

# CSC 311

## Lectures on **Data Structures**

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# CSC 311

## Lecture 10 ADT Stack and ADT Queue

**Definitions, Applications, Implementations,  
Analysis**

# Linear Data Structures— Stacks and Queues

LIFO

using stacks to process  
nested structures

FIFO

context-free  
grammar

parsing, evaluation,  
and backtracking

queues, clients, and servers

## 6.2 Some Background on Stacks

### LEARNING OBJECTIVES

1. To learn some of the common terminology.
2. To become informally acquainted with a significant use of stacks in handling recursion, in the theory of push-down automata, and in the theory of parsing, translation, and compiling.

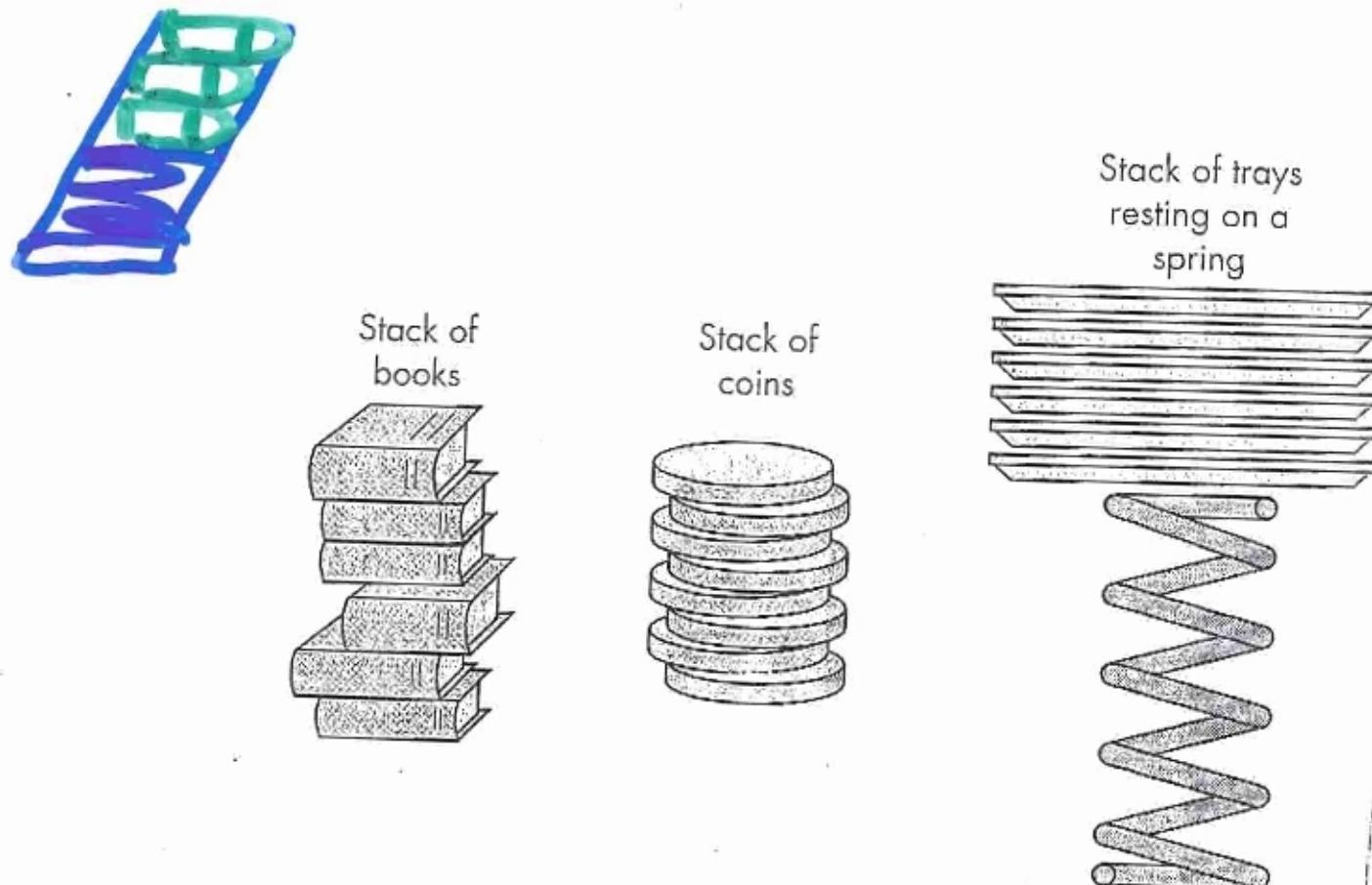


Figure 6.1 Various Kinds of Stacks

# Stacks

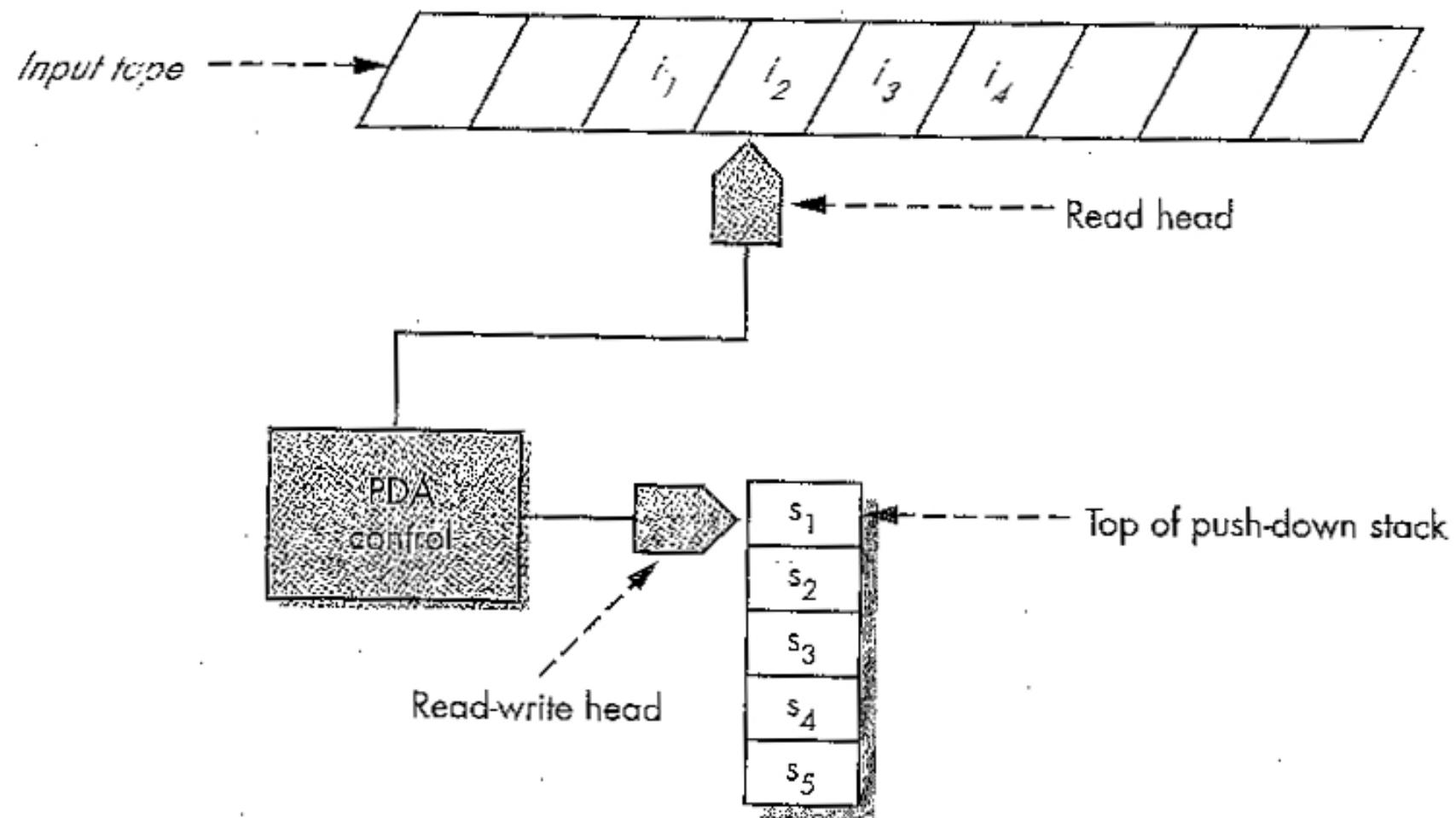


Figure 6.2 Schematic Diagram of a Push-Down Automaton

# Stacks

## Stacks

A Stack,  $S$ , is a sequence of items on which the following operations are defined:

1. Construct an initially empty stack,  $S$ .
2. Determine whether or not the stack,  $S$ , is empty.
3. Push a new item onto the top of the stack,  $S$ .
4. If  $S$  is nonempty, pop an item from the top of stack,  $S$ .
5. If  $S$  is nonempty, peek at the top of stack  $S$  by reading a copy of the top item of  $S$  without removing it.

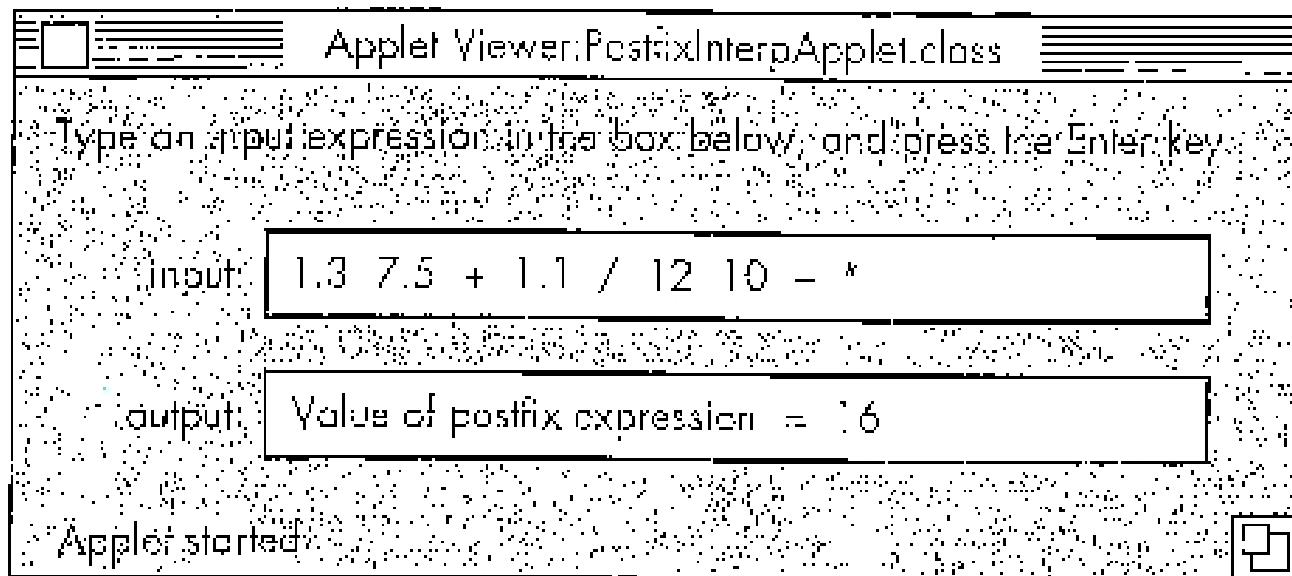
# Stacks

```
1   S = new Stack();           // creates an initially empty Stack S
2
3   S.empty();                // a boolean expression that is true if and
                           // only if Stack S contains no items
4
5   S.push(X);                // pushes an item X onto the top of Stack S
6
7   X = S.pop();              // removes the top item of S and puts it in X
8
9   X = S.peek();             // puts a copy of the top item of S in X
                           // without removing it from S
10
11
```

**Figure 6.3** Informal Method Calls in a Stack ADT Interface

**Table 6.8** Translations of Expressions into Corresponding Postfix Expressions

Prefix Expression	Postfix Expression
$(a + b)$	$a b +$
$(x - y - z)$	$x y - z -$
$(x - y - z) / (u + v)$	$x y - z - u v + /$
$(a^2 + b^2) * (m - n)$	$a^2 b^2 + m n - *$



**Figure 6.9** Evaluating a Postfix Expression

# Queues

## Queues

A Queue,  $Q$ , is a sequence of items on which the following operations are defined:

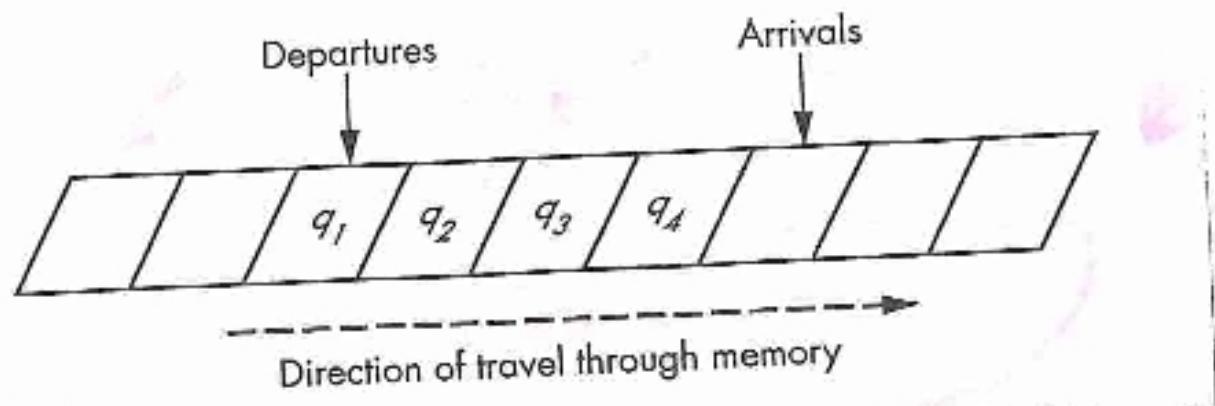
1. Construct an initially empty queue,  $Q$ .
2. Determine whether or not the queue,  $Q$ , is empty.
3. Insert a new item onto the rear of the queue,  $Q$ .
4. Provided  $Q$  is nonempty, remove an item from the front of  $Q$ .



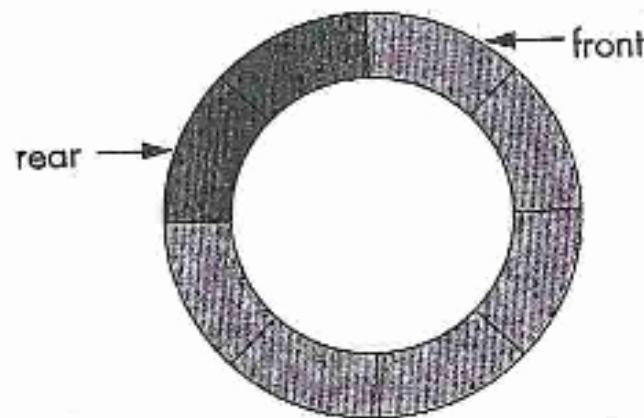
# Queues

```
Q = new Queue( );           // creates an initially empty Queue Q  
10   Q.empty( );           // a boolean expression that is true if and  
    // only if the Queue Q contains no items  
    enqueue  
    Q.insert(X);           // inserts an item X onto the rear of Queue Q  
15   X = Q.remove( );       // removes an item from the front of Q  
    // and puts it in X  
    dequeue
```

Figure 6.4 Informal Method Calls for a Queue ADT Interface



**Figure 6.20** Queue on a Circular Track



$\text{front} = (\text{front} + 1) \% n;$   
 $\text{rear} = (\text{rear} + 1) \% n;$

**Figure 6.21** Incrementing Indexes Using Modular Arithmetic

```

class Queue {
    private int front;           // the array index of the front item of the queue
    private int rear;            // the array index for the next item to insert
    private int count;           // the number of items in the queue
    private int capacity;        // the number of available array positions
    private int capacityIncrement; // the amount to increase the capacity
                                  // during array expansion
    private Object[] itemArray;   // the array that holds queue items
}

// here, we need the no-arg constructor

public Queue() {
    front = 0;
    rear = 0;
    count = 0;
    capacity = 10;
    capacityIncrement = 5;
    itemArray = new Object[capacity];
}

public boolean empty() {
    return (count == 0);
}

/*-----*/
public void insert(Object newItem) {
    // if the itemArray does not have enough capacity,
    // expand the itemArray by the capacity increment
    if (count == capacity) {
        capacity += capacityIncrement;
        Object[] tempArray = new Object[capacity];
        if (front < rear) { // if the items are in itemArray[front:rear-1]
            for (int i = front; i < rear; i++) {
                tempArray[i] = itemArray[i];
            }
        } else { // otherwise, move the items in two separate sections
            for (int i = 0; i < rear; i++) { // section one:
                tempArray[i] = itemArray[i]; // itemArray[0:rear-1]
            }
            for (int i = front; i < count; i++) { // and section two:
                tempArray[i+capacityIncrement] = itemArray[i];
            }
        }
        front += capacityIncrement; // then change front to point to
                                   // its new position
        itemArray = tempArray;
    }
}

```