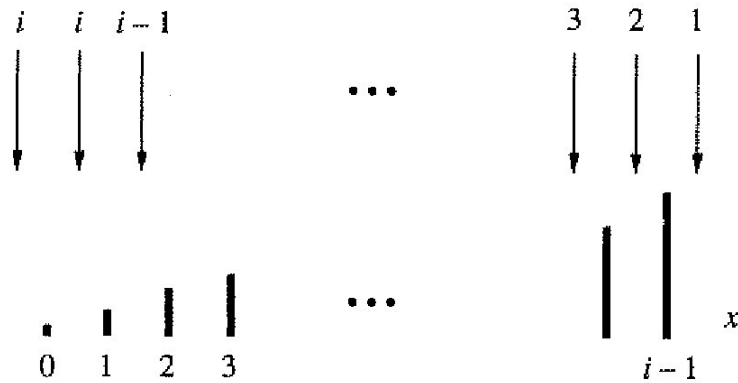


# Average preformance of InsertionSort

Average number of comparisons while inserting  $(i + 1)$  st element  $x$  to a sorted sequence



**Figure 4.5** Number of comparisons needed to determine the position for  $x$

$$E[0] < E[1] < \dots < E[i-1] \quad \leftarrow x$$

of  $i$  elements

The algorithm

Try at index  $i$ :

$$E[0] < E[1] < \dots < E[i-1] < \underset{\substack{\uparrow? \\ x}}{\_}$$

If yes then

$$E[0] < E[1] < \dots < E[i-1] < x$$

(done after 1 comparison)

If no then try at index  $(i-1)$ :

$$E[0] < E[1] < \dots < E[i-2] < \underset{\substack{\uparrow? \\ x}}{\_} < E[i]$$

If yes then

$$E[0] < E[1] < \dots < E[i-2] < x < E[i]$$

(done after 2 comps)

If no then try at index  $(i-2)$  :

$$E[0] < E[1] < \dots < E[i-3] < \underset{\substack{\uparrow? \\ x}}{x} < E[i-1] < E[i]$$

If yes then

$$E[0] < E[1] < \dots < E[i-3] < x < E[i-1] < E[i]$$

(done after 3 comps)

etc.

...

$$E[0] < E[1] < \dots < E[j-1] < x < E[j+1] < \dots < E[i-1] < E[i]$$

(done after  $i-j+1$  comps ; we will derive this formula below)

etc.

...

If no then try at index 1 :

$$E[0] < \underset{\substack{\uparrow? \\ x}}{x} < E[2] < \dots < E[i]$$

If yes then

$$E[0] < x < E[2] < \dots < E[i]$$

(done after  $i$  comps)

If no then

$$x < E[1] < E[2] < \dots < E[i]$$

(done after **how many** comps?)

Let

- $\Pr(j)$  be the probability that the index of  $x$  after insertion in

$$E[0] < E[1] < \dots < E[j-1] < \underset{\substack{\uparrow \\ x}}{\_} < E[j+1] < \dots < E[i-1] < E[i]$$

is  $j$ , where  $j \in \{0, \dots, i\}$

- $\# \text{comp}(j)$  be the number of comparisons necessary for the program to determine that  $x$  should be inserted at index  $j$

We have (assuming that all

$i+1$  possible points of insertion are equally likely,  
see Figure 4.5 in your textbook) :

$$\Pr(j) = \frac{1}{(i+1)}$$

When  $j$  gets **decremented** after each "no" in the above algorithm,

$\# \text{comp}(j)$  gets **incremented**, except for the last step when  $\# \text{comps}(1) = \# \text{comps}(0) = i$ .

Hence,

$$j + \# \text{comp}(j) = \text{constant}$$

for all  $0 < j \leq i$ .

Putting  $j = i$  (the first step of the algorithm) yields

$$i + 1 = \text{constant}$$

because  $\# \text{comps}(i) = 1$ .

Therefore, for all  $0 < j \leq i$ ,

$$j + \# \text{comp}(j) = i + 1$$

or, (subtracting  $j$  from both sides)

$$\# \text{comp} (j) = i - j + 1.$$

So,

$$\# \text{comp} (j) = \begin{cases} i - j + 1 & \text{for } 0 < j \leq i \\ i & \text{for } j = 0 \end{cases}$$

The average number of comparisons while inserting  $(i + 1)$  st element (where  $i \geq 0$ ) is :

$$\begin{aligned} \sum_{j=0}^i \text{Pr} (j) \times \# \text{comp} (j) &= \text{Pr} (0) \times \# \text{comp} (0) + \sum_{j=1}^i \text{Pr} (j) \times \# \text{comp} (j) = \\ \frac{1}{(i+1)} \times i + \sum_{j=1}^i \frac{1}{(i+1)} \times (i - j + 1) &= \\ \frac{i}{2} + \frac{i}{1+i} &= \\ \frac{i}{2} + \frac{i}{1+i} = \frac{i}{2} + \frac{i+1}{1+i} - \frac{1}{1+i} = \frac{i}{2} + 1 - \frac{1}{1+i} \end{aligned}$$

Verification of the above

$$\begin{aligned} \frac{i}{2} + \text{Factor} \left[ 1 - \frac{1}{1+i} \right] &= \\ \frac{i}{2} + \frac{i}{1+i} &= \end{aligned}$$

The average number of comparisons for all insertions is

(remember that  $\text{Pr} (0) \times \# \text{comp} (0) = 0$  so the first insertion at  $i = 0$  may be ignored) :

$$\begin{aligned} T_{\text{avg}} (n) &= \sum_{i=1}^{n-1} \sum_{j=0}^i \text{Pr} (j) \times \# \text{comp} (j) = \\ &= \sum_{i=1}^{n-1} \left( \frac{i}{2} + \frac{i}{1+i} \right) = \\ &= \sum_{i=1}^{n-1} \left( \frac{i}{2} + 1 - \frac{1}{1+i} \right) = \sum_{i=1}^{n-1} \left( \frac{i}{2} + 1 \right) - \sum_{i=1}^{n-1} \frac{1}{1+i} \end{aligned}$$

$$\sum_{i=1}^{n-1} \frac{1}{1+i} = (\text{put } j = i+1; j \in \{2, \dots, n\})$$

$$= \sum_{j=2}^n \frac{1}{j} = \sum_{j=1}^n \frac{1}{j} - \frac{1}{1} = \sum_{j=1}^n \frac{1}{j} - 1.$$

So,

$$T_{\text{avg}}(n) = \sum_{i=1}^{n-1} \left( \frac{i}{2} + 1 \right) - \left( \sum_{j=1}^n \frac{1}{j} - 1 \right) = \sum_{i=1}^{n-1} \left( \frac{i}{2} + 1 \right) + 1 - \sum_{j=1}^n \frac{1}{j}.$$

$$\sum_{i=1}^{n-1} \left( \frac{i}{2} + 1 \right) + 1$$

$$1 + \frac{1}{4} (-4 + 3n + n^2)$$

$$\text{Expand} \left[ 1 + \frac{1}{4} (-4 + 3n + n^2) \right]$$

$$\frac{3n}{4} + \frac{n^2}{4}$$

$$\text{FullSimplify} \left[ \frac{3n}{4} + \frac{n^2}{4} \right]$$

$$\frac{1}{4} n (3 + n)$$

$$\text{Thus } \sum_{i=1}^{n-1} \left( \frac{i}{2} + 1 \right) + 1 = \frac{1}{4} n (n + 3)$$

Thus

$$T_{\text{avg}}(n) = \frac{1}{4} n (3 + n) - \sum_{j=1}^n \frac{1}{j}$$

Recall that

$$\sum_{j=1}^n \frac{1}{j} \approx \text{Log}[n] + 0.577216$$

(Euler's formula for harmonic number, see file Summationa.nb)

Hence,

$$T_{\text{avg}}(n) \approx \frac{1}{4} n (n + 3) - \text{Log}[n] - 0.577216 =$$

$$-0.577216 + 0.75n + 0.25n^2 - \text{Log}[n] \in \Theta(n^2)$$

**Note.** The difference between  $\frac{1}{4} n (n + 3)$  and textbook's

lower bound for average number of comps  $\frac{1}{4} n (n - 1)$  is  $n$ .

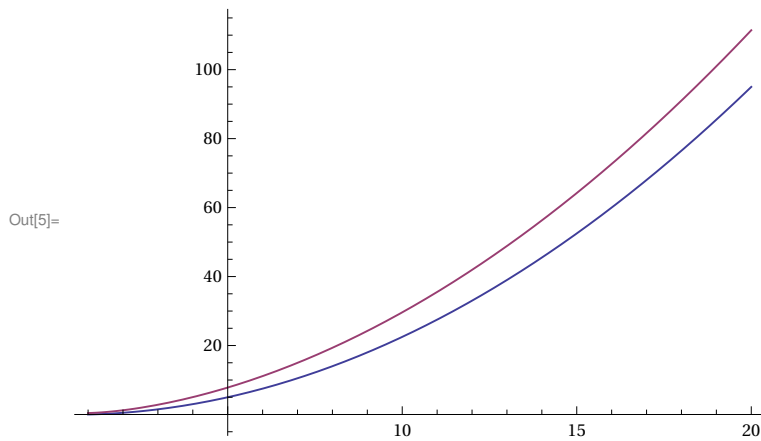
$$\text{Simplify}\left[\frac{1}{4}n(n+3) - \frac{1}{4}n(n-1)\right]$$

$n$

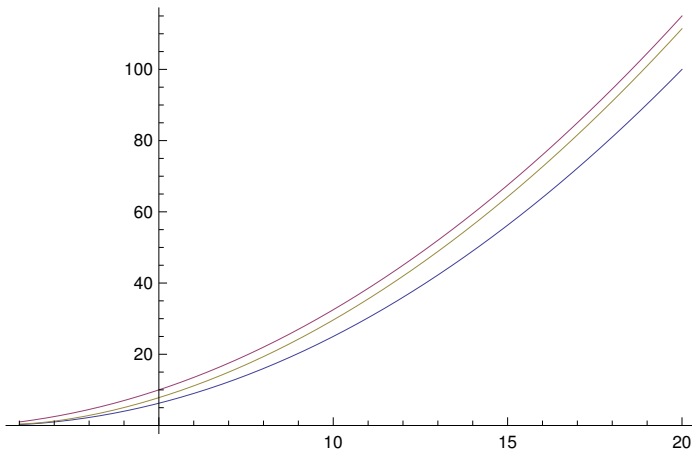
Thus the difference between the upper bound and textbook's lower bound is

$$n - \sum_{j=1}^n \frac{1}{j} \approx n - (\text{Log}[n] + 0.577216)$$

In[5]:= Plot[Tooltip[{ $\frac{1}{4}n(n-1)$ ,  $\frac{1}{4}n(3+n) - \text{Log}[n] - 0.577216$ }],  
{n, 1, 20}, PlotTheme → "Classic"]



Plot[Tooltip[{ $0.25n^2$ ,  $\frac{1}{4}n(3+n)$ ,  $\frac{1}{4}n(3+n) - \text{Log}[n] - 0.577216$ }], {n, 1, 20}]

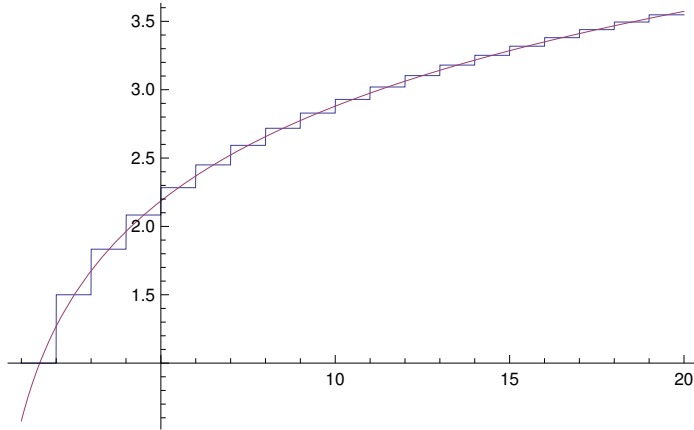


$$\text{Limit}\left[\frac{1}{n^2}(-0.577216 + 0.75n + 0.25n^2 - 1 \cdot \text{Log}[n]), n \rightarrow \infty\right]$$

0.25

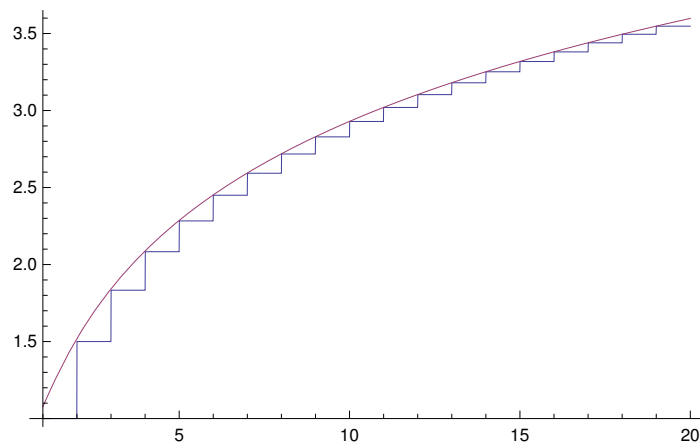
Euler's formula, again:

`Plot[{ $\sum_{j=1}^n \frac{1}{j}$ ,  $\text{Log}[n] + 0.577216$ }, {n, 1, 20}]`



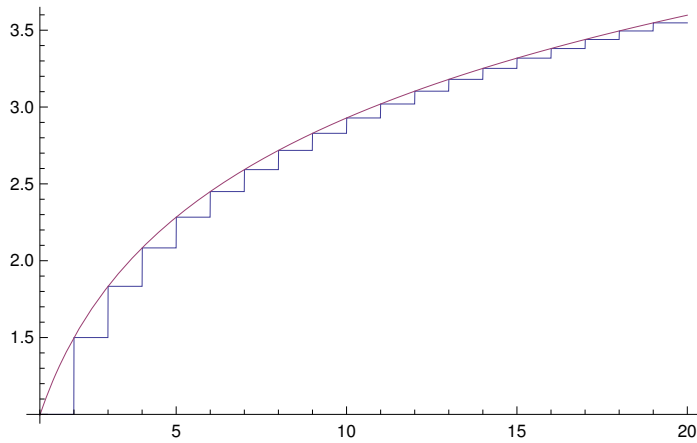
**More accurate :**

`Plot[{ $\sum_{j=1}^n \frac{1}{j}$ ,  $\text{Log}[n] + 0.577216 + \frac{1}{2n}$ }, {n, 1, 20}]`



**Even more accurate**

Plot $\left[\left\{\sum_{j=1}^n \frac{1}{j}, \text{Log}[n] + 0.577216 + \frac{1}{2n} - \frac{1}{12n^2}\right\}, \{n, 1, 20\}\right]$

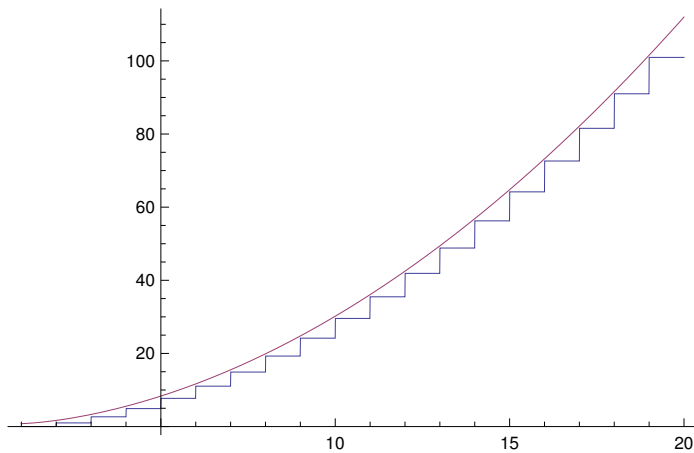


They yield more precise (optional for all students) approximation of the average number of comparisons

$$T_{\text{avg}}(n) \approx \frac{1}{4} n (n+3) - \text{Log}[n] - 0.577216 - \frac{1}{2n} + \frac{1}{12n^2}$$

Here is a visual verification of the above :

Plot $\left[\left\{\sum_{i=1}^{n-1} \left(\frac{i}{2} + \frac{i}{1+i}\right), \frac{1}{4} n (n+3) - \left(\text{Log}[n] + 0.577216 + \frac{1}{2n} - \frac{1}{12n^2}\right)\right\}, \{n, 1, 20\}\right]$



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In[11]:= Plot[ $\left\{\frac{1}{4} n (n - 1), \sum_{i=1}^{n-1} \left(\frac{i}{2} + \frac{i}{1 + i}\right), \frac{1}{4} n (n + 3) - \left(\text{Log}[n] + 0.577216 \left(+ \frac{1}{2 n}\right) - \frac{1}{12 n^2}\right)\right\},$ 
  {n, 1, 20}, PlotTheme -> "Classic"]

```

Out[11]=

