P - a program with input I that runs in time t (I)

 ${\tt D}$ - the set of all valid inputs for P (I)

size (I) - size of input I

 $size : D \rightarrow N$

 $\ensuremath{\text{D}}_n$ - the set of all valid inputs of size n

$$D_n = \{I \in D \mid size(I) = n\}$$

Running time as a function of size of input

Time :
$$\mathbb{N} \to \mathbb{R}^+$$

Worst - case running time

$$T(n) = \max\{t(I) \mid I \in D_n\}$$

Average - case running time

Given probability distribution p_n on D_n ,

$$T_{avg} \text{ (n) } = \sum \left\{ \text{t (I)} \times p_n \text{ (I) } \middle| \text{I } \in \text{D}_n \right\}$$

In the case of uniform distribution of probability:

$$T_{avg}(n) = \frac{\sum \{t(I) \mid I \in D_n\}}{\mid D_n \mid}$$

Example.

QuickSort:

 ${\tt D}\,$ - the set of all permutations of some initial interval of the set of natural numbers ${\tt N}\,$

size (I) = number of elements in I (to be sorted)

 D_n - the set of all permutations of $\{0\,,\ \ldots,\ n\,-\,1\}$

$$|D_n| = n!$$

Worst - case running time:

$$T(n) \sim n^2$$

The worst input of size $\ensuremath{\mathbf{n}}$ is the sequence

$$< 0, \ldots, n-1 > .$$

Average - case running time:

Assume that all inputs of size n to QuickSort are equally likely (with probability $\frac{1}{n!}$).

$$T_{avg}$$
 (n) ~ n log_2 n

 $\texttt{Plot}\big[\texttt{Tooltip}\big[\big\{\texttt{x}\,\texttt{Log}\,\texttt{[x]}\,,\,\,\texttt{x}^2\big\}\big]\,,\,\,\{\texttt{x},\,\,0\,,\,\,10\}\big]$

