

2012

E
IExternal path length and
internal path length

$$E = I + 2n$$

I - the total number of
comparisons needed
to build a binary search
tree whose internal
path length is I.

- the total number of
comparisons to run the
binary search sort
via such a tree
(same as for Quicksort)

Avg. number of comparisons
to search successfully
for a key in a B.S.
tree whose internal
path length is I:

$$C = \frac{I+n}{n} = \frac{I}{n} + 1$$

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Avg. number of comparisons
to search unsuccessfully
for a key in a B.S.
tree whose external
path length is E

$$c' = \frac{E}{n+1} = \frac{I+2n}{n+1} \approx \frac{I+2n}{n} =$$
$$= \frac{I}{n} + 2 = c+1$$

Best case

$$I \approx n \lg n - 2n$$

$$E = I + 2n \approx n \lg n$$

$$c = \frac{I}{n} + 1 \approx \lg n - 1$$

$$c' \approx c + 1 = \lg n$$

Aug. case

$$T \approx 1.4 n \lg n - 2.8 n$$

$$E \approx 1.4 n \lg n - 0.8 n$$

$$C \approx 1.4 \lg n - 1.8$$

$$C' \approx 1.4 \lg n - 0.8$$

Worst case

$$T = \frac{n(n-1)}{2}$$

$$E = \frac{n(n+3)}{2}$$

$$C = \frac{n+1}{2}$$

$$C' = \frac{n+3}{2}$$

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