

## Average ipl (n) and epl (n) in any BS tree on n nodes

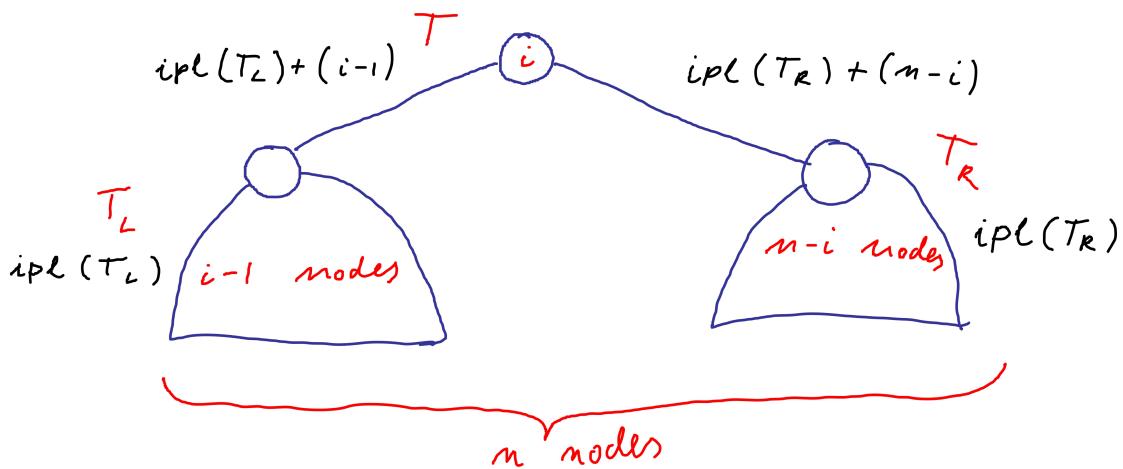
Let T be a BS tree created by n consecutive insertions of numbers

1 through n (not necessarily in that order) into initially empty tree.

There are  $n!$  orders (a.k.a. permutations) of numbers

1 through n. We assume that they are all equally likely.

Let  $i$  be the first number inserted.



$$\begin{aligned} \text{ipl}(T) &= \text{ipl}(T_L) + (i-1) + \text{ipl}(T_R) + (n-i) = \\ &= \text{ipl}(T_L) + \text{ipl}(T_R) + (n-1) \end{aligned}$$

$i$  can be equal to any of the numbers 1 through  $n$ , each with probability  $\frac{1}{n}$ .

Here is the recurrence relation

(refer to your knowledge of Discrete Mathematics on recurrence relations) :

$$\begin{aligned} \text{ipl}_{\text{avg}}[n] &= \frac{1}{n} \sum_{i=1}^n (\text{ipl}_{\text{avg}}[i-1] + \text{ipl}_{\text{avg}}[n-i] + (n-1)) = \\ &= \frac{1}{n} \sum_{i=1}^n (\text{ipl}_{\text{avg}}[i-1] + \text{ipl}_{\text{avg}}[n-i]) + \frac{1}{n} \sum_{i=1}^n (n-1) = \\ &= \frac{1}{n} \sum_{i=1}^n (\text{ipl}_{\text{avg}}[i-1] + \text{ipl}_{\text{avg}}[n-i]) + \frac{1}{n} n (n-1) = \\ &= \frac{1}{n} \sum_{i=1}^n (\text{ipl}_{\text{avg}}[i-1] + \text{ipl}_{\text{avg}}[n-i]) + (n-1) \end{aligned}$$

So,

$$ipl_{avg}[n_] := n - 1 + \frac{1}{n} \sum_{i=1}^n (ipl_{avg}[i - 1] + ipl_{avg}[n - i])$$

$$ipl_{avg}[0] := 0$$

$$ipl_{avg}[1] := 0$$

$$ipl_{avg}[2]$$

$$\frac{1}{1}$$

$$ipl_{avg}[3]$$

$$\frac{8}{3}$$

$$ipl_{avg}[4]$$

$$\frac{29}{6}$$

It may be verified by direct calculation that the function

$$ipl_{avg}[n] = 2(n+1) \left( \sum_{i=1}^n \frac{1}{i} \right) - 4n$$

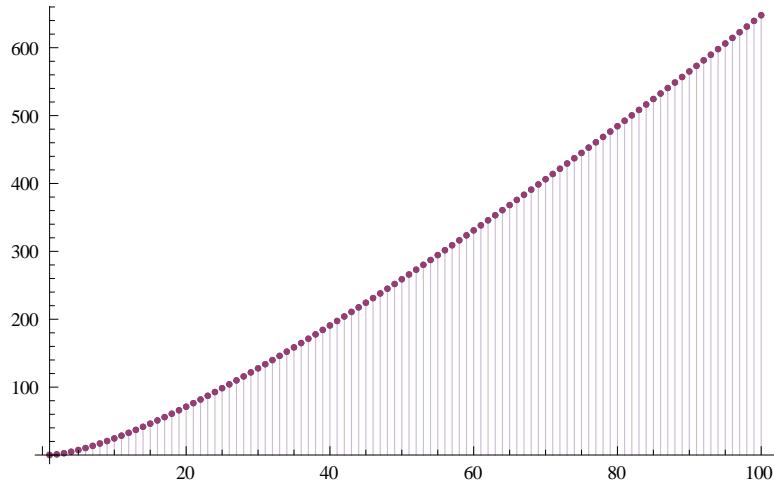
satisfies the above recurrence relation for every  $n \geq 1$ .

To that end it suffices to show that for every  $n \geq 1$ :

$$2(n+1) \left( \sum_{j=1}^n \frac{1}{j} \right) - 4n = \\ n - 1 + \frac{1}{n} \sum_{i=0}^{n-1} \left( 2(i+1) \left( \sum_{j=1}^i \frac{1}{j} \right) - 4i + 2(n-i) \left( \sum_{j=1}^{n-1-i} \frac{1}{j} \right) - 4(n-1-i) \right)$$

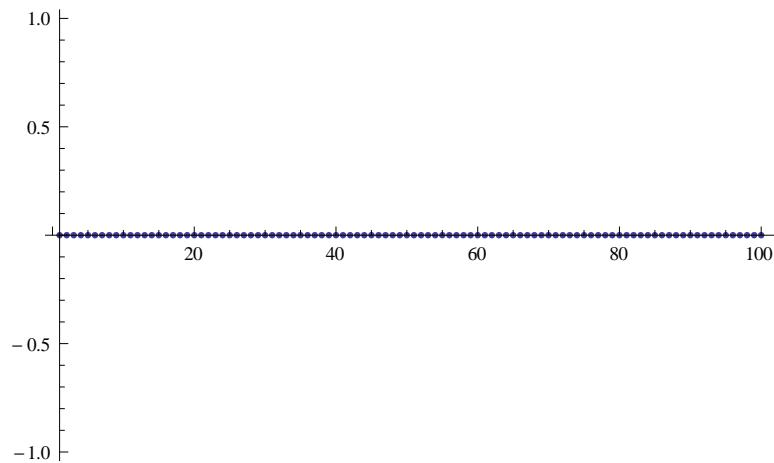
Here is an experimental verification by graphing

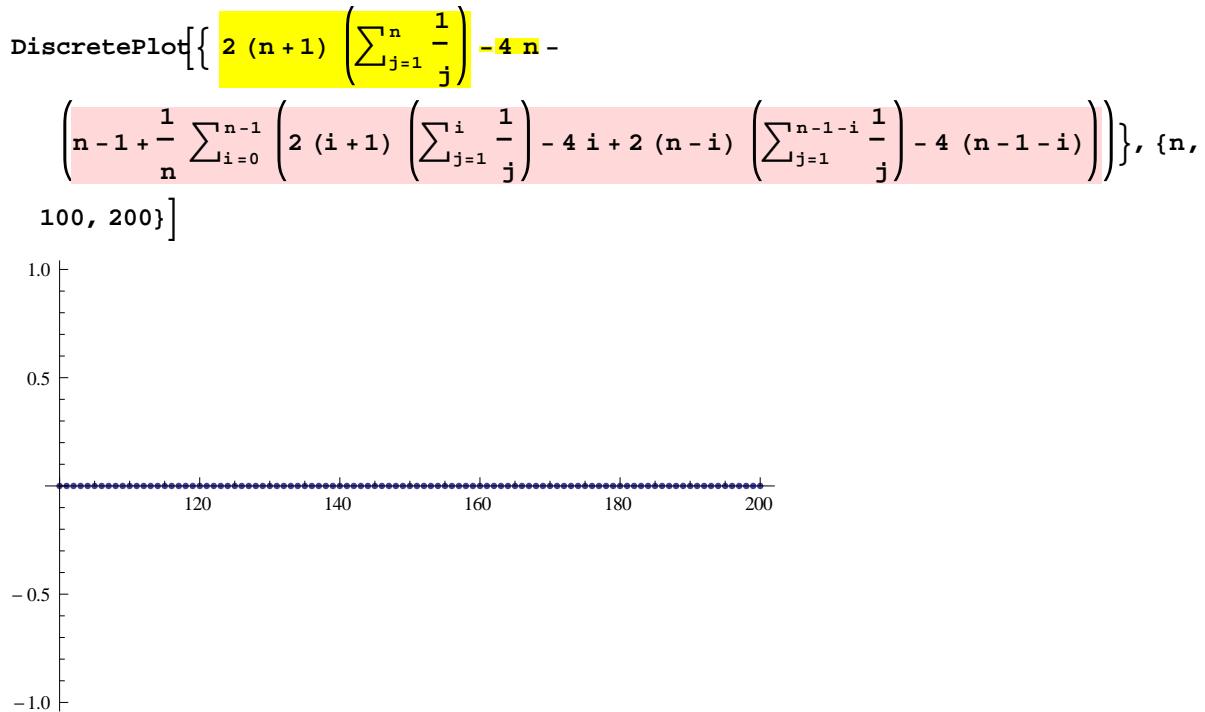
$$\text{DiscretePlot}\left[\left\{ 2(n+1) \left(\sum_{j=1}^n \frac{1}{j}\right) - 4n, n-1 + \frac{1}{n} \sum_{i=0}^{n-1} \left( 2(i+1) \left(\sum_{j=1}^i \frac{1}{j}\right) - 4i + 2(n-i) \left(\sum_{j=1}^{n-1-i} \frac{1}{j}\right) - 4(n-1-i) \right) \right\}, \{n, 1, 100\} \right]$$



The difference between the sides of the equation is 0

$$\text{DiscretePlot}\left[\left\{ 2(n+1) \left(\sum_{j=1}^n \frac{1}{j}\right) - 4n - \left( n-1 + \frac{1}{n} \sum_{i=0}^{n-1} \left( 2(i+1) \left(\sum_{j=1}^i \frac{1}{j}\right) - 4i + 2(n-i) \left(\sum_{j=1}^{n-1-i} \frac{1}{j}\right) - 4(n-1-i) \right) \right) \right\}, \{n, 1, 100\} \right]$$





Substituting the Euler's approximation

$\text{Log}[n] + .577$

of

$$\sum_{i=1}^n \frac{1}{i}$$

in the above formula (see file Summationa.nb) we obtain:

```
ipl_avg[n] ≈ 2(n+1)(Log[n] + .577) - 4n =
Simplify[2(n+1)(Log[n] + .577) - 4n]
1.154 - 2.846 n + 2(1+n) Log[n]
1.154` - 2.846` n + 2 (1+n)  $\frac{\text{Log}[n]}{\text{Log}[2]}$ 
2 N[Log[2]]
1.38629
1.154` - 2.846` n + 1.386 (1+n) Log2[n]
ipl_avg[n] ≈ 1.386 (n+1) Log2[n] - 2.846 n + 1.154
```

Note.

$1.386 (n+1) \text{Log2}[n] - 2.846 n + 1.154$   
is the number of comparison of keys that Quicksort will perform  
on average while sorting an  $n$ -element array with no duplicate keys.

$$epl_{avg}[n] = ipl_{avg}[n] + 2n \approx 1.386(n+1)\log_2[n] - 0.846n + 1.154$$

$$epl_{avg}[n] \approx 1.386(n+1)\log_2[n] - 0.846n + 1.154$$

Average number of comparisons for successful search in an average BS tree :

$$c_n = \frac{ipl_{avg}[n]}{n} + 1 \approx \frac{1}{n} (1.386(n+1)\log_2[n] - 2.846n + 1.154) + 1 \approx \\ 1.386\log_2[n] - 2.846 + 1 = 1.386\log_2[n] - 1.846$$

$$c_n \approx 1.386\log_2[n] - 1.846$$

Average number of comparisons for unsuccessful search in an average BS tree :

$$c'_n = \frac{epl_{avg}[n]}{n+1} \approx (1.386(n+1)\log_2[n] - 0.846n + 1.154) / (n+1) \approx \\ 1.386\log_2[n] - 0.846$$

$$c'_n \approx 1.386\log_2[n] - 0.846$$

NB\*