CSC 501/401

Lectures on
Analysis of Algorithms

by

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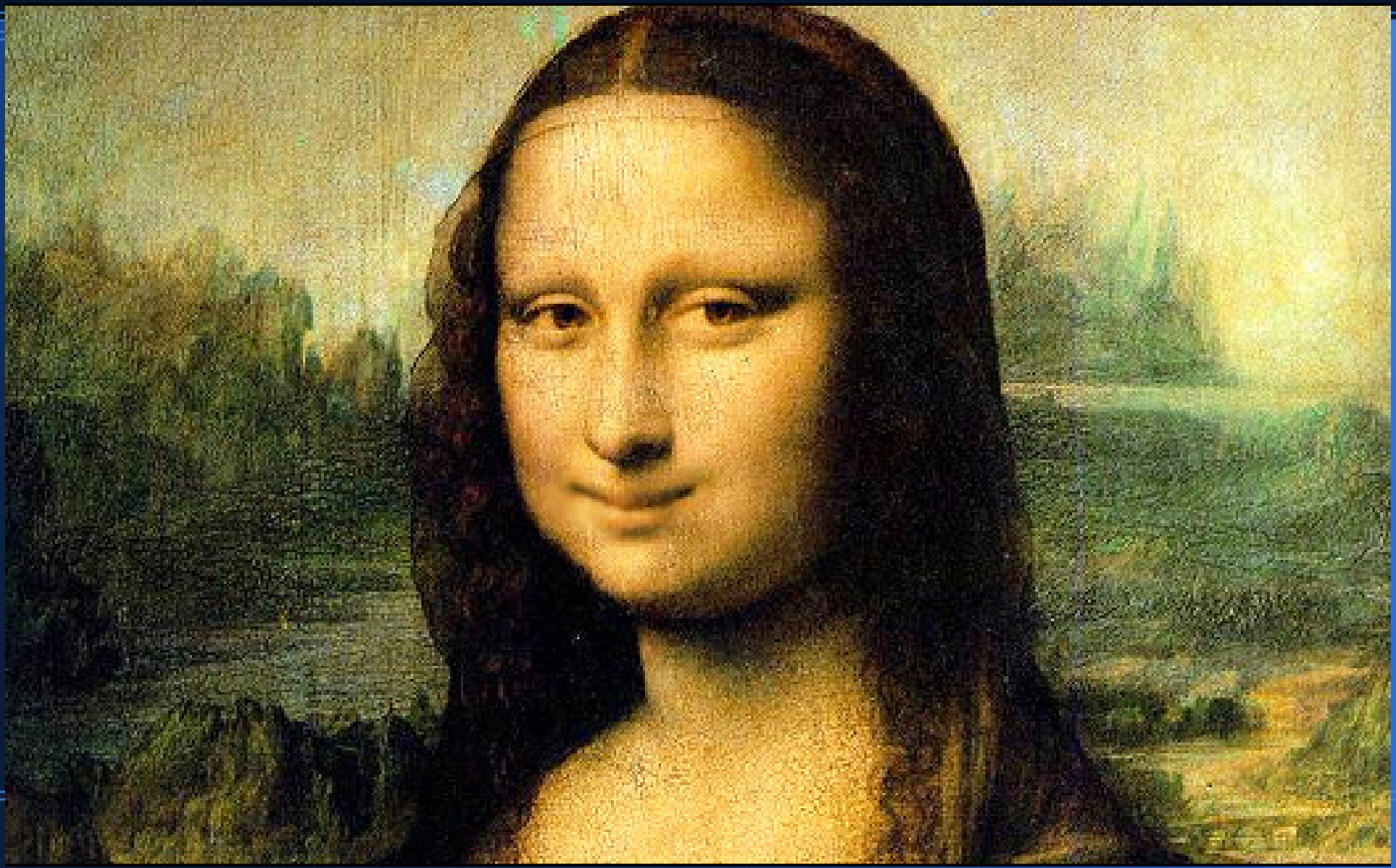
Computer Science
CSUDH

CSC 501/401

Chapter 7
Graphs and Graph Traversals

Depth-First Search
Breadth-First Search

This will not be covered by Test 2



Graphs and Digraphs

Definitions and Representations

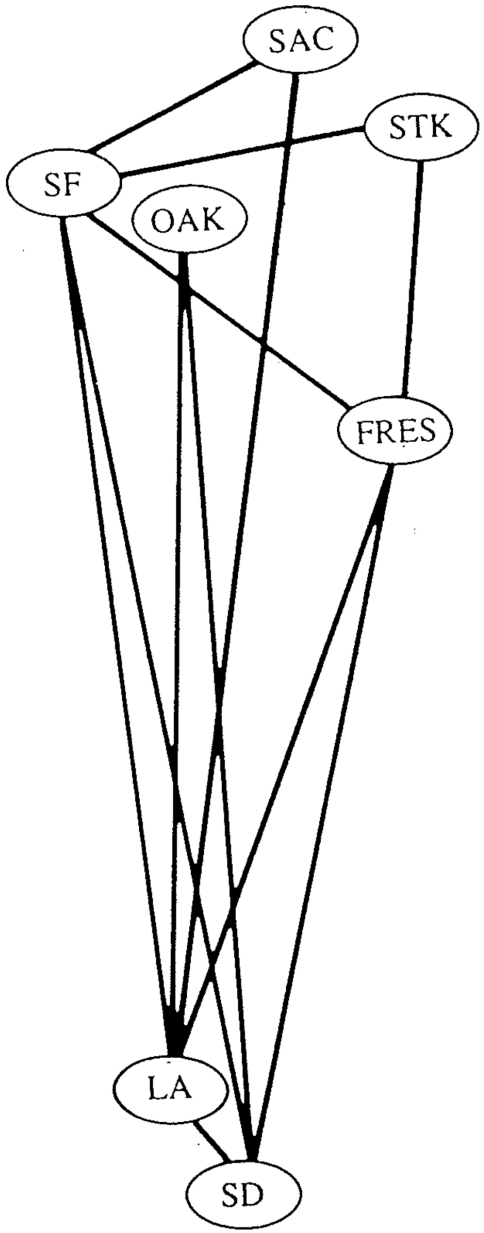
A Minimum Spanning Tree Algorithm

A Shortest-Path Algorithm

Traversing Graphs and Digraphs

Definitions and Examples

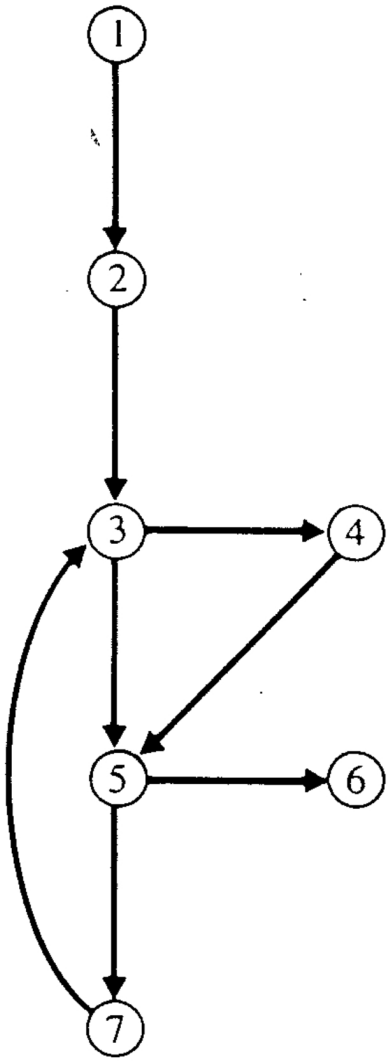
graph, G, is a pair (V, E)



$V = \{SF, OAK, SAC, STK, FRES, LA, SD\}$

$E = \{\{SF, STK\}, \{SF, SAC\}, \{SF, LA\}, \{SF, SD\}, \{SF, FRES\}, \{SD, OAK\}, \{SAC, LA\}, \{LA, OAK\}, \{LA, FRES\}, \{LA, SD\}, \{FRES, STK\}, \{SD, FRES\}\}$.

A *digraph*, G, is a pair (V, E)



$$V = \{1, 2, \ldots, 7\}$$

$$E = \{(1,2),\ (2,3),\ (3,4),\ (3,5),\ (4,5),\ (5,6),\ (5,7),\ (7,3)\}$$

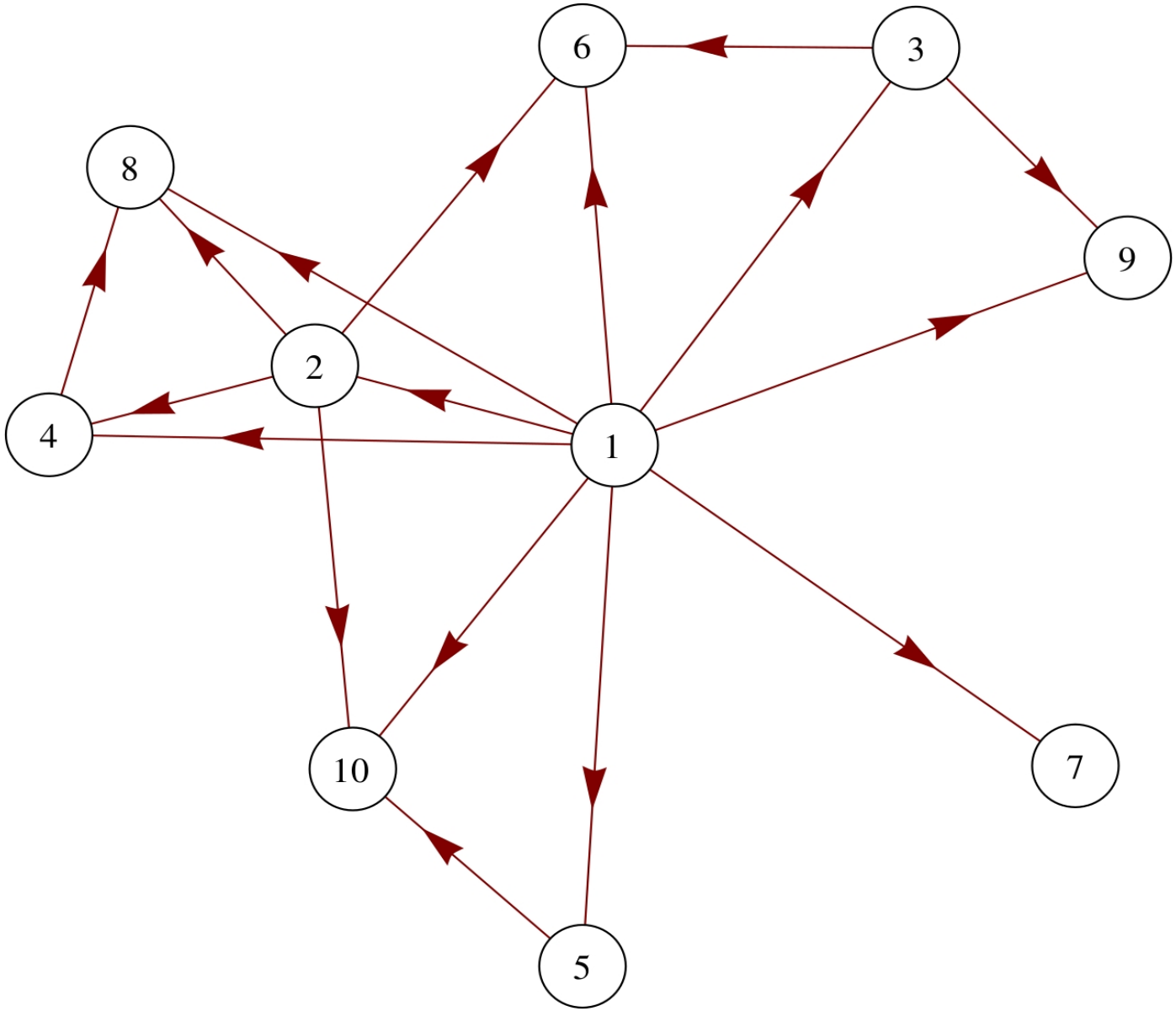
(v, w) is represented in the diagrams as $v \rightarrow w$.

(x,y)

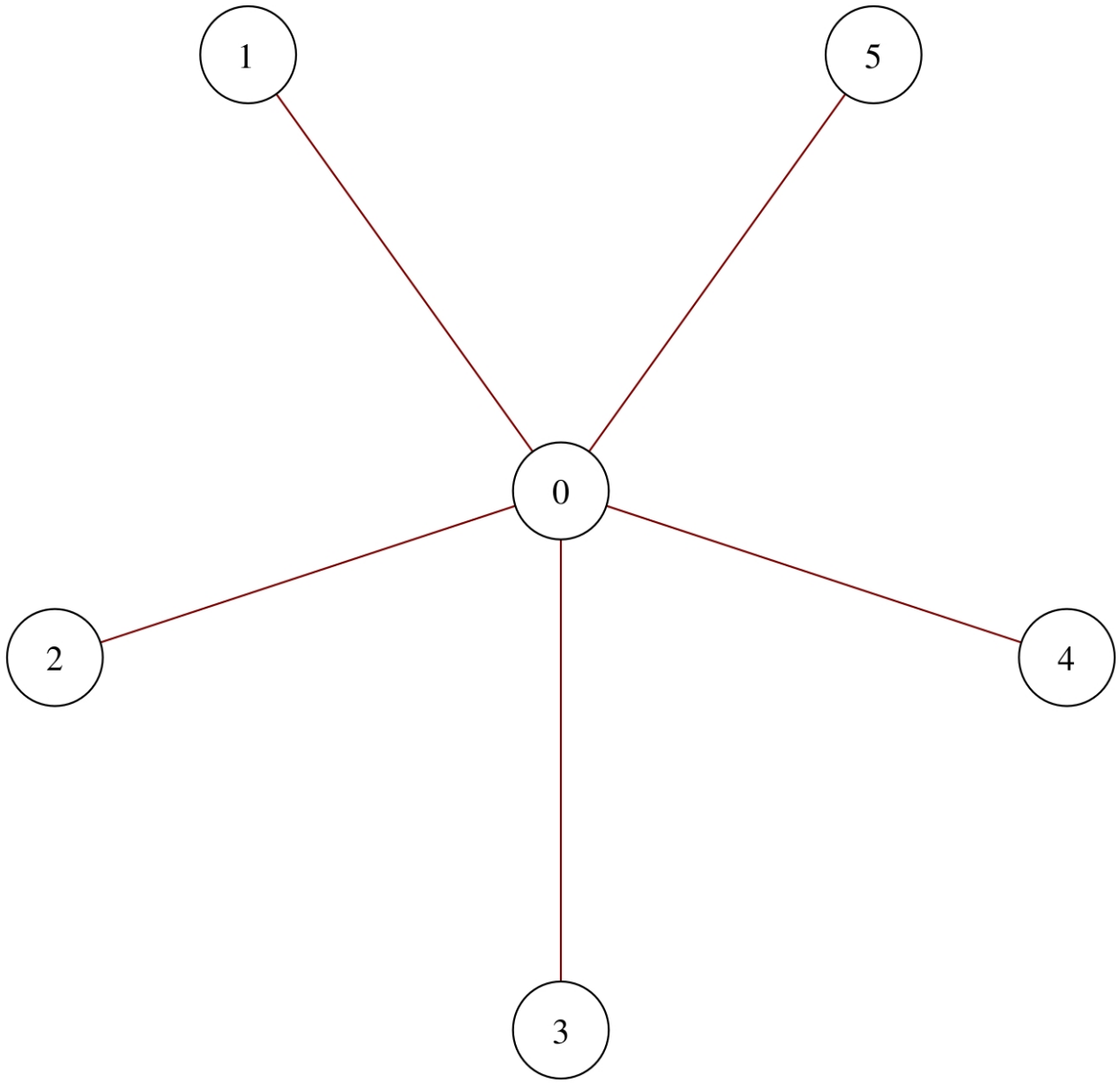
in E

iff

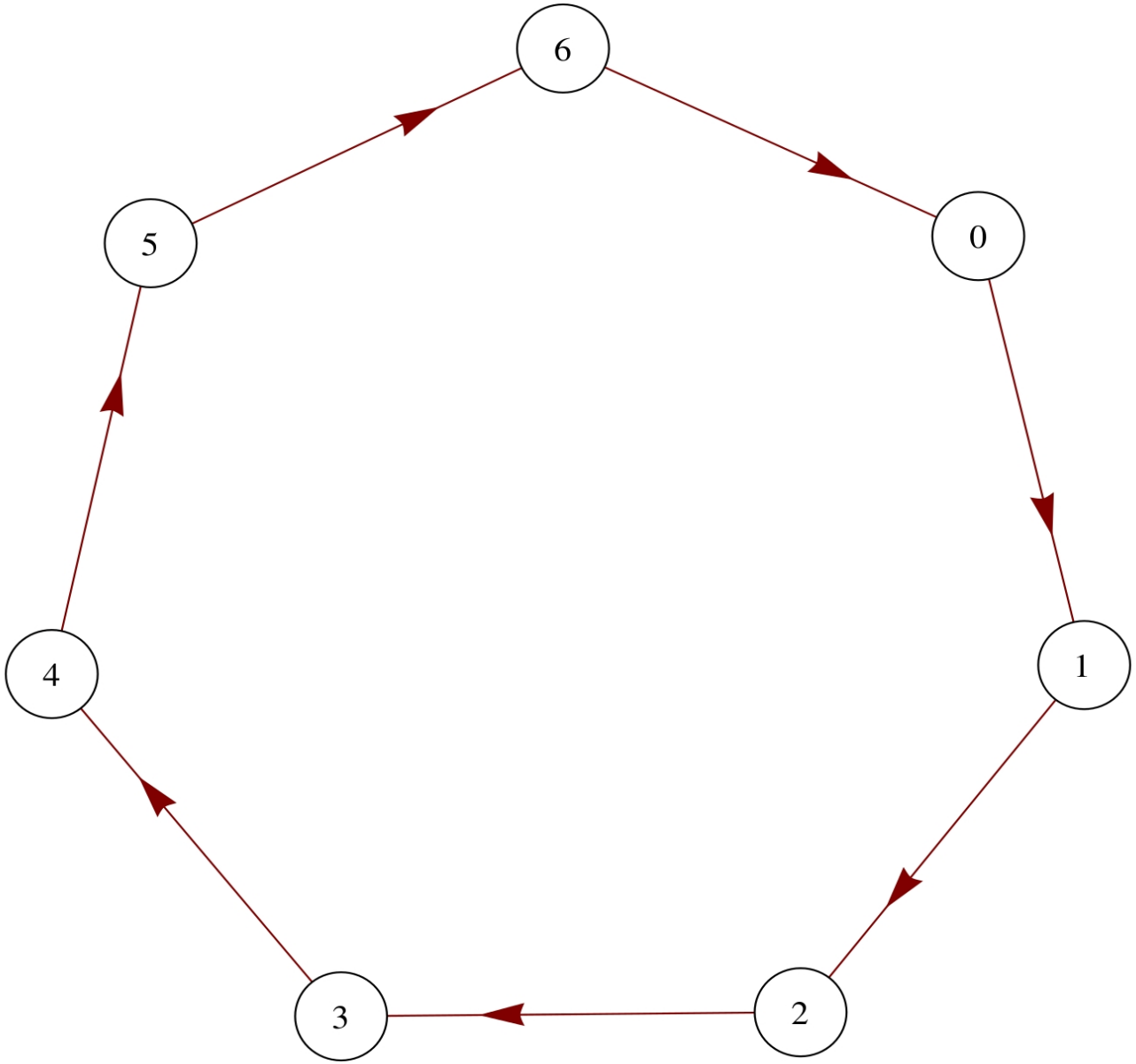
y|x



Star

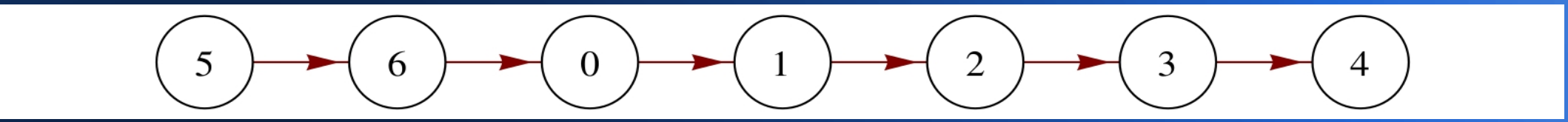


Cycle



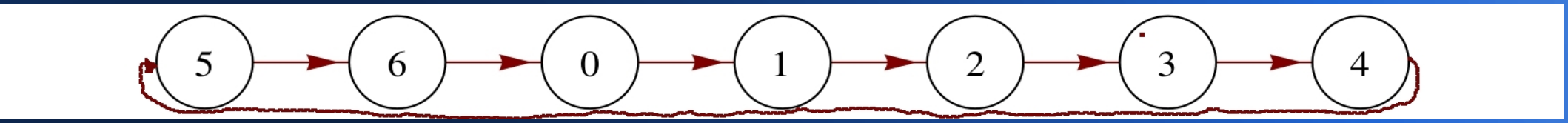
Graphs and Digraphs

Cycle



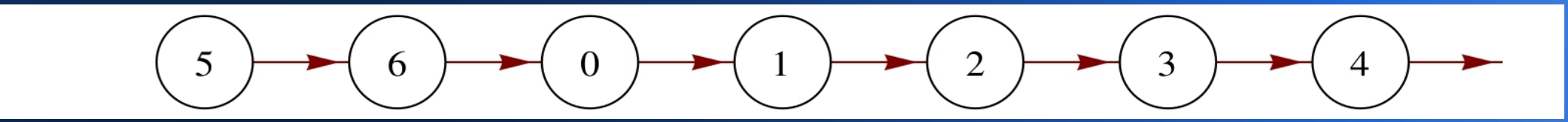
Graphs and Digraphs

Cycle



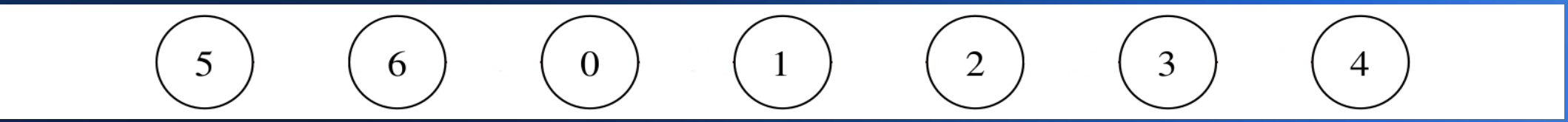
Graphs and Digraphs

Cycle



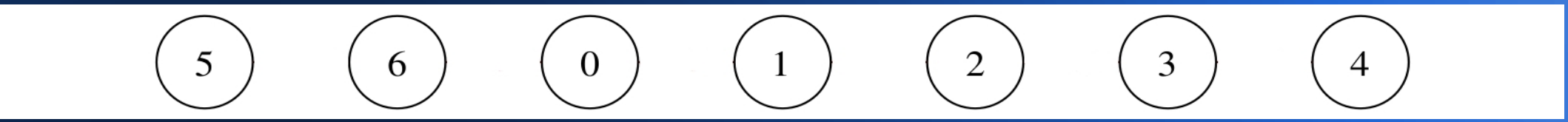
Graphs and Digraphs

Cycle



Graphs and Digraphs

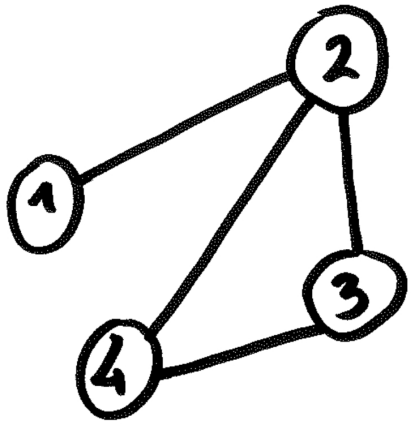
Cycle



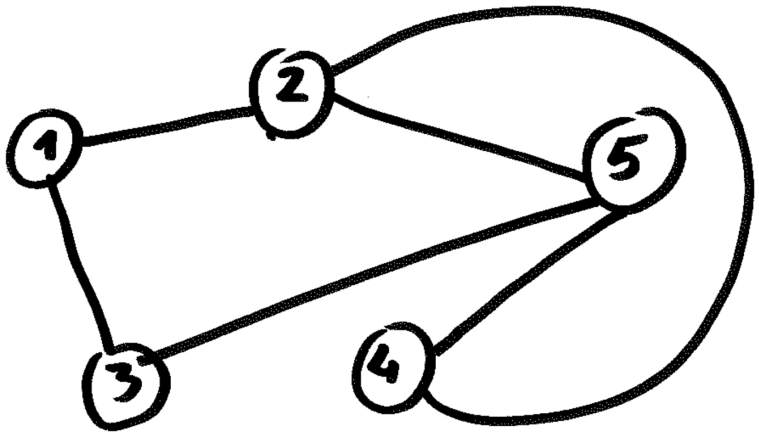
(5, 6, 0, 1, 2 ,3 ,4)

Gra ns

GRAPH



SUBGRAPH



PATH: 5 3 1 2 4

CYCLE: 5 3 1 2 4

Finding Connected Components of a Graph

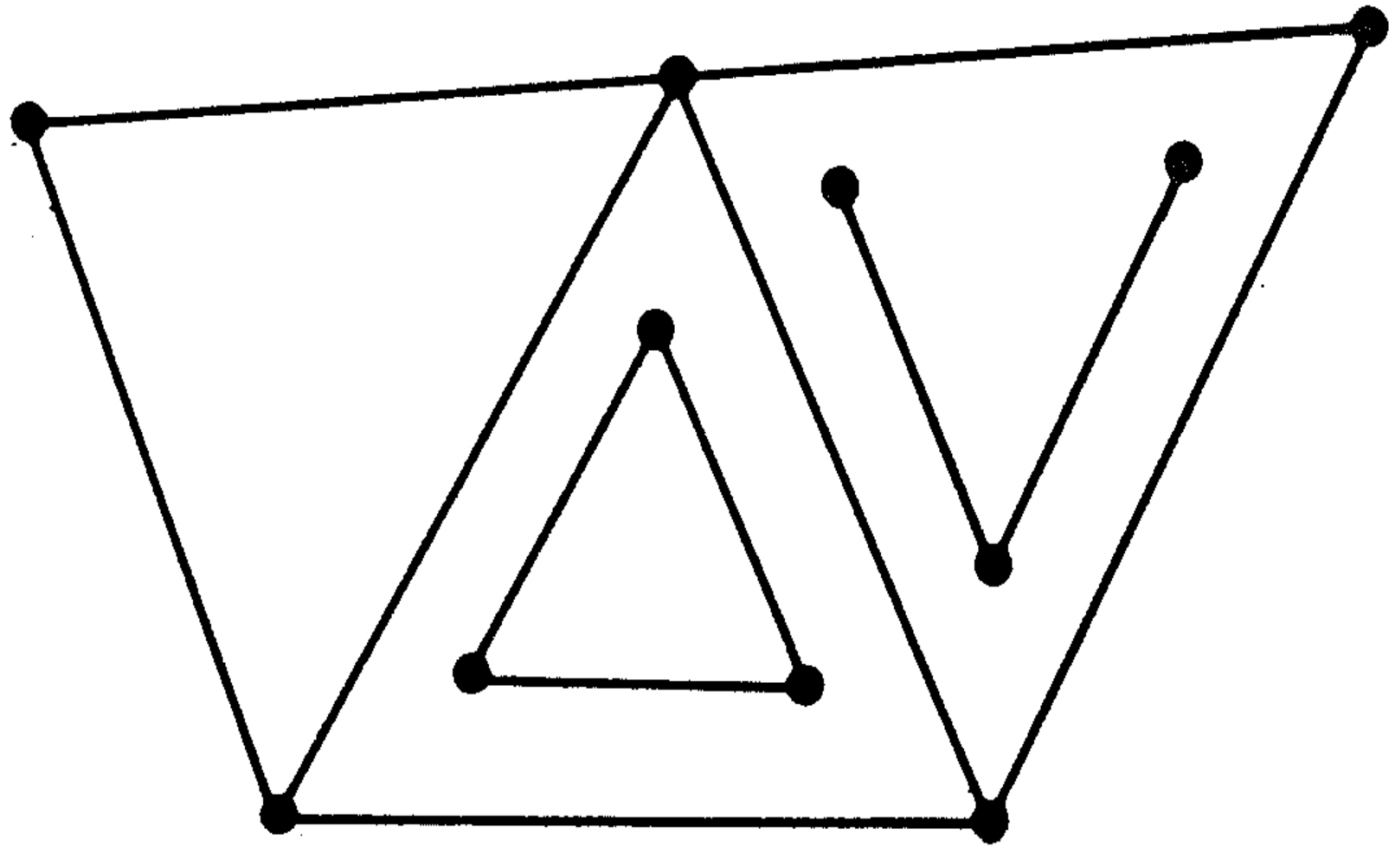


Figure 4.7 A graph with three connected components.

Computer Representation of Graphs and Digraphs

adjacency matrix $A = (a_{ij})$

Let $G = (V, E)$

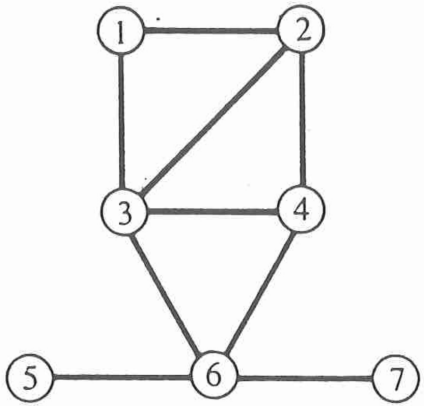
$$a_{ij} = \begin{cases} 1 & \text{if } (v_i, v_j) \in E \\ 0 & \text{otherwise} \end{cases}$$

If $G = (V, E, W)$

$$a_{ij} = \begin{cases} W(v_i v_j) & \text{if } v_i v_j \in E \\ c & \text{otherwise} \end{cases}$$

Graphs and Digraphs

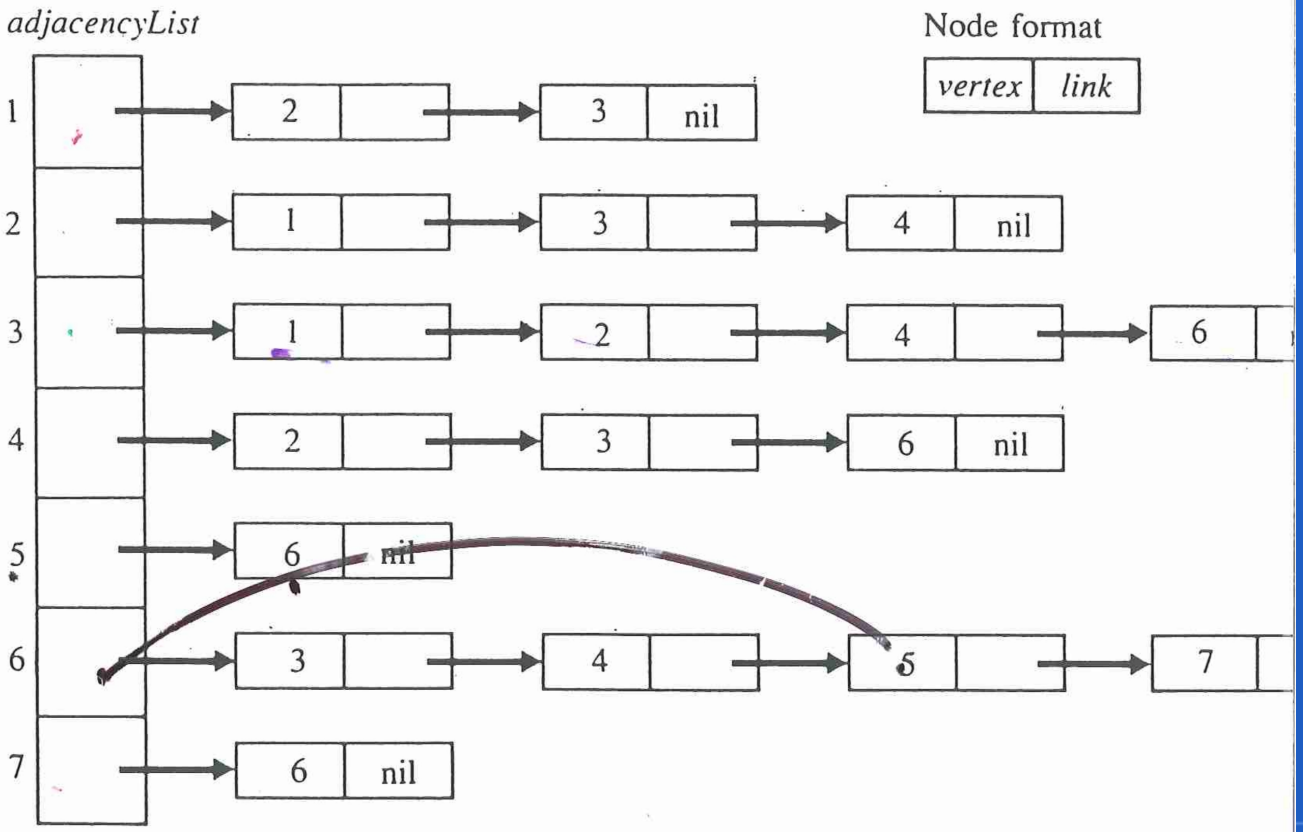
```
19  class DiGraph {
20
21  private int           size;    //number of vertices
22  private boolean[][] AdjMtrx; //adjacency matrix
23  private boolean[]    mark;    //to mark "visited"
```



(a) A graph

$$\begin{pmatrix} 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

(b) Its adjacency matrix.

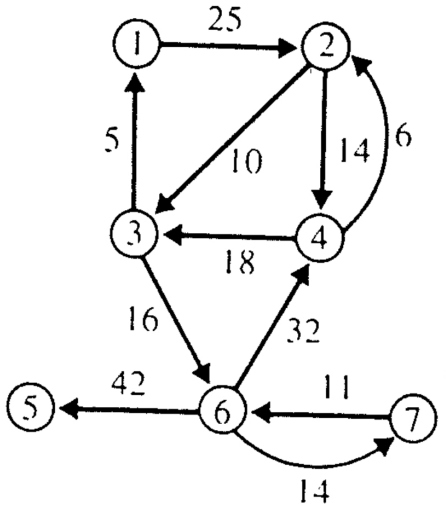


(c) Its adjacency list structure.

Figure 4.10 Representations for a graph.

Graphs and Digraphs

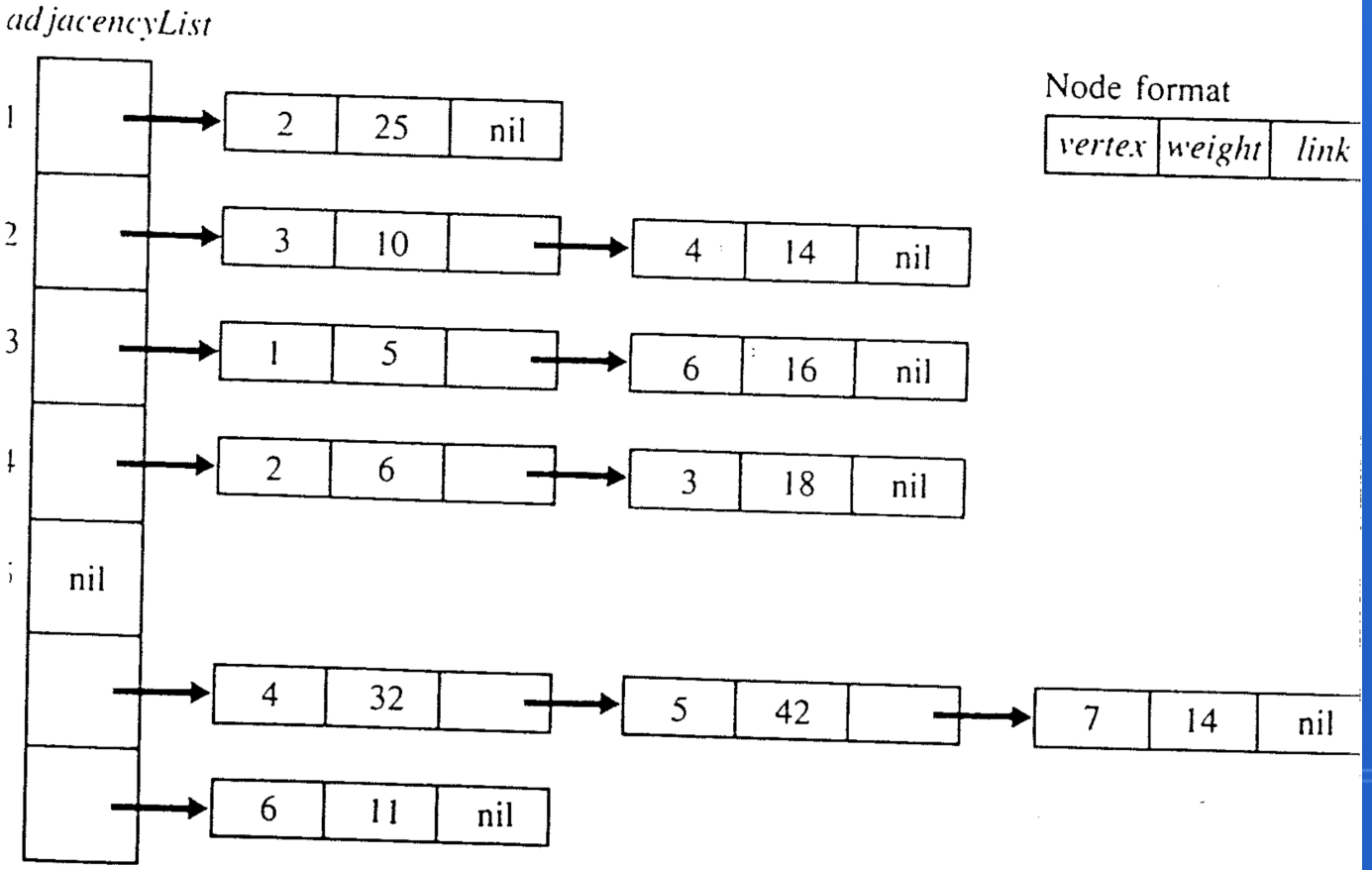
```
18   class DiGraph {
19
20   private int          size;     //number of vertices
21   private LIST[] AdjLists; //array of adjacency lists
22   private boolean[]    mark;     //to mark "visited"
```



(a) A weighted digraph.

$$\begin{pmatrix} 0 & 25 & \infty & \infty & \infty & \infty & \infty \\ \infty & 0 & 10 & 14 & \infty & \infty & \infty \\ 5 & \infty & 0 & \infty & \infty & 16 & \infty \\ \infty & 6 & 18 & 0 & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & 0 & \infty & \infty \\ \infty & \infty & \infty & 32 & 42 & 0 & 14 \\ \infty & \infty & \infty & \infty & \infty & 11 & 0 \end{pmatrix}$$

(b) Its adjacency matrix.



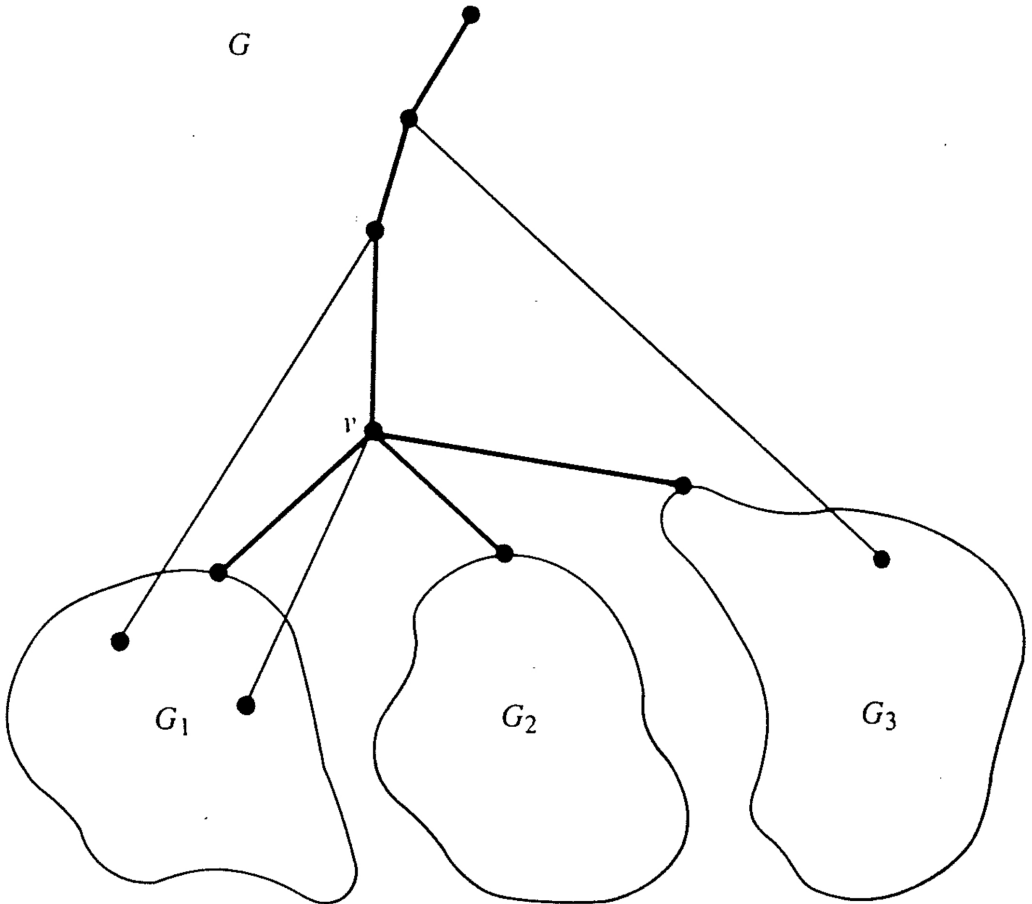
(c) Its adjacency list structure.

Traversing Graphs and Digraphs

Depth-first and Breadth-first Searches

Depth-first search is a generalization of preorder traversal of trees.

In a breadth-first search, vertices are visited in order of increasing distance from the starting point



Graphs and Digraphs

```
138 public static void DFS(DiGraph G, int v) {
139 int w;
140 if (!G.IsVisited(v)) {
141 G.Visit(v);
142 for (w = G.FirstAdj(v); w != -1; w = G.NextAdj(v, w)) DFS(G, w);
143 System.out.println("Back out of " + v);
144 }
145 }
```

```
    void preOrderTraversal(TreeNode T)  {

        Stack  S = new Stack( );                    // let S be an initially e
        TreeNode  N;                                // N points to nodes durin
5
        S.push(T);                                  // push the pointer T onto the em

        while ( !S.empty( ) ) {

0           N = (TreeNode)S.pop( );                 // pop top pointer

            if (N != null) {
                System.out.print(N.info);           // print N
                S.push(N.rlink);                    // push the right poi
5               S.push(N.llink);                    // push the left poi
            }

        }
    }
```

Graphs and Digraphs

```
268  public static void DFSnrec(DiGraph G, int v) {
269      //Adjacency is reversed by pushing on a stack
270      if (!G.IsVisited(v)) {
271          STACK S = new STACK();
272          S.push(v);
273          while (!S.isEmpty())
274          {
275              v = S.pop();
276              G.Visit(v);
277              for (int w = G.LastAdj(v); w != -1; w = G.PreviousAdj(v, w))
278                  if (!G.IsVisited(w)) S.push(w);
279          }
280      }
281  }
```

```
    void levelOrderTraversal(TreeNode T)  {

        Queue Q = new Queue( );                    // let Q be a
        TreeNode N;                                // N points to
5
        Q.insert(T);                               // insert the

        while ( ! Q.empty( ) ) {

10          N = (TreeNode) Q.remove( );                    // r

            if (N != null ) {
                System.out.print(N.info);
                Q.insert(N.llink)                  // insert
15              Q.insert(N.rlink)                  // insert ri
            }
        }
    }
```

Graphs and Digraphs

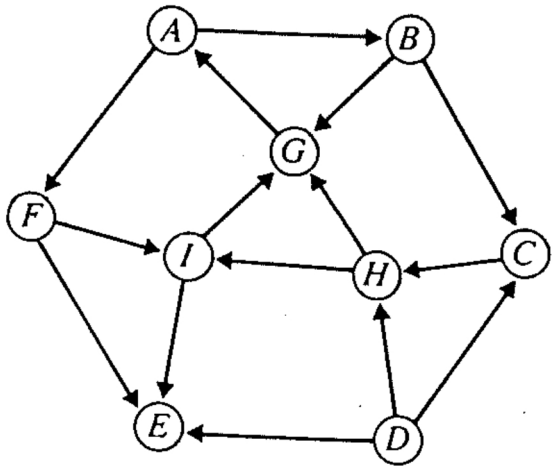
```
177 public static void BFS(DiGraph G, int v) {
178     if (!G.IsVisited(v)) {
179         QUEUE Q = new QUEUE();
180         int w;
181         Q.enqueue(v);
182         while (!Q.isEmpty())
183         {
184             v = Q.dequeue();
185             if (!G.IsVisited(v)) G.Visit(v);
186             for (w = G.FirstAdj(v); w != -1; w = G.NextAdj(v, w))
187                 if (!G.IsVisited(w)) Q.enqueue(w);
188         }
189 }
190 }
```

Graphs and Digraphs

```
177 public static void BFS(DiGraph G, int v) {
178     if (!G.IsVisited(v)) {
179         QUEUE Q = new QUEUE();
180         int w;
181         Q.enqueue(v);
182         while (!Q.isEmpty())
183         {
184             v = Q.dequeue();
185             if (!G.IsVisited(v)) G.Visit(v);
186             for (w = G.FirstAdj(v); w != -1; w = G.NextAdj(v, w))
187                 if (!G.IsVisited(w)) Q.enqueue(w);
188         }
189 }
190 }
```

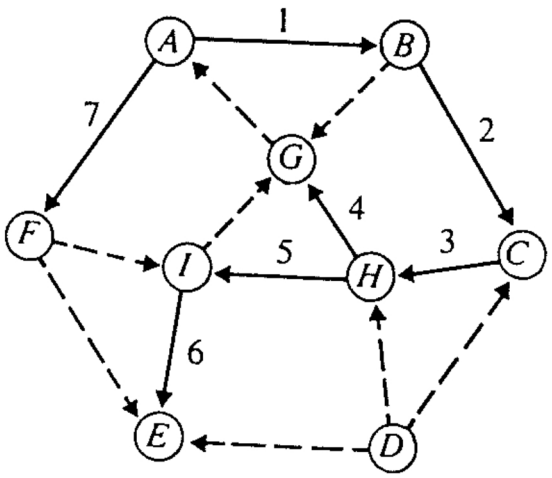
Graphs and Digraphs

```
177 public static void BFS(DiGraph G, int v) {
178     if (!G.IsVisited(v)) {
179         QUEUE Q = new QUEUE();
180         int w;
181         Q.enqueue(v);
182         while (!Q.isEmpty())
183         {
184             v = Q.dequeue();
185             if (!G.IsVisited(v)) G.Visit(v);
186             for (w = G.FirstAdj(v); w != -1; w = G.NextAdj(v, w))
187                 if (!G.IsVisited(w)) Q.enqueue(w);
188         }
189 }
190 }
```

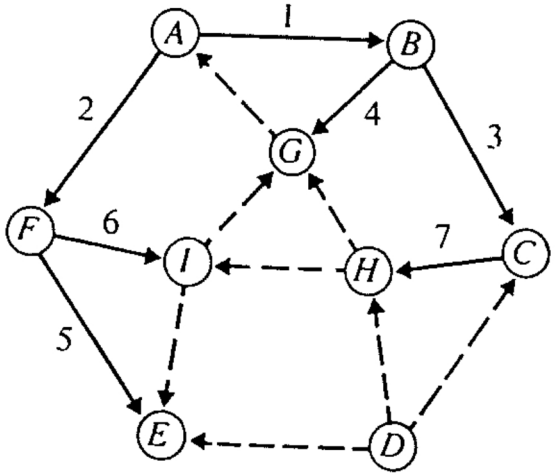


(a) A digraph.

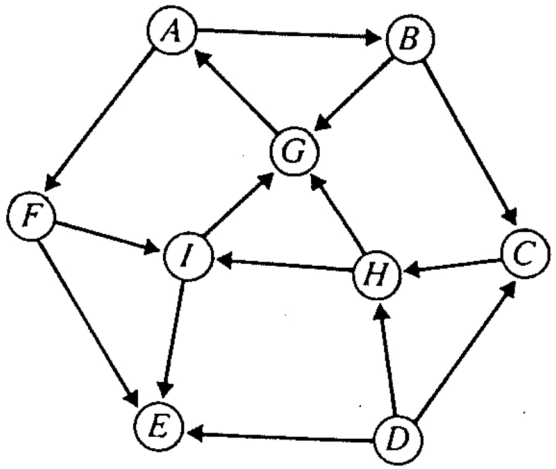
Edges are numbered in the order traversed.



(b) Depth-first search beginning at *A*; order in which vertices are visited: *A B C H G I E F*

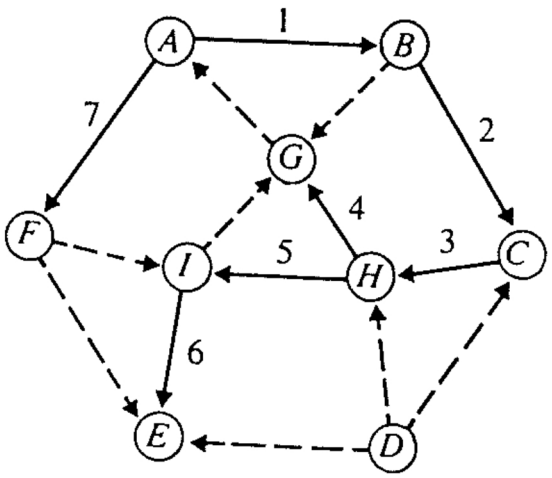


(c) Breadth-first search beginning at *A*; order in which vertices are visited: *A B F C G E I H*

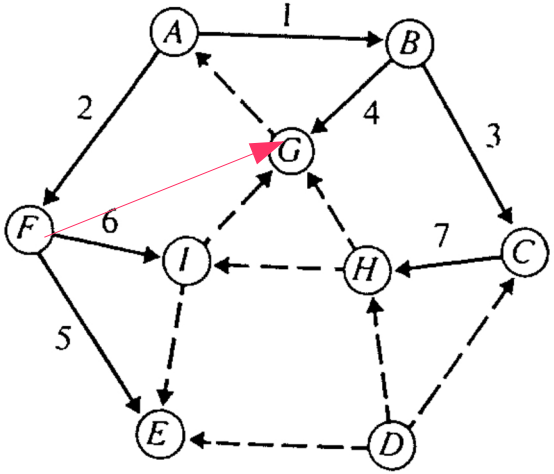


(a) A digraph.

Edges are numbered in the order traversed.

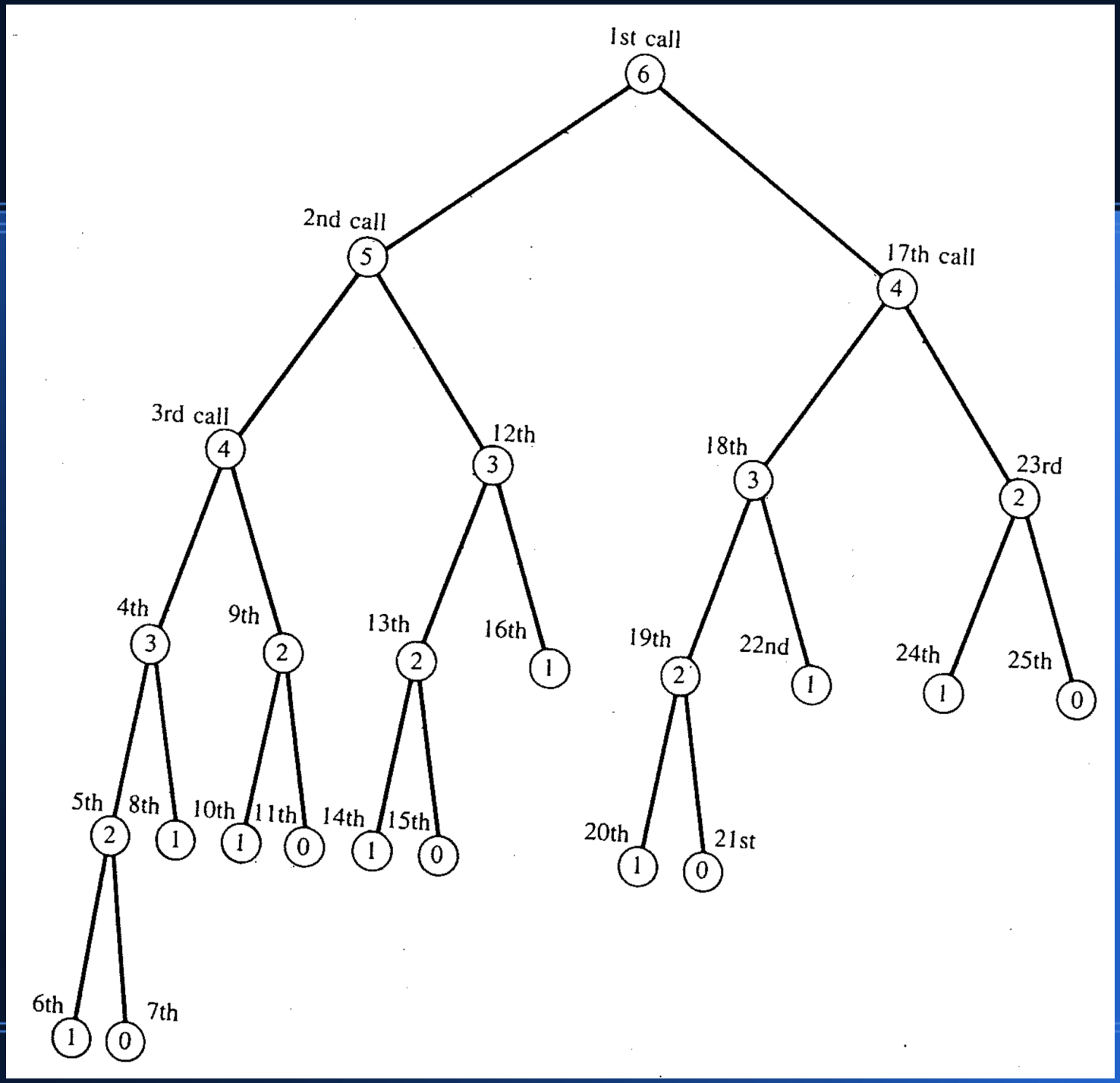


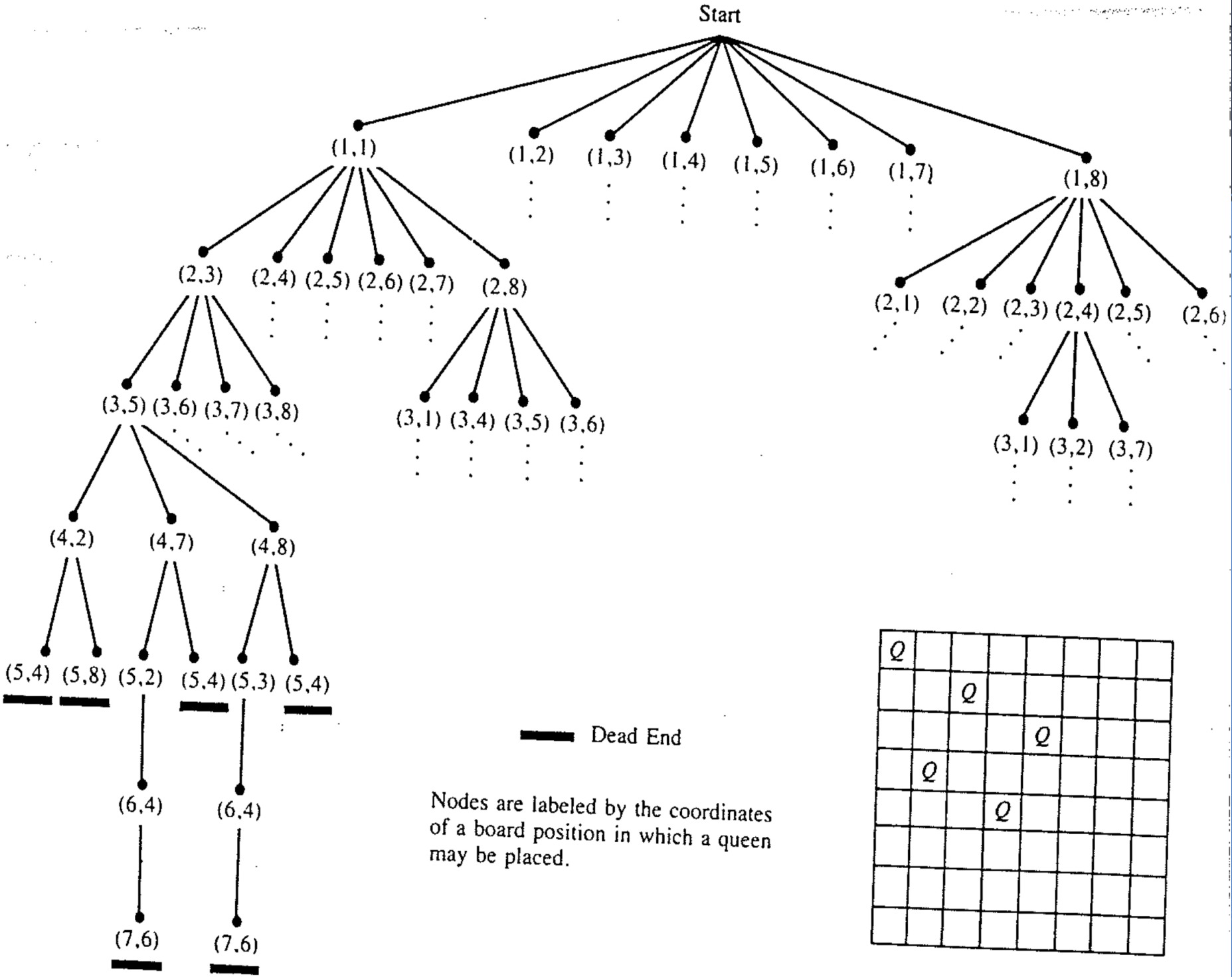
(b) Depth-first search beginning at A; order in which vertices are visited: $A\ B\ C\ H\ G\ I\ E\ F$

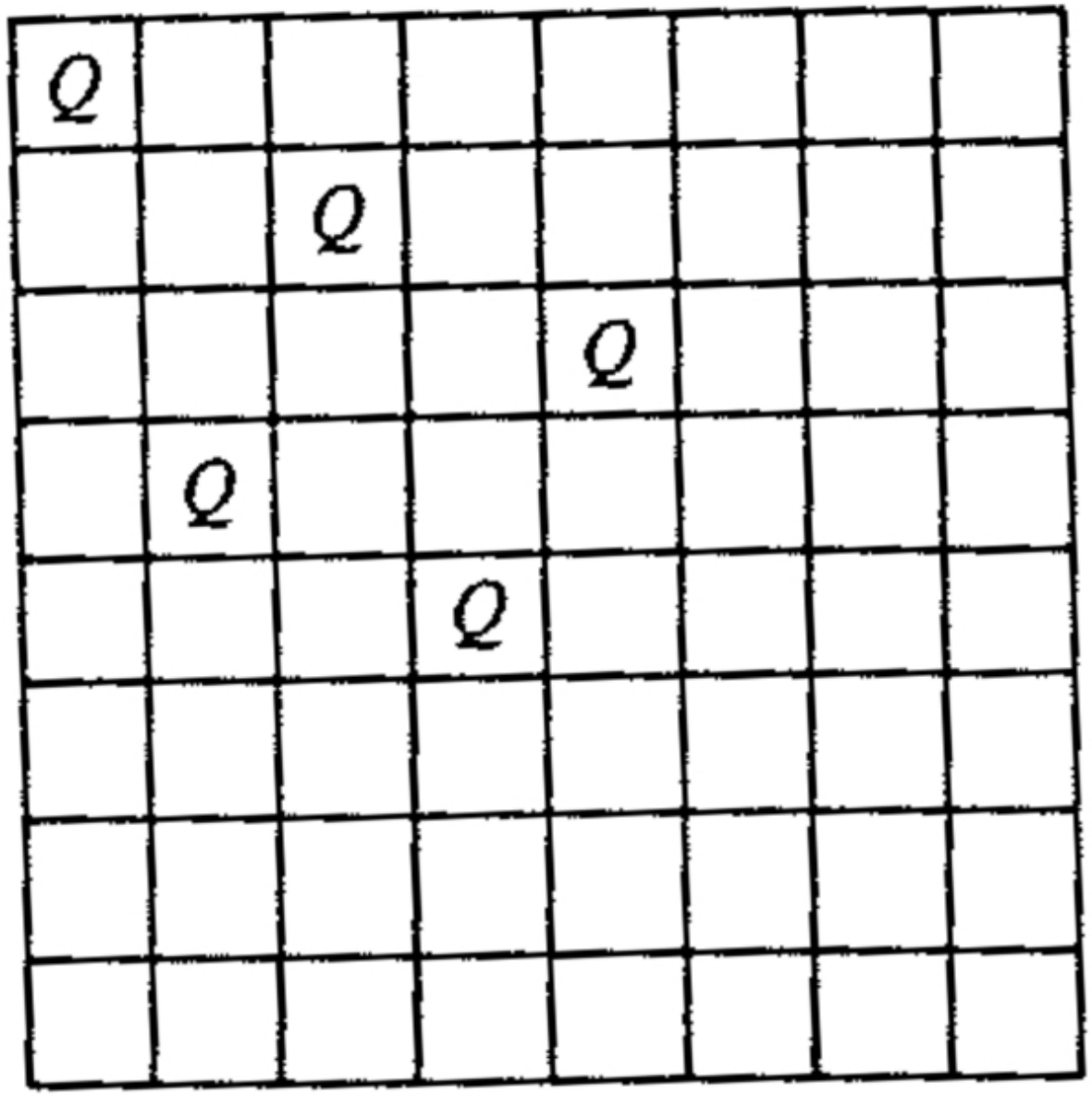


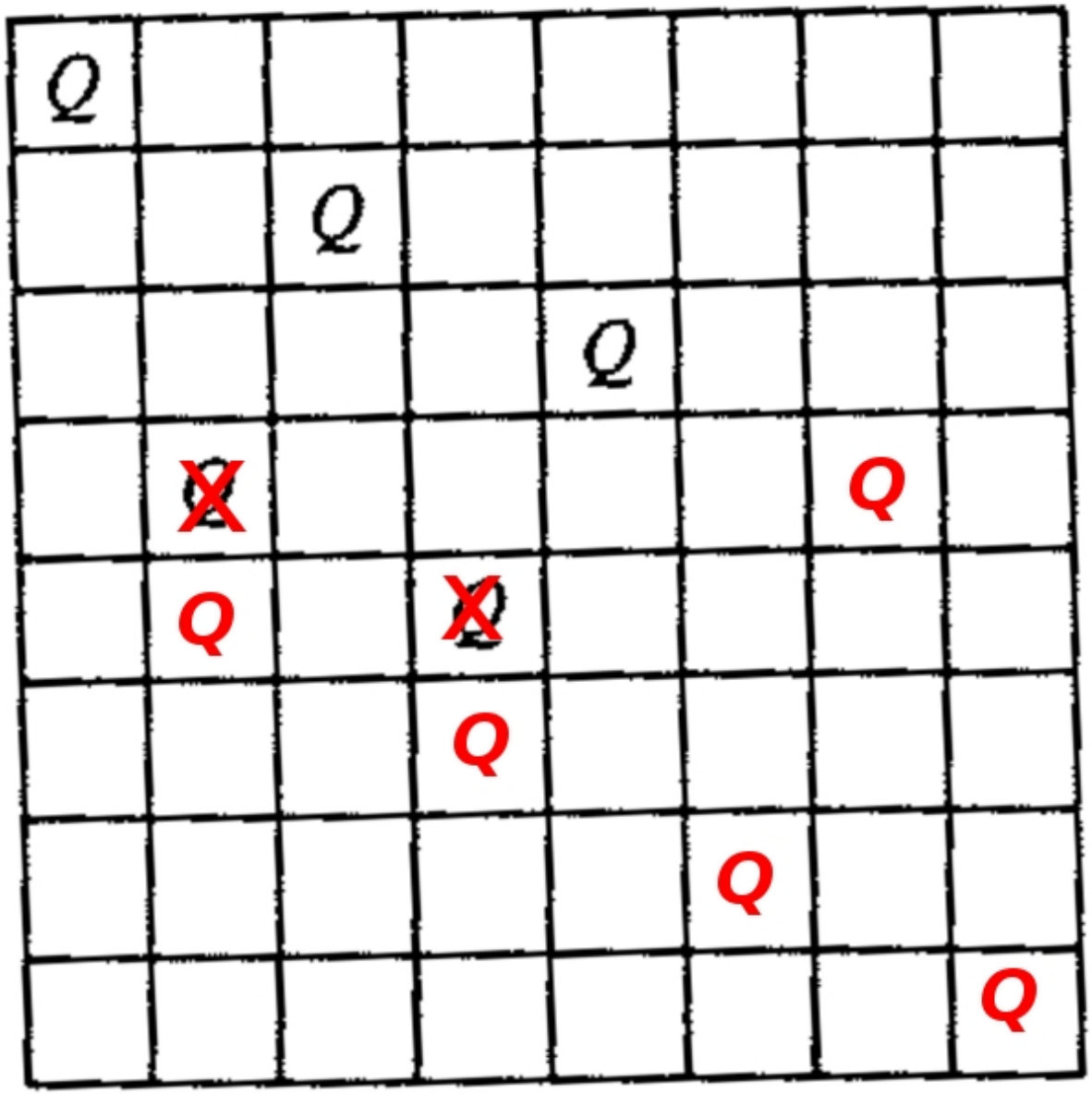
Variant

(c) Breadth-first search beginning at A; order in which vertices are visited: $A\ B\ F\ C\ G\ E\ I\ H$

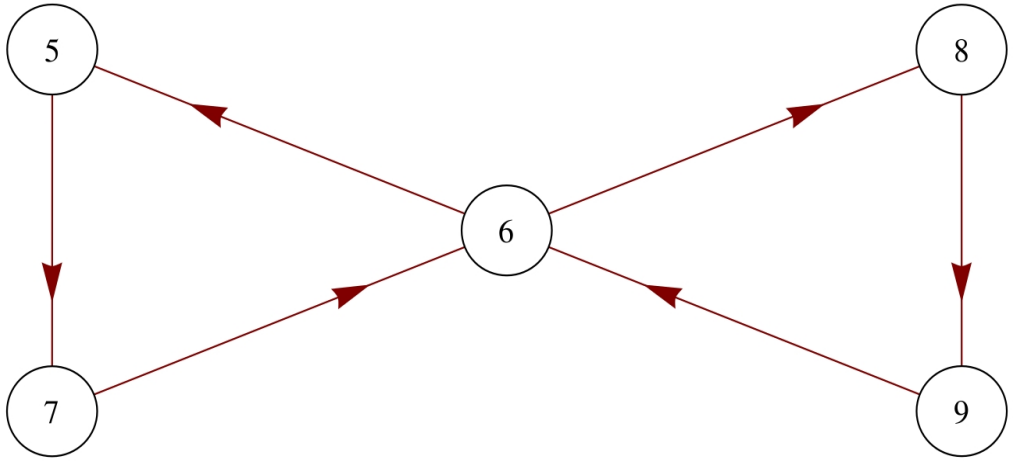

1st call
6
2nd call
5
17th call
4
3rd call
4
12th
3
18th
3
23rd
2
4th
3
9th
2
13th
2
16th
1
19th
2
22nd
1
24th
1
25th
0
5th
2
8th
1
10th
1
11th
0
14th
1
15th
0
20th
1
21st
0
6th
1
7th
0

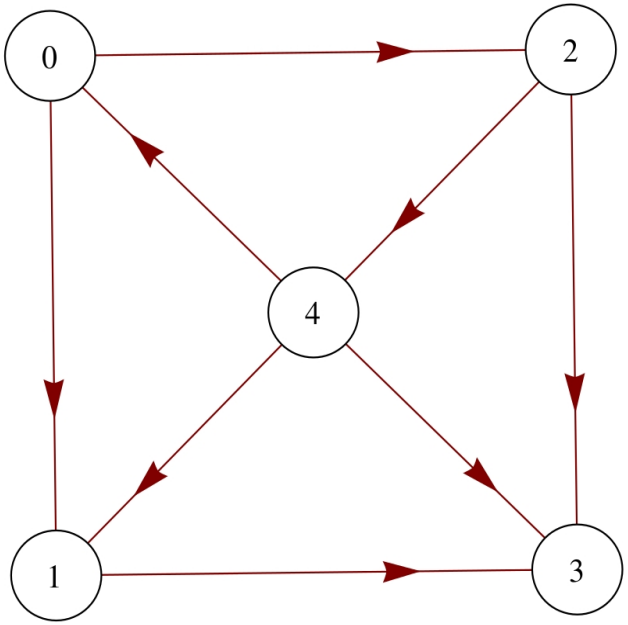

Start
(1,1)
(1,2)
(1,3)
(1,4)
(1,5)
(1,6)
(1,7)
(1,8)
(2,3)
(2,4)
(2,5)
(2,6)
(2,7)
(2,8)
(2,1)
(2,2)
(2,3)
(2,4)
(2,5)
(2,6)
(3,5)
(3,6)
(3,7)
(3,8)
(3,1)
(3,4)
(3,5)
(3,6)
(3,1)
(3,2)
(3,7)
(4,2)
(4,7)
(4,8)
(5,4)
(5,8)
(5,2)
(5,4)
(5,3)
(5,4)
(6,4)
(6,4)
(7,6)
(7,6)
Dead End
Nodes are labeled by the coordinates of a board position in which a queen may be placed.
Q
Q
Q
Q
Q



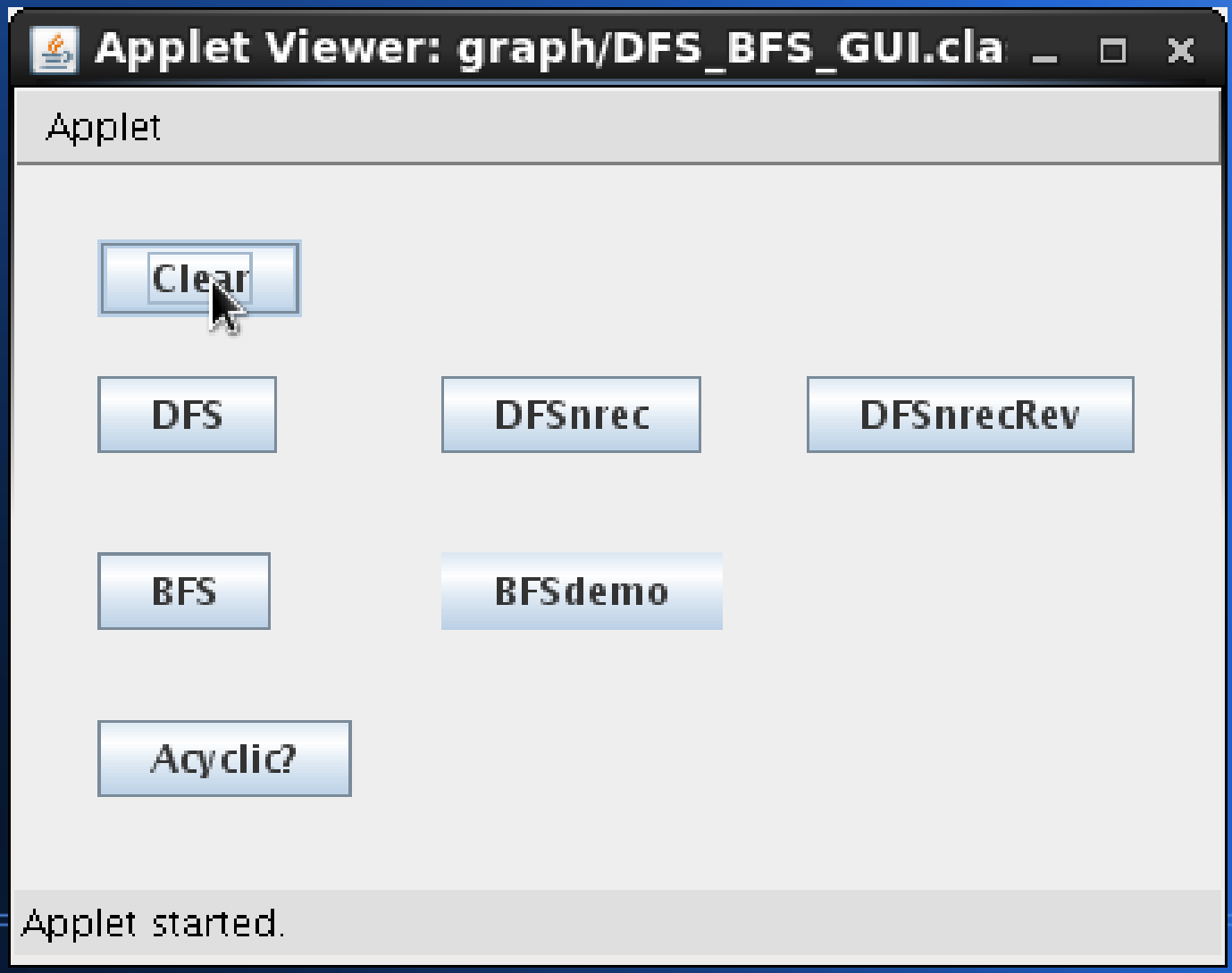

Q
Q
Q
Q
Q
Q
Q
Q
Q
Q

Demo of Graph progr. in java

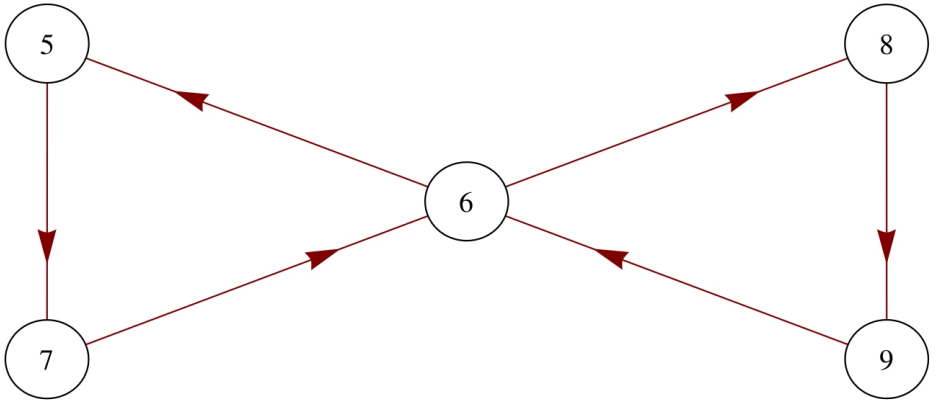


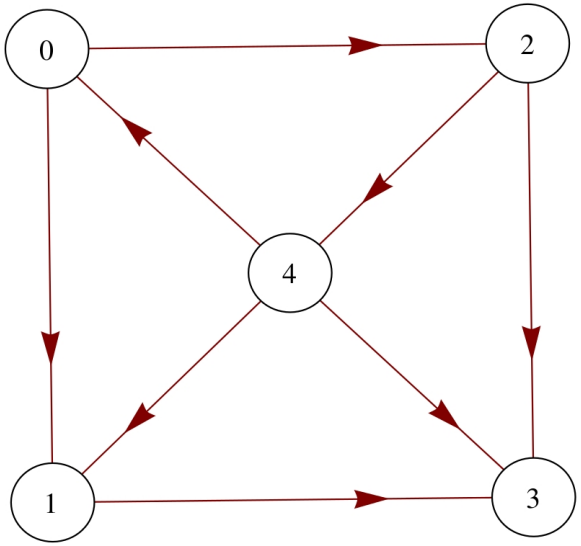


DFS & BFS

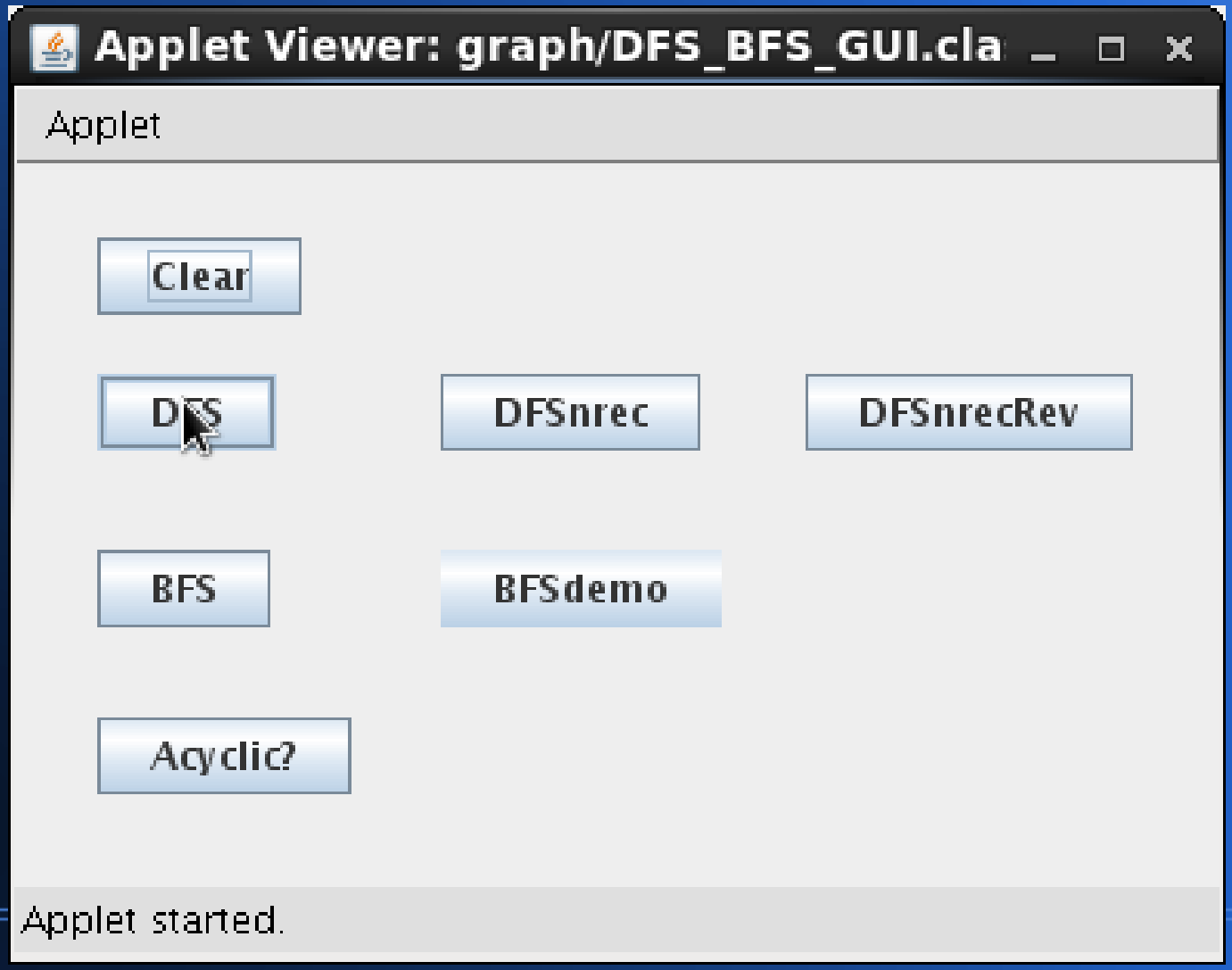

Applet Viewer: graph/DFS_BFS_GUI.cla
Applet
Clear
DFS
DFSnrec
DFSnrecRev
BFS
BFSdemo
Acyclic?
Applet started.

```
compile-single:
run-applet:
Dump start
AdjList[0]: 1 2
AdjList[1]: 3
AdjList[2]: 3 4
AdjList[3]:
AdjList[4]: 0 1 3
AdjList[5]: 7
AdjList[6]: 5 8
AdjList[7]: 6
AdjList[8]: 9
AdjList[9]: 6
Dump end
```



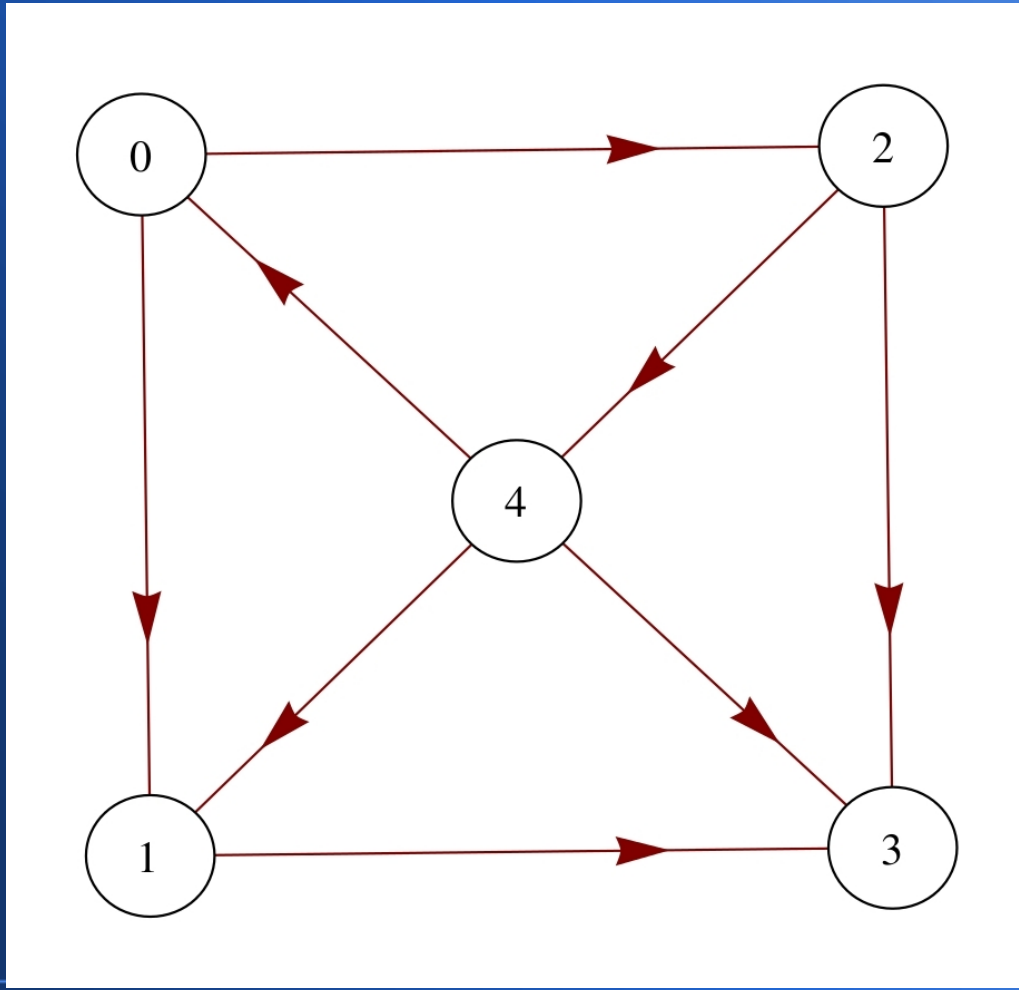


DFS & BFS

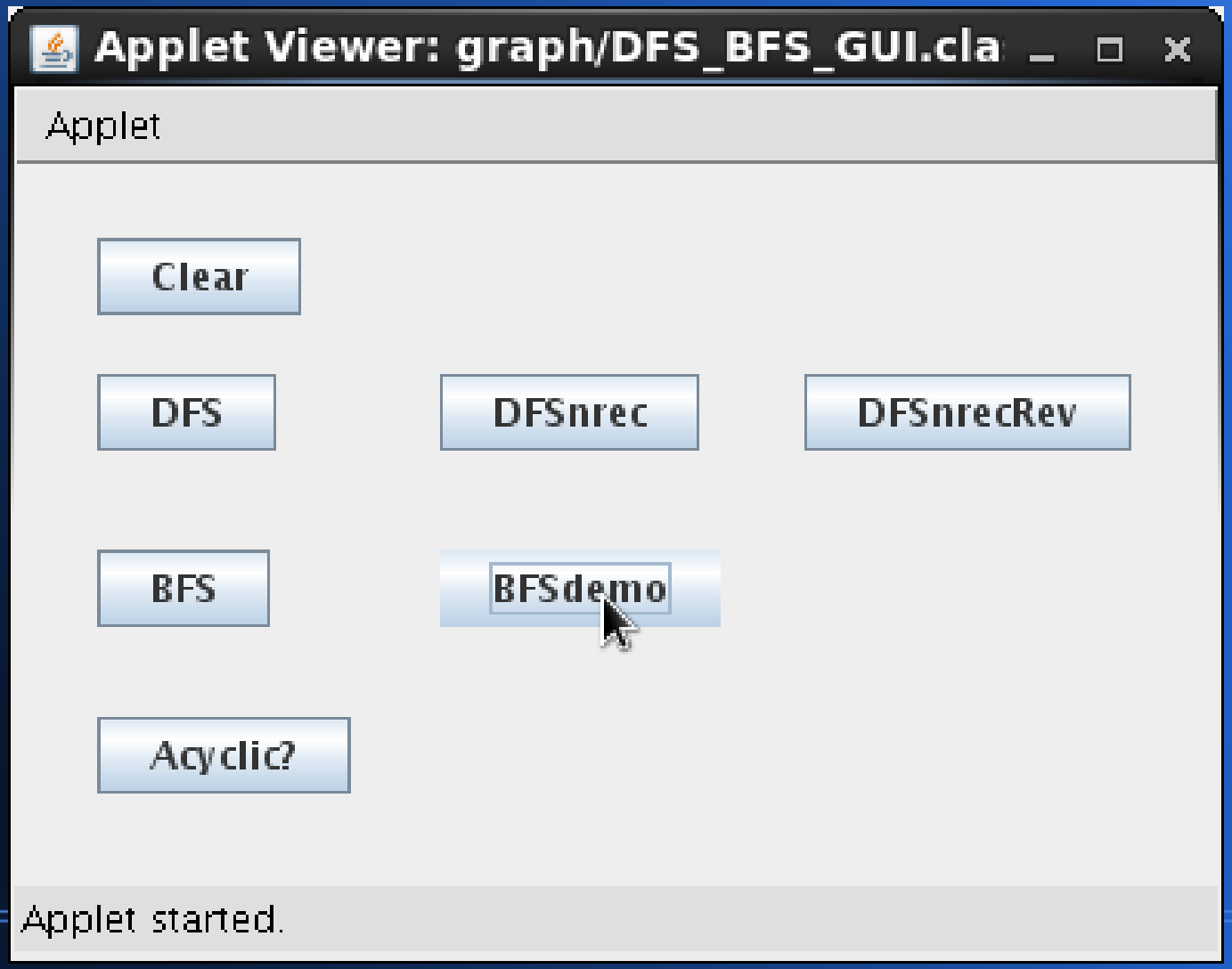

Applet Viewer: graph/DFS_BFS_GUI.cla
Applet
Clear
DFS
DFSnrec
DFSnrecRev
BFS
BFSdemo
Acyclic?
Applet started.

DFS & BFS

```
0 1 3 Back out of 3
Back out of 1
2 4 Back out of 4
Back out of 2
Back out of 0
```

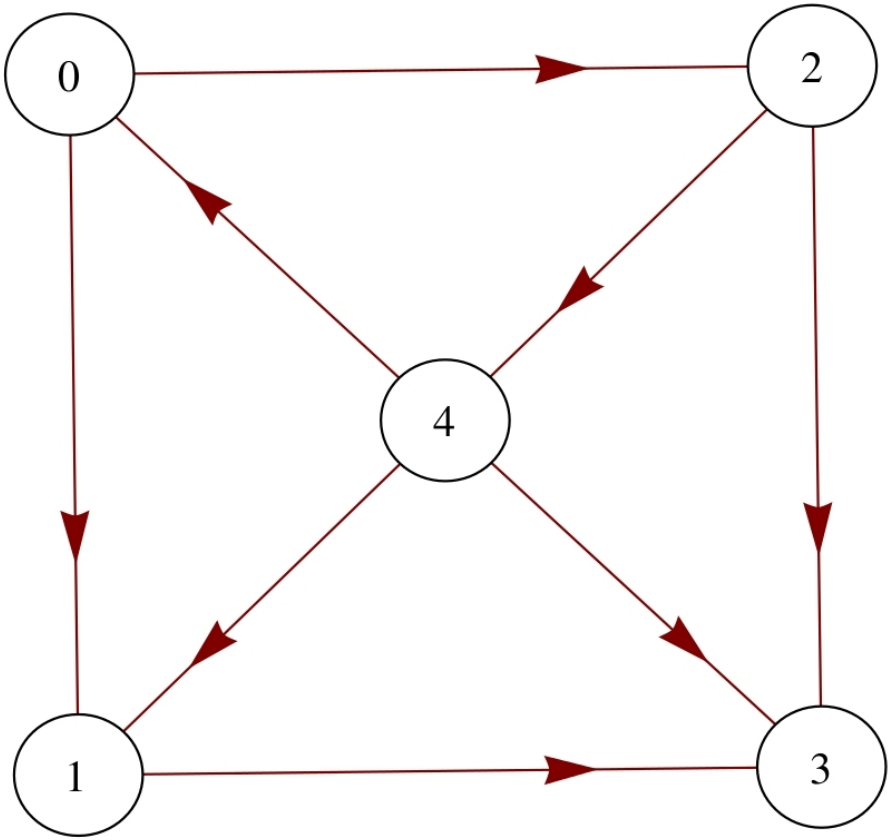


DFS & BFS

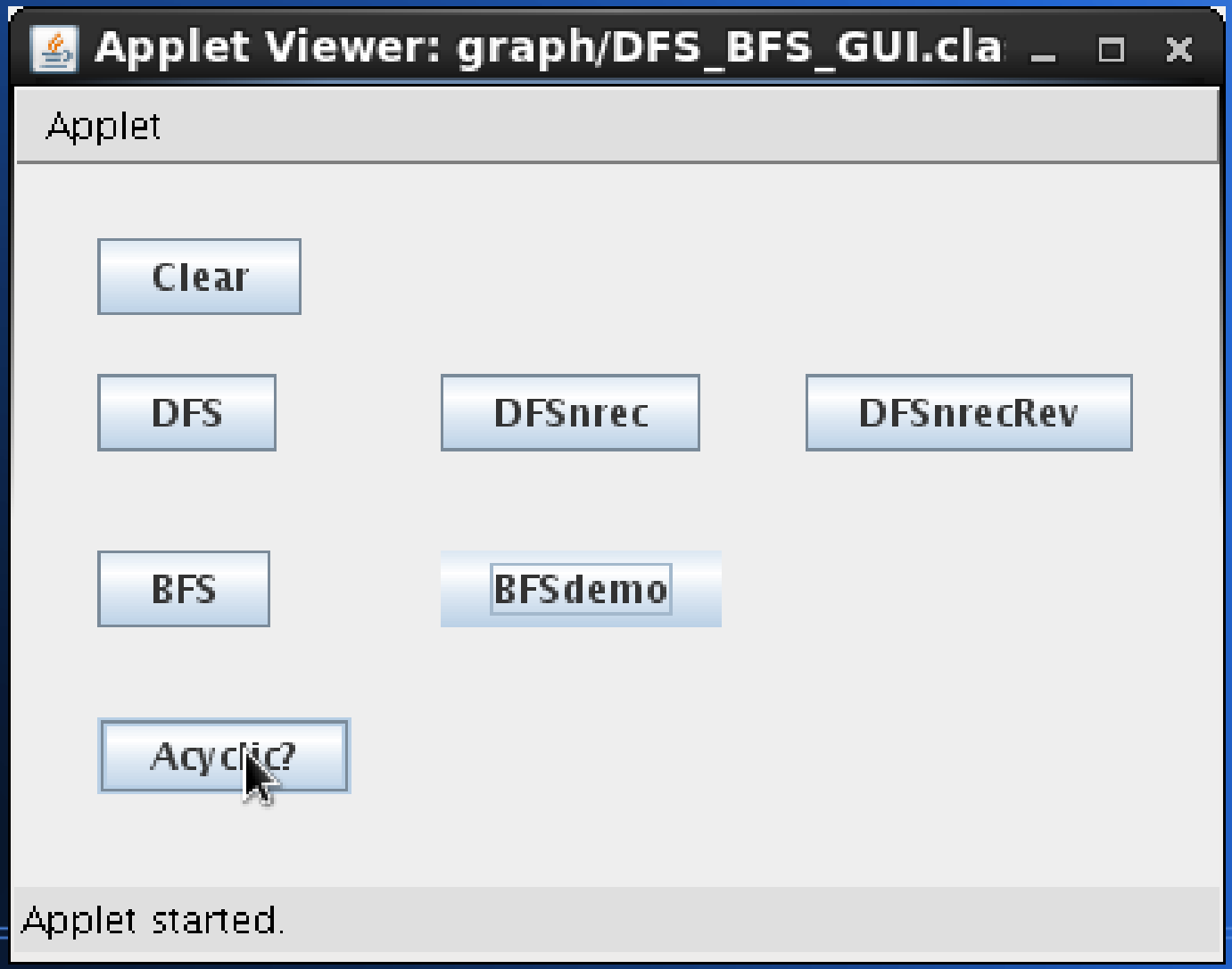

Applet Viewer: graph/DFS_BFS_GUI.cla
Applet
Clear
DFS
DFSnrec
DFSnrecRev
BFS
BFSdemo
Acyclic?
Applet started.

DFS & BFS

```
Level 0: 0
Level 1: 1 2
Level 2: 3 4
```



DFS & BFS

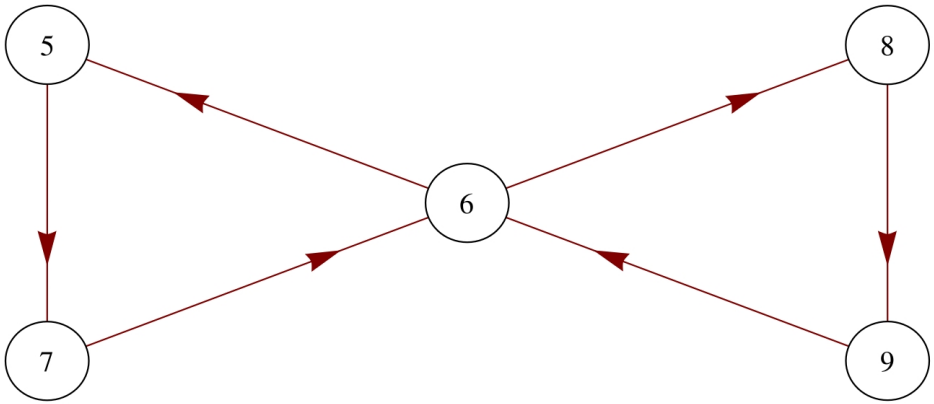

Applet Viewer: graph/DFS_BFS_GUI.cla
Applet
Clear
DFS
DFSnrec
DFSnrecRev
BFS
BFSdemo
Acyclic?
Applet started.

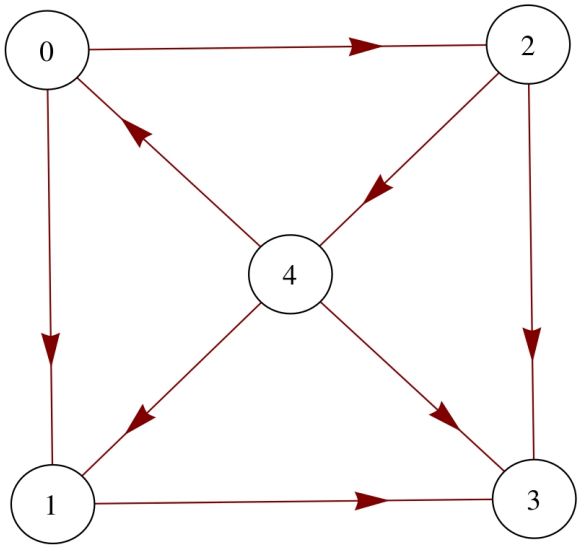
DFS & BFS

```
0 1 3 Back out of 3
Back out of 1
2 4
A cycle: 0 2 4  there is an edge from here to 0
Back out of 4
Back out of 2
Back out of 0
5 7 6
A cycle: 5 7 6  there is an edge from here to 5
8 9
A cycle: 6 8 9  there is an edge from here to 6
Back out of 9
Back out of 8
Back out of 6
Back out of 7
Back out of 5
```

DFS & BFS

```
0 1 3 Back out of 3
Back out of 1
2 4
A cycle: 0 2 4  there is an
Back out of 4
Back out of 2
Back out of 0
5 7 6
A cycle: 5 7 6  there is an
8 9
A cycle: 6 8 9  there is an
Back out of 9
Back out of 8
Back out of 6
Back out of 7
Back out of 5
```





This all will be covered on Final