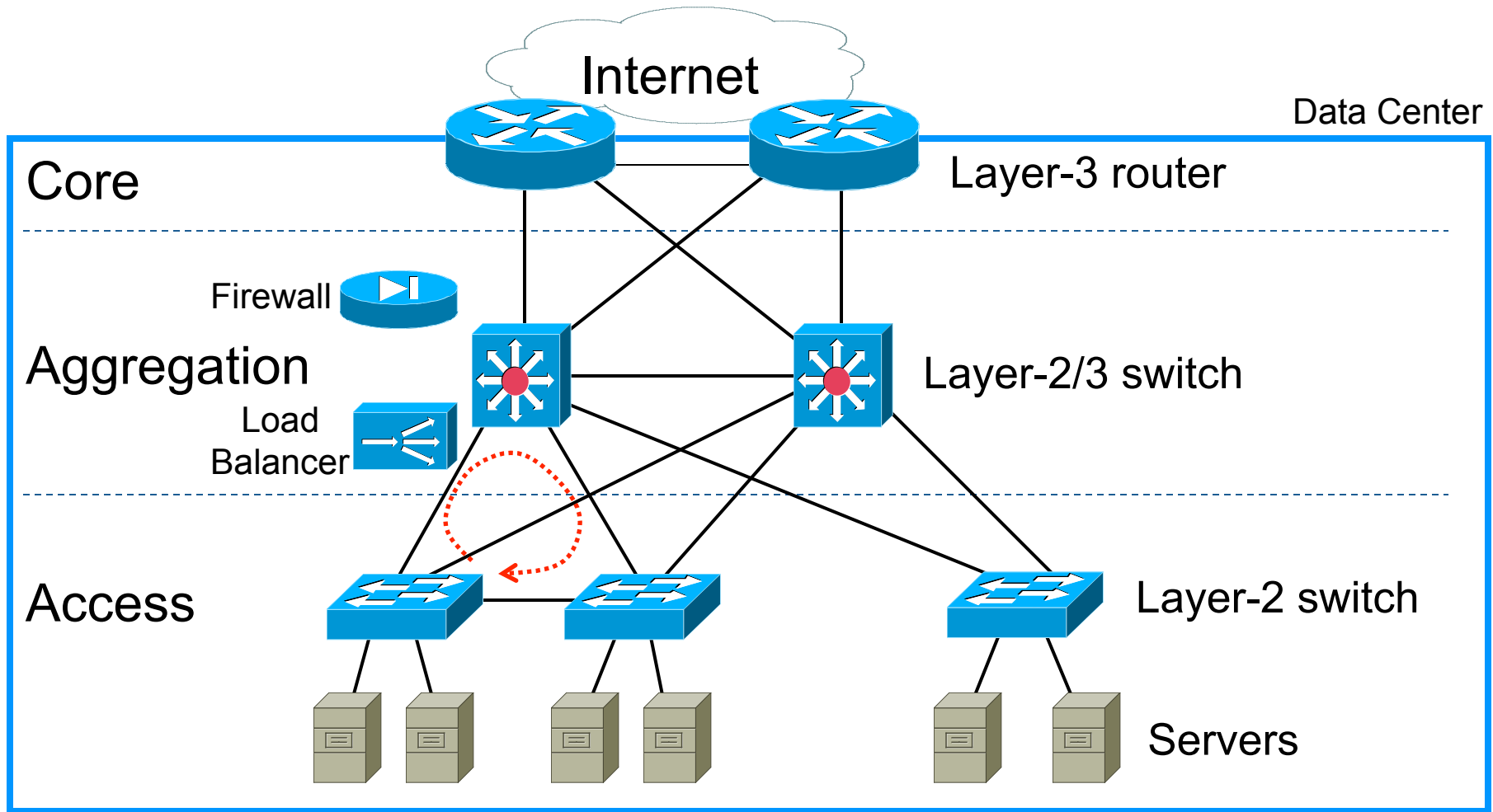
- On path placement fails to achieve

Flexibility | (Re)configurable network topology

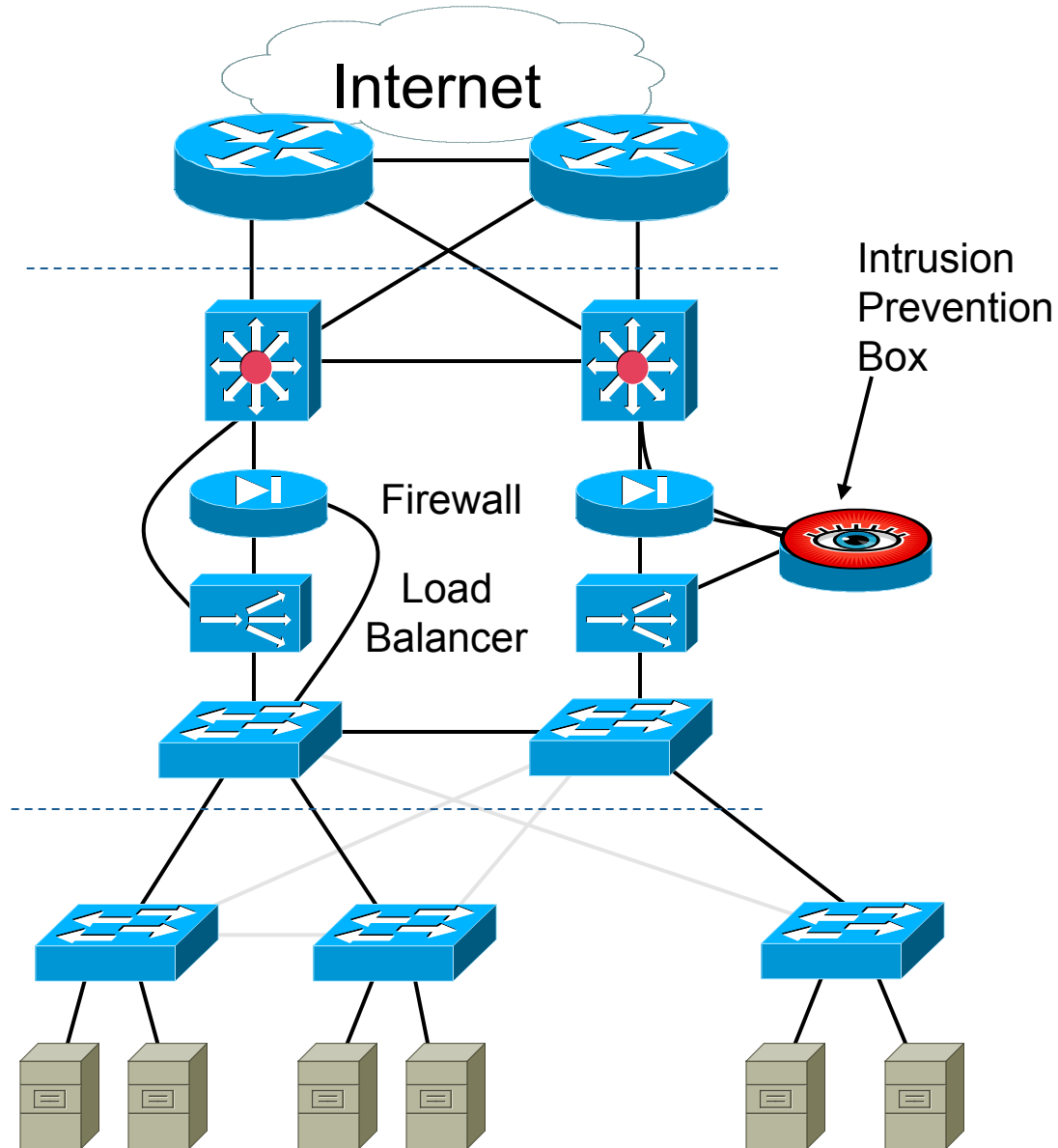
Efficiency | No middlebox resource wastage

Correctness | Guaranteed middlebox traversal

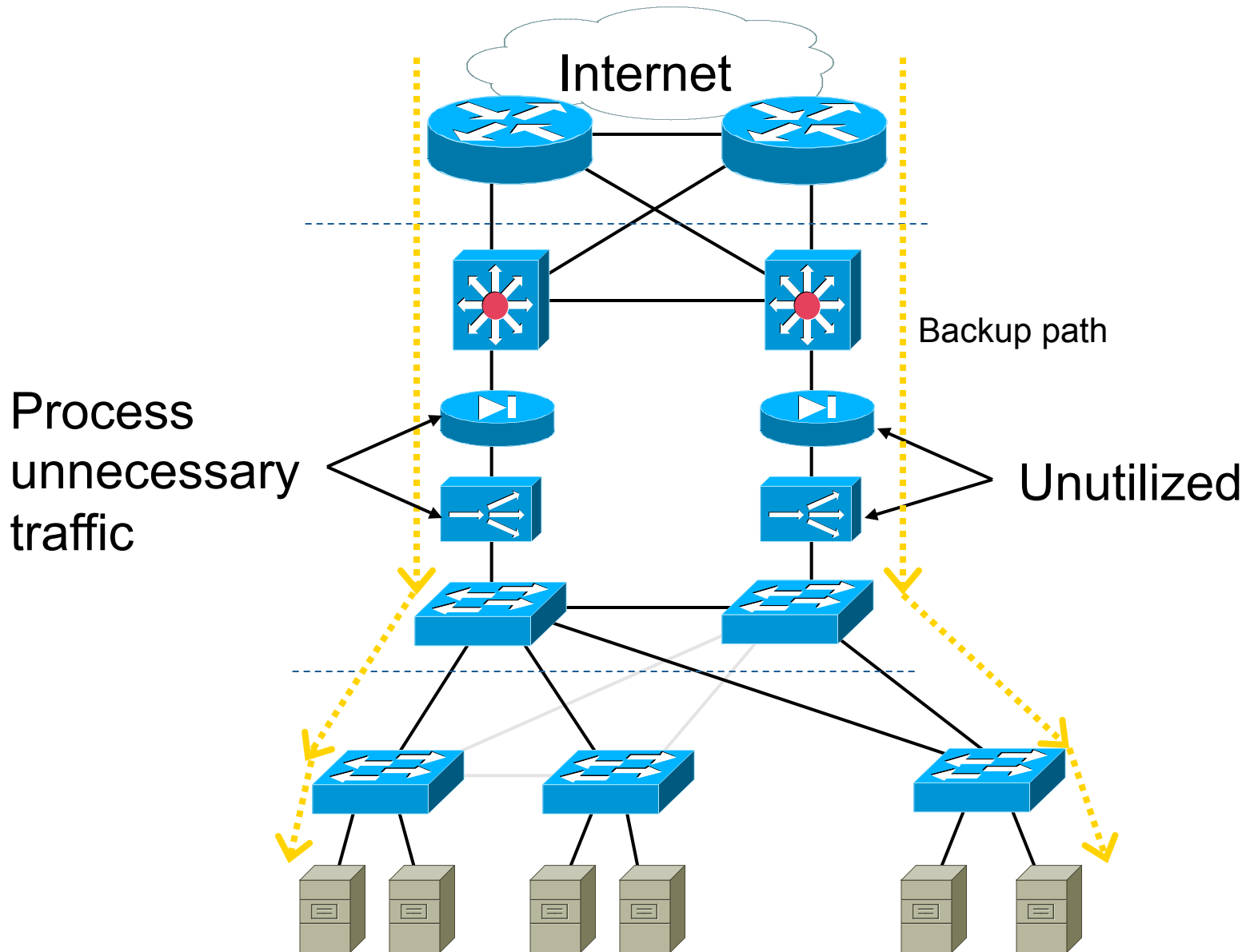
Common data center topology



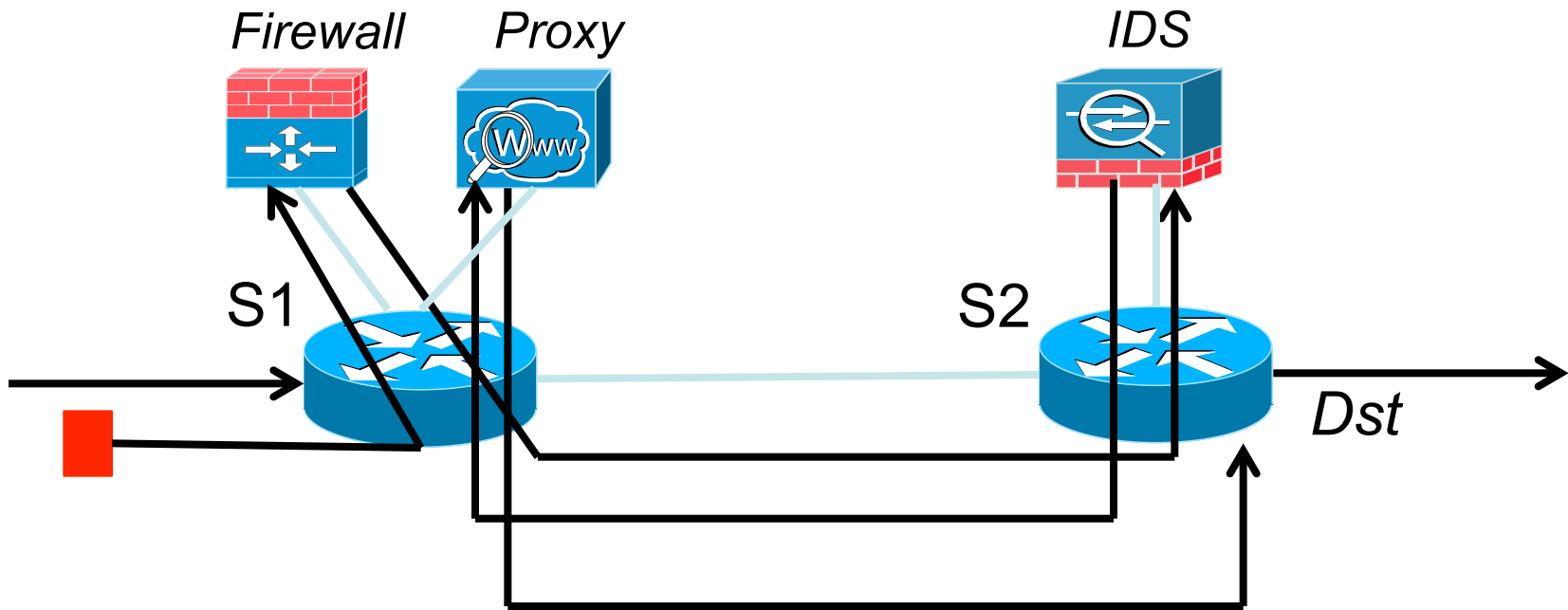
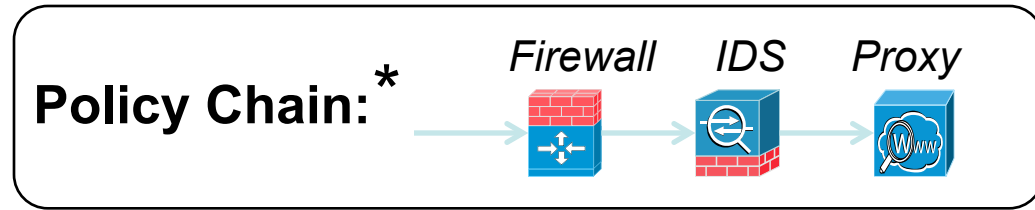
Inflexible topology



Inefficient - middlebox resource wastage



Policy-Preserving of MBs



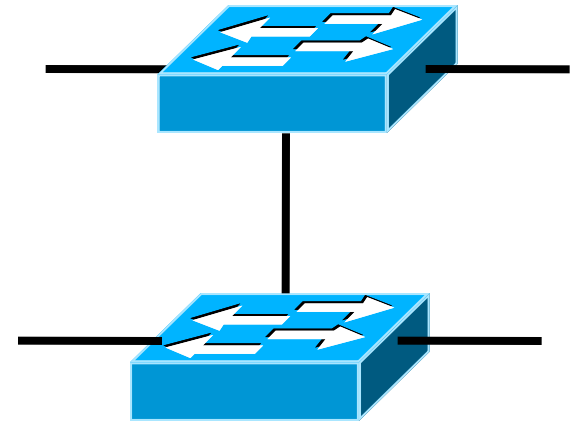
The Internet: A Remarkable Story

- Tremendous success
 - From research experiment to global infrastructure
- Brilliance of under-specifying
 - Network: best-effort packet delivery
 - Hosts: arbitrary applications
- Enables innovation in applications
 - Web, P2P, VoIP, social networks, virtual worlds
- But, change is easy only at the edge... ☹️



Inside the 'Net: A Different Story...

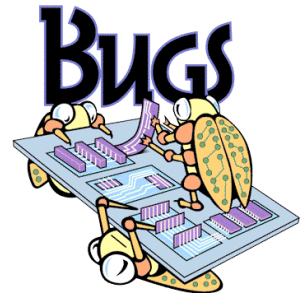
- Closed equipment
 - Software bundled with hardware
 - Vendor-specific interfaces
- Over specified
 - Slow protocol standardization
- Few people can innovate
 - Equipment vendors write the code
 - Long delays to introduce new features



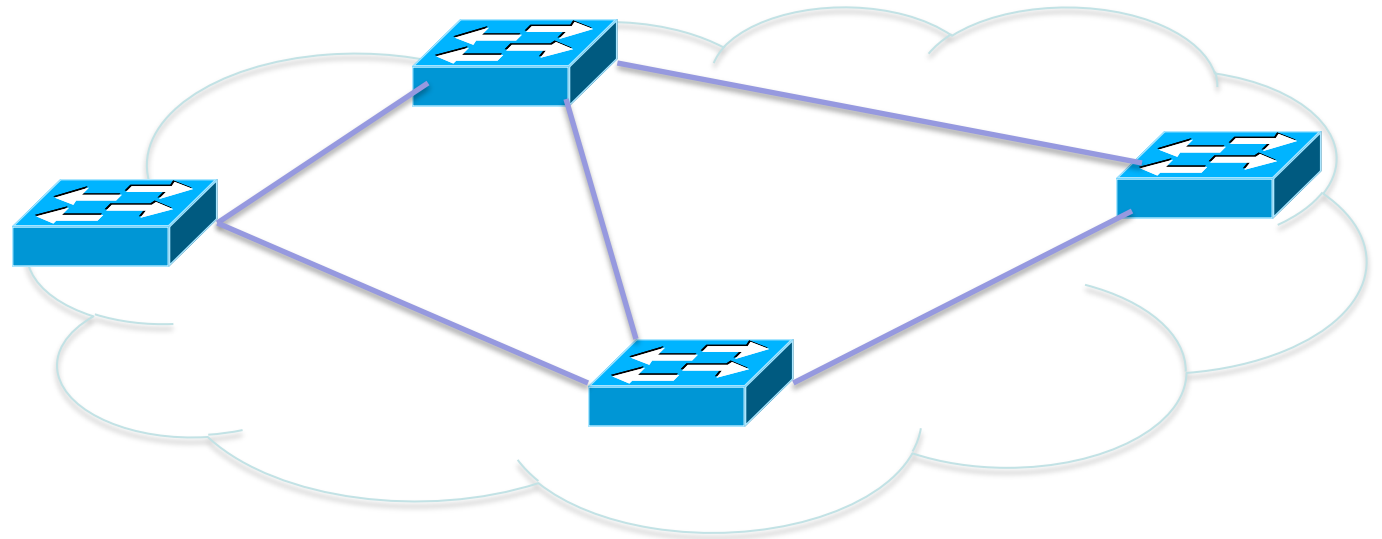
Impacts performance, security, reliability, cost...

Networks are Hard to Manage

- Operating a network is expensive
 - More than half the cost of a network
 - Yet, operator error causes most outages
- Buggy software in the equipment
 - Routers with 20+ million lines of code
 - Cascading failures, vulnerabilities, etc.
- The network is “in the way”
 - Especially a problem in data centers
 - ... and home networks



Traditional Computer Networks

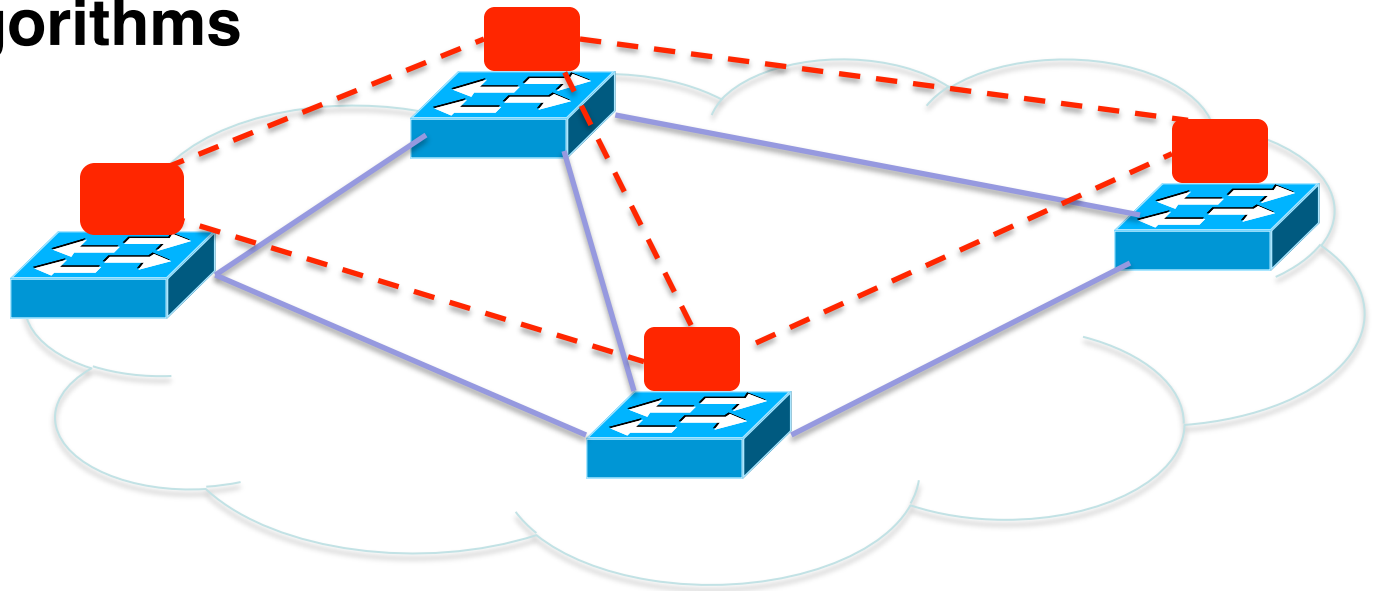


Data plane:
Packet
streaming

Forward, filter, buffer, mark,
rate-limit, and measure packets

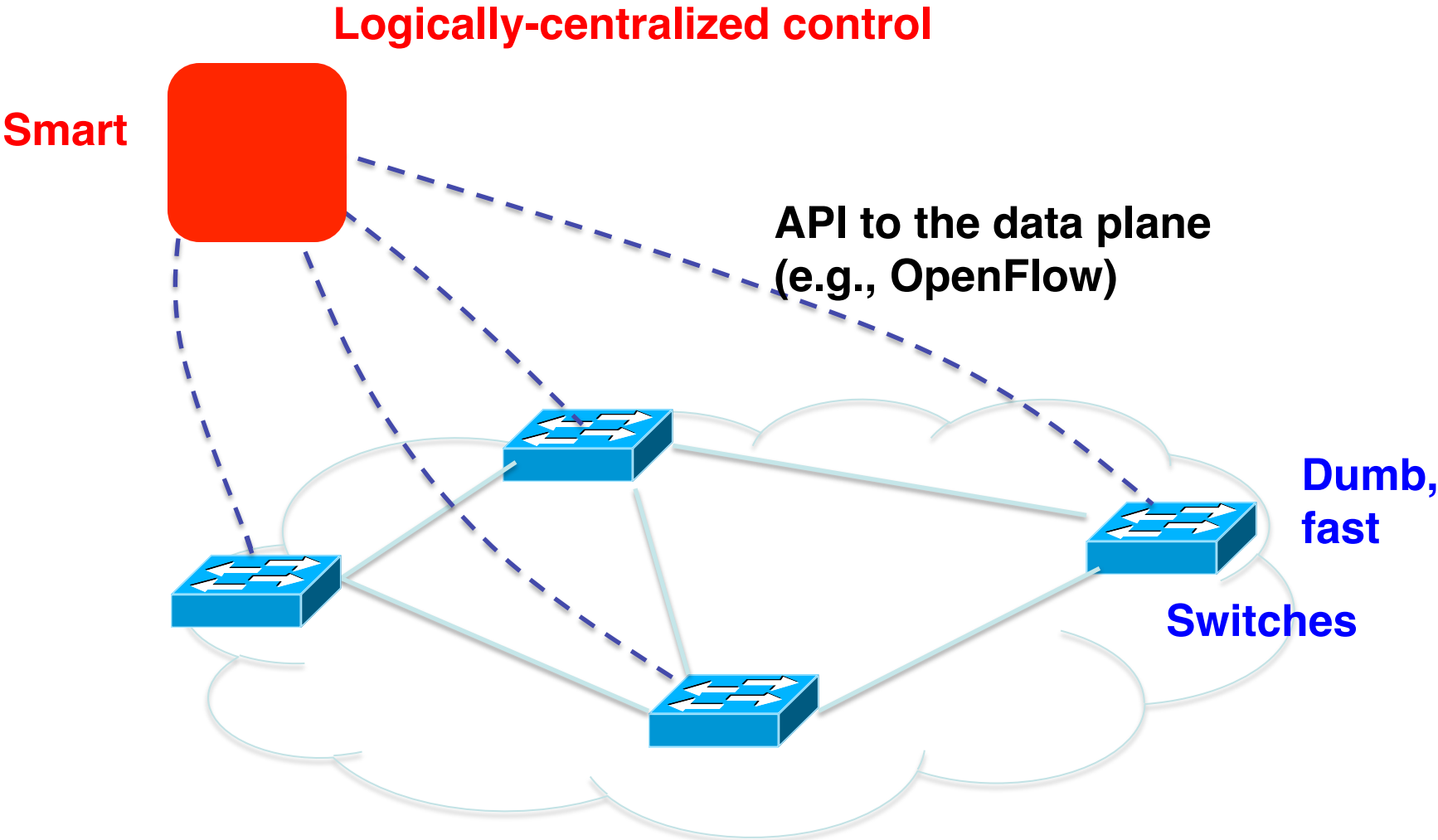
Traditional Computer Networks

Control plane:
Distributed algorithms

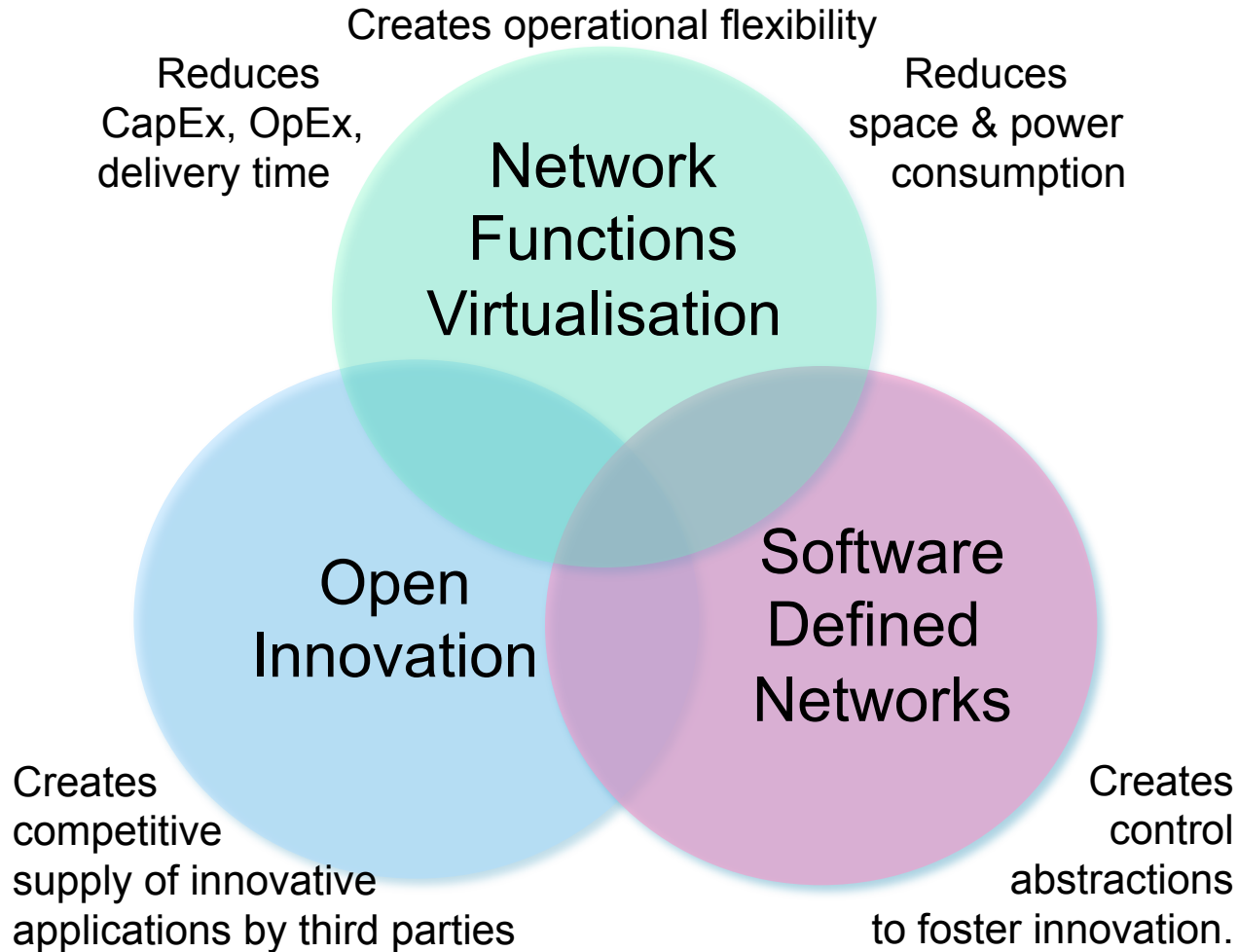


Track topology changes, compute routes, install forwarding rules

Software Defined Networking (SDN)

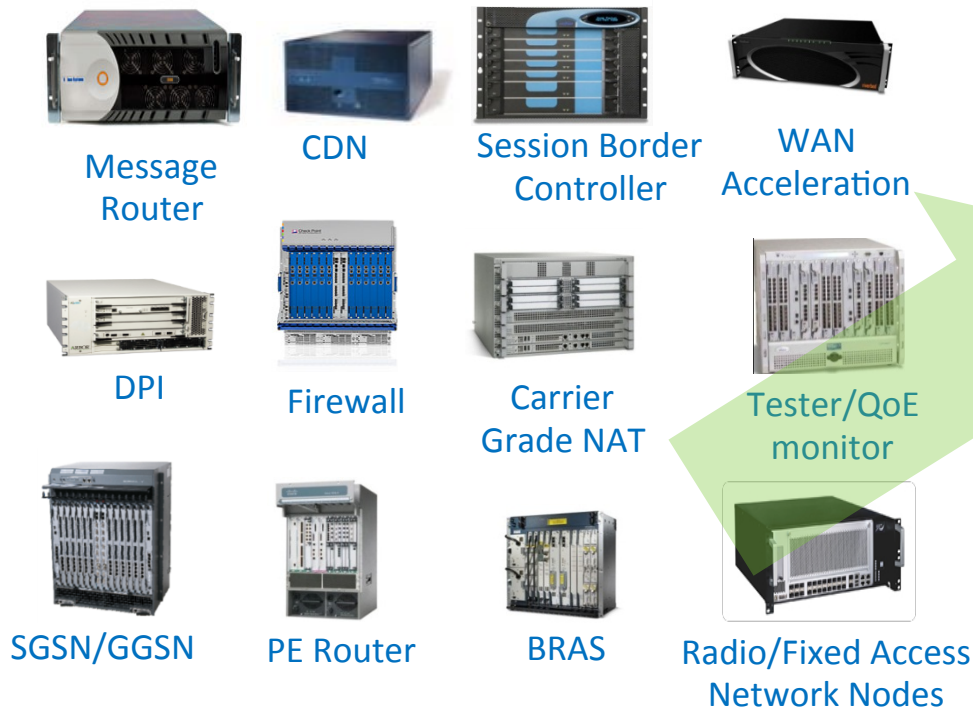


3 Complementary but Independent Networking Developments



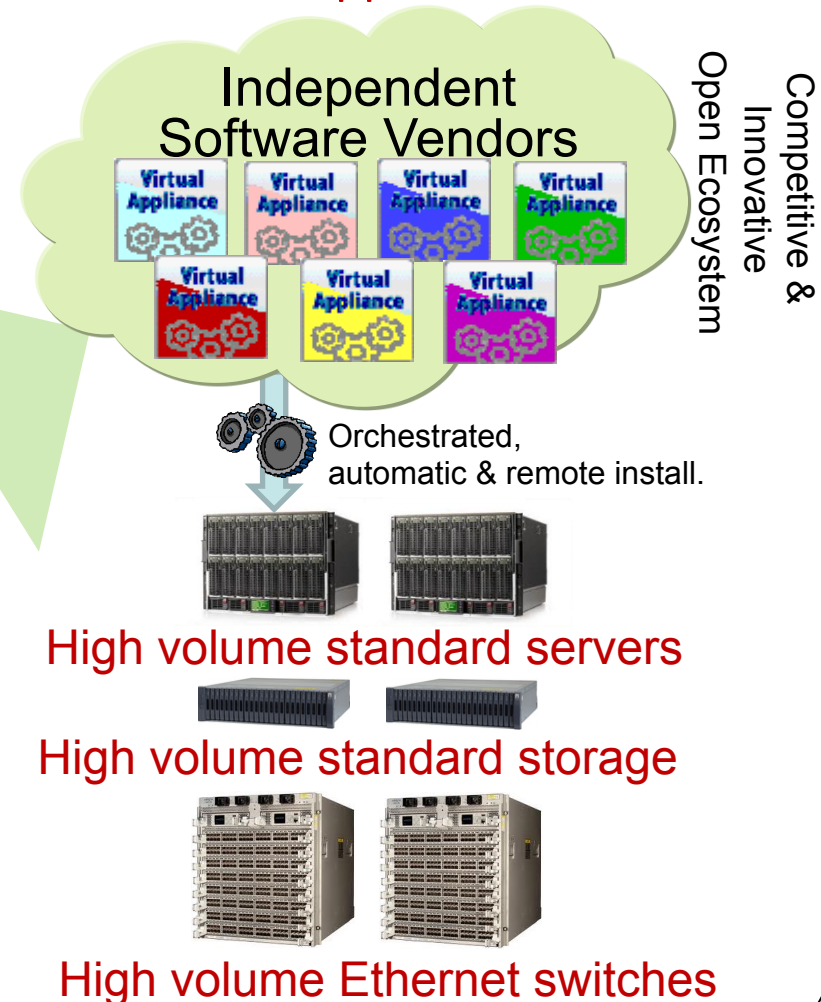
Network Functions Virtualisation: Vision

Classical Network Appliance Approach

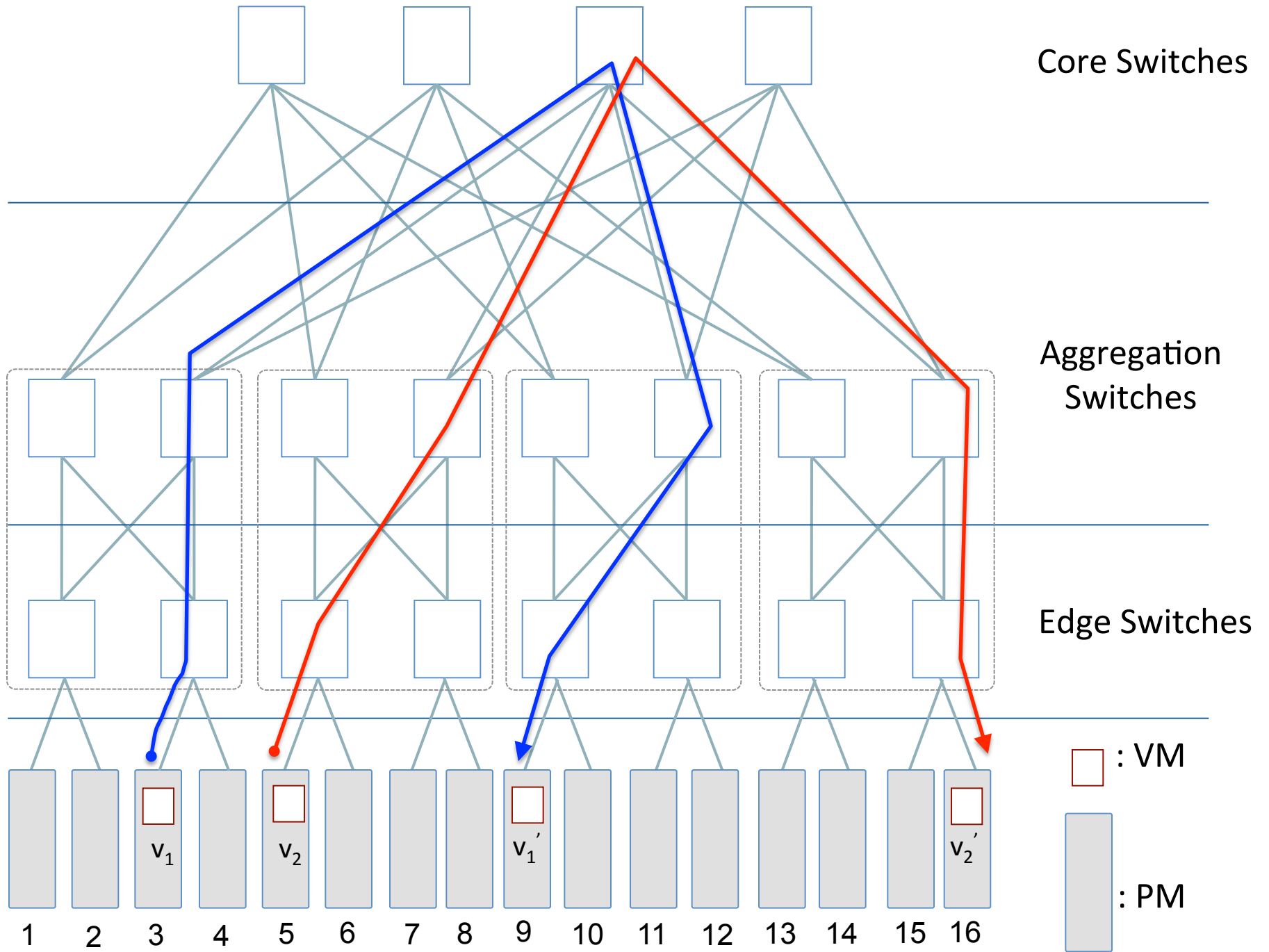


- Fragmented, purpose-built hardware.
- Physical install per appliance per site.
- Hardware development large barrier to entry for new vendors, constraining innovation & competition.

Network Functions Virtualisation Approach



Policy-Preserving MB Placement Problem in Data Centers



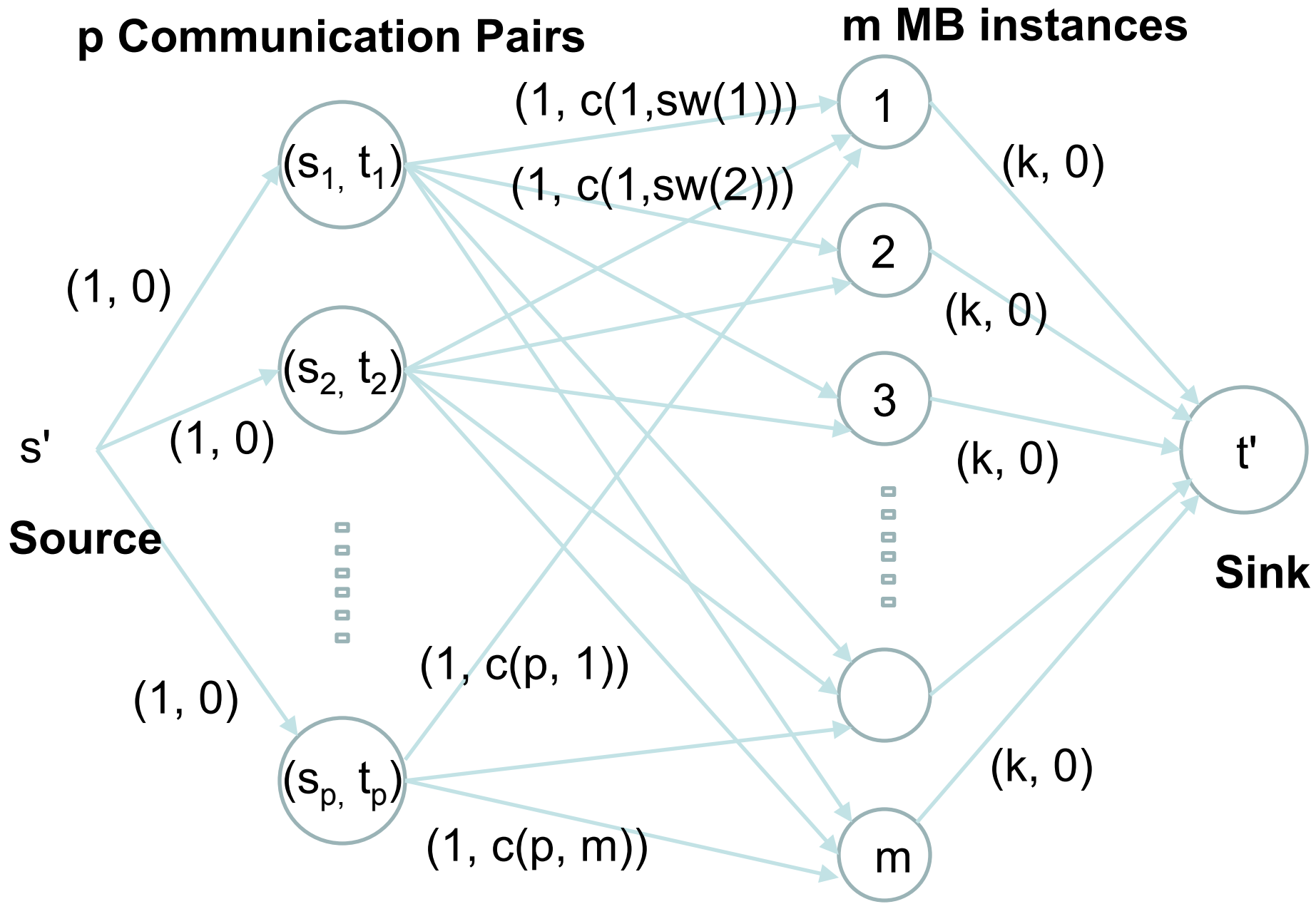
MB Placement Problems

- Many communication pairs in the network
- **Single MB Type**
 - One MB type, say firewall, but multiple instances
- **Multiple MBs Type**
 - each has one instance
 - Ordered Service Chaining
 - Unordered Server Chaining
- **Goal:** Minimize total communication cost
- **Constraint:** Capacity of MB (each can only process limited number of pairs)

Single MB Case

- Given a data center graph $G(V,E)$
- There are m instances of a MB, placed at different node in V
- A set of p communicating node pairs P , each pair (s,t) in P needs to traverse to an instance of a MB
- Each middlebox can only be traversed by at most k pairs
- When $p = (s,t)$ traverses an MB instance m , its cost $c(p,m) = d(s,sw(m)) + d(sw(m),t)$
- Goal: assign all the pairs in P , each traverses one MB instance, s.t. the total cost is minimized, subject to that each MB instance takes at most k pairs.

Solution – minimum cost flow



Ordered Multiple MBs Case

- Given a data center graph $G(V, E)$
- There are m MBs $M = \{mb_1, mb_2, \dots, mb_m\}$ to be placed inside the data center
- A set of p communicating node pairs P , each pair (s, t) in P needs to traverse mb_1, mb_2, \dots, mb_m in that order
- The cost for $p = (s, t)$ is $c(p) = d(s, mb_1) + d(mb_1, mb_2) + \dots + d(mb_{m-1}, mb_m) + d(mb_m, t)$
- **Goal: where to place the m MBs, s.t. the total cost of all p pairs is minimized**

Ordered Multiple MBs Case: Solution

- NP-hard
- **Random:** randomly place the m MBs inside the data center
- **Greedy:** takes place in m rounds
 - In round i , it places mb_i at a node that minimizes the total communication cost so far
- **Load Balancing:** each switch can only accommodate limited number of communication pairs

Un-Ordered Multiple MBs Case

- Given a data center graph $G(V, E)$
- There are m MBs $M = \{mb_1, mb_2, \dots, mb_m\}$ to be placed inside the data center
- A set of p communicating node pairs P , each pair (s, t) in P needs to traverse mb_1, mb_2, \dots, mb_m , but not necessarily in that order
- The cost for $p = (s, t)$ is $c(p) = d(s, mb_{i,1}) + d(mb_{i,1}, mb_{i,2}) + \dots + d(mb_{i,m-1}, mb_{i,m}) + d(mb_{i,m}, t)$
- **Goal: where to place the m MBs, s.t. the total cost of all p pairs is minimized**

Un-Ordered Multiple MBs Case: Solution

- Even more complicated than Ordered Multiple MB case

MB Migration Problems

- Many communication pairs in the network
- Move MBs from their initial location to other locations
- **Goal:** Minimize total communication cost
- **Constraint:** Capacity of MB (each can only process limited number of pairs)

MB Replication Problems

- Many communication pairs in the network
- Multiple MB types, each has one instance
- Goal: How to replicate the MBs, in order to minimize total communication cost
- Constraint: Capacity of switch (each can only store limited number of MB instances)

Conclusions

- Deploying middleboxes is hard, but SDN and NFV makes it easier
- Middleboxes management in SDN-enabled data center is a new and exciting research fields
- Many new algorithmic problems that have not been solved
- Need your participation!

Questions?