

Minimum $ipl(n)$ and $ep1(n)$ in any BS tree on n nodes

A binary tree T on n nodes has the minimum internal path length among all binary trees on n nodes,

that is,

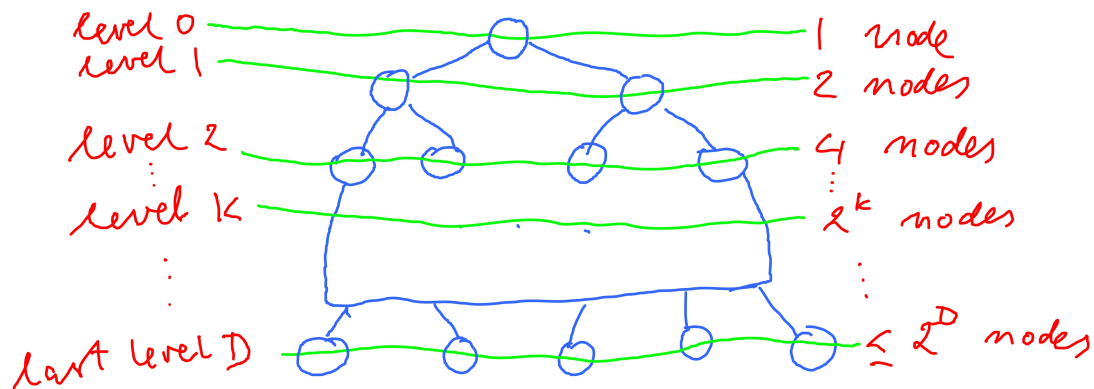
$$ipl(T) = ipl_{\min}(n)$$

iff

T is empty (so $n = 0$)

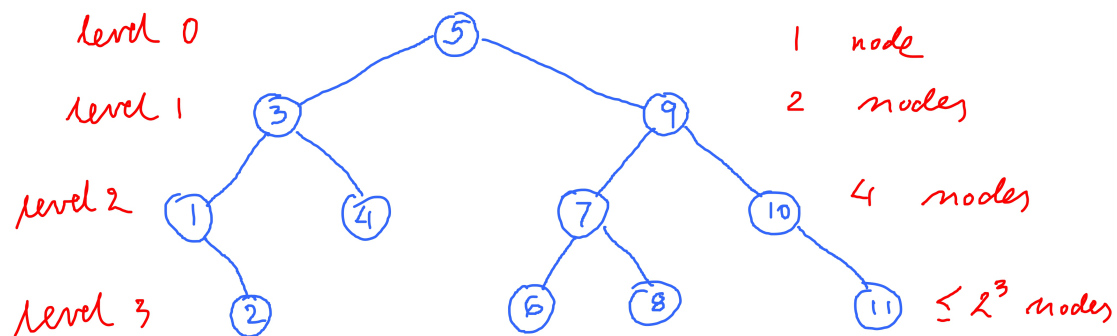
or

all levels of T , except perhaps for the last level of T , have the maximum number of nodes.



Recall that the maximum number of nodes in level k of a binary tree is 2^k .

Here is example of such a BS tree on nodes 1 through 11 :



Hence, the minimum internal path length $ipl_{\min}(n)$ is equal to the internal path length $ipl(H)$ in a heap on n nodes.

Thus,

$$\begin{aligned} \text{ipl}_{\min}(n) &= \sum_{i=1}^n \lfloor \text{Log2}[i] \rfloor = \\ &= (n+1) \lfloor \text{Log2}[n+1] \rfloor - 2^{\lfloor \text{Log2}[n+1] \rfloor + 1} + 2 = \\ &= (n+1) (\text{Log2}[n+1] + \epsilon) - 2n, \end{aligned}$$

where ϵ is a function of $n+1$ oscillating between:

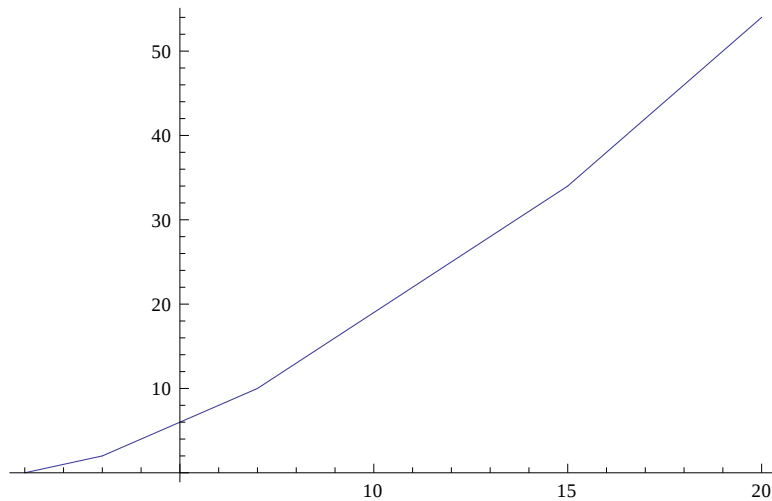
$$0 \leq \epsilon < 0.086071332056$$

Note.

$$(n+1) (\text{Log2}[n+1] + \epsilon) - 2n$$

is the number of comparisons of keys that Quicksort will perform in the best case while sorting an n -element array with no duplicate keys.

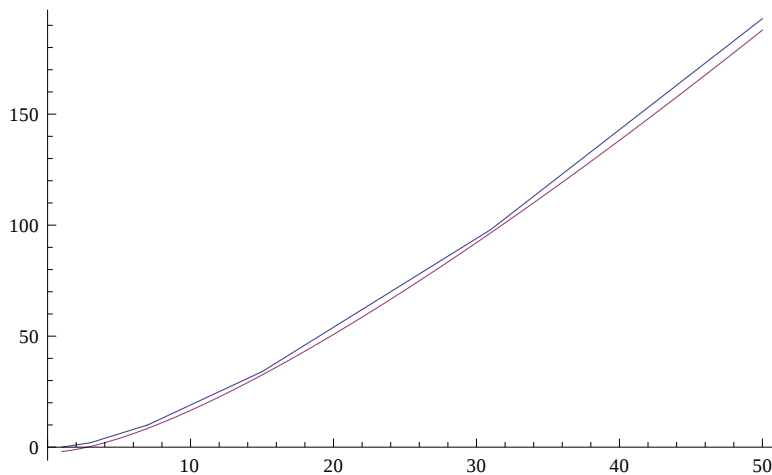
Plot[$(n+1) \lfloor \text{Log2}[n+1] \rfloor - 2^{\lfloor \text{Log2}[n+1] \rfloor + 1} + 2, \{n, 1, 20\}$]



So,

$$\text{ipl}_{\min}(n) \approx (n+1) \text{Log2}[n] - 2n.$$

Plot[$\{(n+1) \lfloor \text{Log2}[n+1] \rfloor - 2^{\lfloor \text{Log2}[n+1] \rfloor + 1} + 2, (n+1) \text{Log2}[n] - 2n\}, \{n, 1, 50\}$]



Therefore,

$$epl_{\min}(n) = ipl_{\min}(n) + 2n \approx (n+1) \log_2 n$$

or

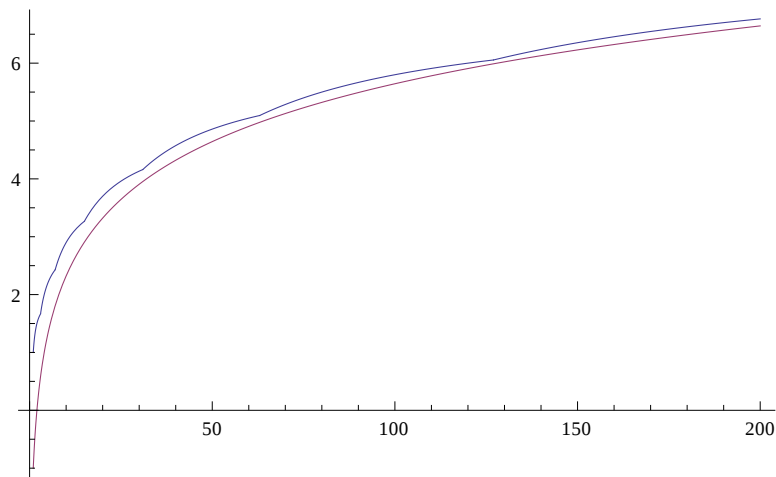
$$epl_{\min}(n) \approx (n+1) \log_2 n$$

Average number of comparisons for **successful** search in an average BS tree :

$$c_n = \frac{ipl_{\text{avg}}[n]}{n} + 1 \approx \frac{1}{n} ((n+1) \log_2 n - 2n) + 1 \approx \log_2 n - 2 + 1 = \log_2 n - 1.$$

$$c_n \approx \log_2 n - 1$$

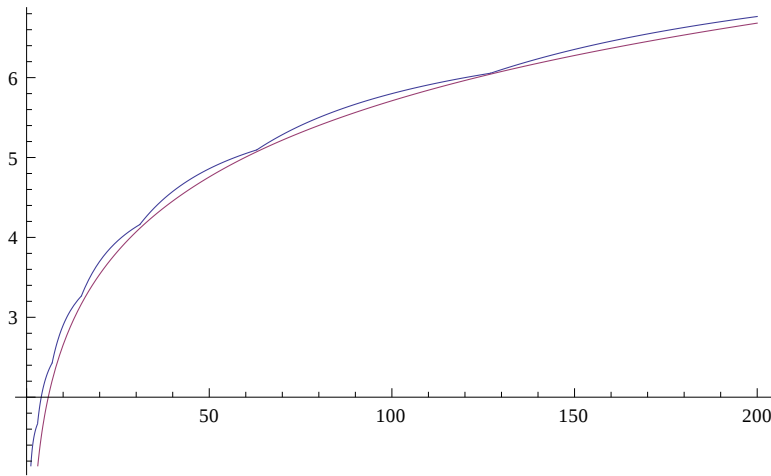
Plot[$\{1 + ((n+1) \lfloor \log_2(n+1) \rfloor - 2^{\lfloor \log_2(n+1) \rfloor + 1} + 2) / n, \log_2 n - 1\}$,
 $\{n, 1, 200\}$, PlotPoints $\rightarrow$ 200]



Here is a bit more accurate approximation:

$$c_n \approx \log_2 n - 1 + \frac{\log_2 n}{n}$$

Plot[$\left\{1 + \frac{(n+1) \lfloor \text{Log2}[n+1] \rfloor - 2^{\lfloor \text{Log2}[n+1] \rfloor + 1} + 2}{n}, \text{Log2}[n] - 1 + \frac{\text{Log2}[n]}{n}\right\},$
 $\{n, 1, 200\}, \text{PlotPoints} \rightarrow 200]$



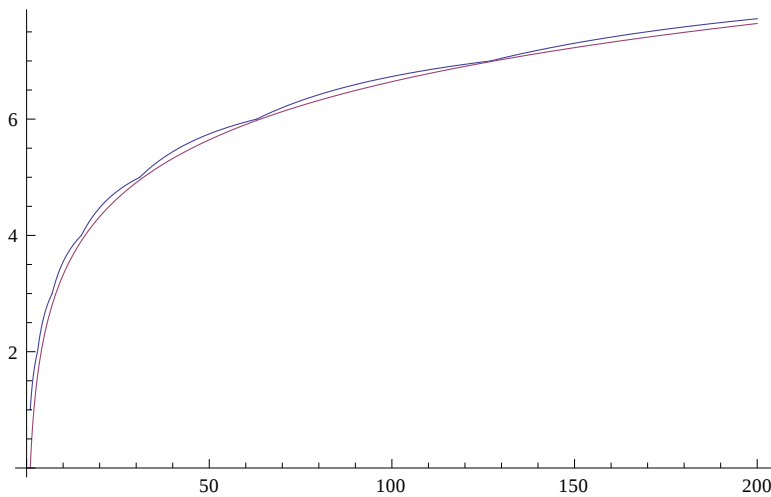
Average number of comparisons for **unsuccessful** search in an average BS tree :

$$c_n' = \frac{\text{epI}_{\text{avg}}[n]}{n+1} \approx \frac{(n+1) \text{Log2}[n]}{n+1} = \text{Log2}[n]$$

$$1.386 \text{Log2}[n] - 0.846$$

$$c_n \approx \text{Log2}[n].$$

Plot[$\left\{\frac{(n+1) \lfloor \text{Log2}[n+1] \rfloor - 2^{\lfloor \text{Log2}[n+1] \rfloor + 1} + 2 + 2n}{(n+1)}, \text{Log2}[n]\right\},$
 $\{n, 1, 200\}, \text{PlotPoints} \rightarrow 200]$



NB*